

## Lecture 17:

Recall:

Power method:

Step 1: Initialize  $\vec{x}^{(0)}$  (such that  $\vec{x}^{(0)} = a_1 \vec{x}_1 + \dots + a_n \vec{x}_n$   
with  $a_1 \neq 0$ )

Step 2: Compute  $\vec{x}^{(k+1)} = A \vec{x}^{(k)} / \|A \vec{x}^{(k)}\|_\infty$

Step 3: Compute  $\lambda^{(k+1)} \stackrel{\text{def}}{=} \|A \vec{x}^{(k+1)}\|_\infty$

Then:  $\lambda^{(k)} \rightarrow |\lambda_1|$  as  $k \rightarrow \infty$

## Generalization of Power method

Consider an invertible matrix  $A$ . Suppose  $A$  has eigenvalues:

$$|\lambda_1| > |\lambda_2| > \dots > |\lambda_n| (> 0)$$

Consider  $A^{-1}$  (exist as all eigenvalues are non-zero). Then  $A^{-1}$  has eigenvalues:

$$\frac{1}{\lambda_1}, \frac{1}{\lambda_2}, \dots, \frac{1}{\lambda_n} \text{ with } \left| \frac{1}{\lambda_n} \right| > \left| \frac{1}{\lambda_{n-1}} \right| > \dots > \left| \frac{1}{\lambda_1} \right|$$

Extension: Apply Power's method on  $A^{-1}$  to obtain  $\left| \frac{1}{\lambda_n} \right|$ .

$\therefore$  the minimal eigenvalue can be determined! (Inverse Power method)

Remark: Computing  $A^{-1}$  is difficult! We solve:  $A\vec{y} = \vec{x}^{(n)}$  in each iteration to determine  $A^{-1}\vec{x}^{(n)}$ .

Finding  $A^{-1}$  is equivalent to solving:

$$A\vec{y} = \begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix}, A\vec{y} = \begin{pmatrix} 0 \\ 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix}, \dots, A\vec{y} = \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ 1 \end{pmatrix}.$$

$$\underbrace{A \begin{pmatrix} \frac{1}{\lambda_1} \vec{v}_1 & \frac{1}{\lambda_2} \vec{v}_2 & \dots & \frac{1}{\lambda_n} \vec{v}_n \end{pmatrix}}_{A^{-1}} = \begin{pmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{pmatrix}$$

Algorithm: (Inverse Power method)

Step 1: Pick  $\vec{x}^{(0)}$  with  $\|\vec{x}^{(0)}\|_\infty = 1$

Step 2: For  $k = 1, 2, \dots$ , solve  $A\vec{w} = \vec{x}^{(k-1)}$ .

$$\text{Let } \vec{x}^{(k)} = \frac{\vec{w}}{\|\vec{w}\|_\infty}.$$

$$\text{Let } \rho_k = \|A\vec{x}^{(k)}\|_\infty$$

Remark: Again,  $\rho_k \rightarrow |\lambda_n|$  as  $k \rightarrow \infty$ . ( $\vec{x}^{(k)} \approx$  eigenvector of eigenvalue  $\lambda_n$ )

Inverse power method with shift

Goal: Take  $\mu \in \mathbb{R}$ . Find the eigenvalue of  $A$  closest to  $\mu$ .

Observation: Consider  $B = A - \mu I$ . Then  $B$  has eigenvalues:  
 $\{\lambda_1 - \mu, \lambda_2 - \mu, \dots, \lambda_n - \mu\}$

Inverse Power method find eigenvalues such that  $|\lambda_j - \mu|$  is the smallest.

$\therefore \lambda_j$  closest to  $\mu$  can be found.

Algorithm: (Inverse power method with shift)

Step 1: Take  $\mu \in \mathbb{R}$ . Pick  $\vec{x}^{(0)}$  such that  $\|\vec{x}^{(0)}\|_\infty = 1$ .

Step 2: For  $k = 1, 2, \dots$

Solve:  $(A - \mu I) \vec{w} = \vec{x}^{(k-1)}$  for  $\vec{w}$ .

Let:  $\vec{x}^{(k)} = \frac{\vec{w}}{\|\vec{w}\|_\infty}$ .

Let  $\rho_k = \|A \vec{x}^{(k)}\|_\infty$  ( $\rho_k \rightarrow |\lambda_j|$  as  $k \rightarrow \infty$ )



Recall:

Power's method reads:  $\vec{x}^{(k+1)} = \frac{A \vec{x}^{(k)}}{\|A \vec{x}^{(k)}\|_\infty}$  for  $k=0, 1, \dots$

$$\Rightarrow \vec{x}^{(k)} = \frac{A^k \vec{x}^{(0)}}{\|A^k \vec{x}^{(0)}\|_\infty}$$

Suppose  $A$  is diagonalizable. That's, we can assume  $\vec{x}_1, \vec{x}_2, \dots, \vec{x}_n$  form a basis for  $\mathbb{C}^n$ .

Take  $\vec{x}^{(0)} = a_1 \vec{x}_1 + a_2 \vec{x}_2 + \dots + a_n \vec{x}_n$  (assuming  $a_1 \neq 0$ )

$$\vec{x}^{(k)} = \frac{a_1 \lambda_1^k \left[ \vec{x}_1 + \sum_{j=2}^n \frac{a_j}{a_1} \left( \frac{\lambda_j}{\lambda_1} \right)^k \vec{x}_j \right]}{\| a_1 \lambda_1^k \left[ \vec{x}_1 + \sum_{j=2}^n \frac{a_j}{a_1} \left( \frac{\lambda_j}{\lambda_1} \right)^k \vec{x}_j \right] \|_\infty} \approx \underbrace{\frac{a_1}{|a_1|} \frac{\vec{x}_1}{\|\vec{x}_1\|_\infty} \frac{\lambda_1^k}{|\lambda_1|^k}}_{\vec{v}}$$

when  $k$  is big

$$(\text{Because } 1 > \left| \frac{\lambda_2}{\lambda_1} \right| \geq \left| \frac{\lambda_3}{\lambda_1} \right| \geq \dots \geq \left| \frac{\lambda_n}{\lambda_1} \right|)$$

Rate of convergence of Power's method depends on  $\left| \frac{\lambda_2}{\lambda_1} \right|$  ← second largest eigenvalue  
 ← largest eigenvalue

In fact,

$$\| A \vec{x}^{(k)} \|_\infty \rightarrow \left\| A \left( \frac{a_1}{|a_1|} \frac{\vec{x}_1}{\|\vec{x}_1\|_\infty} \right) \right\|_\infty = \left\| \frac{a_1}{|a_1|} \lambda_1 \frac{\vec{x}_1}{\|\vec{x}_1\|_\infty} \right\|_\infty = |\lambda_1|$$

∴

$$|\lambda_1| + O\left(\left(\frac{|\lambda_2|}{|\lambda_1|}\right)^k\right)$$

Convergence rate:

$$|\lambda_1| > |\lambda_2| > \dots > |\lambda_n|$$

1. Power method:

Converges if  $\eta \stackrel{\text{def}}{=} \left| \frac{\lambda_2}{\lambda_1} \right| < 1$  and  $\langle \vec{v}_1, \vec{x}^{(0)} \rangle \neq 0$  ( $\vec{v}_1 =$  eigenvector of  $\lambda_1$ )

Also,  $\rho_k = \|A \vec{x}^{(k)}\|_\infty = |\lambda_1| + \mathcal{O}(\eta^k)$  (Slow convergence if  $\eta \approx 1$ !)

2. Inverse Power method:

Converges if  $\left| \frac{1/\lambda_{n-1}}{1/\lambda_n} \right| = \left| \frac{\lambda_n}{\lambda_{n-1}} \right| < 1$  and  $\langle \vec{v}_n, \vec{x}^{(0)} \rangle \neq 0$  ( $\vec{v}_n =$  eigenvector of  $\lambda_n$ )

Also,  $\rho_k = \|A \vec{x}^{(k)}\|_\infty = |\lambda_n| + \mathcal{O}(\eta^k)$  (Slow convergence if  $\eta \approx 1$ !)

3. Inverse Power method with shift, let  $\lambda_j$  be closest to  $\mu$ .

Converges if:  $\eta = \max_{m \neq j} \left| \frac{\lambda_j - \mu}{\lambda_m - \mu} \right| < 1$  and  $\langle \vec{v}_j, \vec{x}^{(0)} \rangle \neq 0$  ( $\vec{v}_j =$  eigenvector of  $\lambda_j$ .)

$\rho_k = \|A \vec{x}^{(k)}\|_\infty = |\lambda_j| + \mathcal{O}(\eta^k)$  (Slow convergence if  $\eta \approx 1$ !)

How to speed up convergence? Let  $A \in M_{n \times n}(\mathbb{R})$

Idea: Use Inverse Power method with shift, update  $\mu$  in each iteration (such that  $\mu$  is closer to a real eigenvalue in each iteration)

Then:  $\eta \stackrel{\text{def}}{=} \max_{m \neq j} \left| \frac{\lambda_j - \mu}{\lambda_m - \mu} \right|$  becomes smaller and smaller  $\Rightarrow$  Converges faster and faster!

Definition: (Rayleigh quotient) Let  $\vec{v} \neq \vec{0} \in \mathbb{R}^n$ ,  $A \in M_{n \times n}$ . Then, the Rayleigh quotient is defined as:  $R(\vec{v}, A) = \frac{\vec{v}^* A \vec{v}}{\vec{v}^* \vec{v}}$ .

Remark: Let  $A$  be symmetric positive definite. Then: all eigenvalues:

$\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$  are real.

Then:  $\lambda_n \leq R(\vec{v}, A) \leq \lambda_1$  and

$R(\vec{v}, A) = \lambda_1$  when  $\vec{v} = \vec{v}_1$  = eigenvector of  $\lambda_1$ .

$R(\vec{v}, A) = \lambda_n$  when  $\vec{v} = \vec{v}_n$  = eigenvector of  $\lambda_n$

$R(\vec{v}, A)$  can be regarded as the approximation of eigenvalue  $\lambda_j$ , given that  $\vec{v}$  is closed to  $\vec{v}_j$ .



## Rayleigh Quotient Iteration

Let  $A \in M_{n \times n}(\mathbb{R})$

Initiate  $\vec{x}^{(0)}$  such that  $\vec{x}^{(0)T} \vec{x}^{(0)} = 1$

Initiate  $\mu_0$  = initial guess of desired eigenvalue.

Solve:  $(A - \mu_0 I) \vec{z}_1 = \vec{x}^{(0)}$

Let  $\vec{x}^{(1)} = \frac{\vec{z}_1}{\|\vec{z}_1\|_2}$  ( $\|\vec{x}\|_2 \stackrel{\text{def}}{=} \sqrt{\vec{x}^T \vec{x}}$ )

Let  $\mu_1 = R(\vec{x}^{(1)}, A) = \vec{x}^{(1)T} A \vec{x}^{(1)}$  (Improve  $\mu_0$  such that it is closer to an actual eigenvalue)

Keep iteration going!

Algorithm: (Rayleigh Quotient Iteration)

Input:  $\vec{x}^{(0)}$  s.t.  $\|\vec{x}^{(0)}\|_2 = 1$  and  $\mu_0$

Output:  $\mu_k$  = eigenvalue

For  $k = 0, 1, 2, \dots$

Step 1: Solve  $(A - \mu_k I) \vec{z}_{k+1} = \vec{x}^{(k)}$

Step 2: Let  $\vec{x}^{(k+1)} = \frac{\vec{z}_{k+1}}{\|\vec{z}_{k+1}\|_2}$

Step 3: Compute  $\mu_{k+1} = R(\vec{x}^{(k+1)}, A)$ .

Example: Let  $A = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 1 \\ 3 & 2 & 1 \end{pmatrix}$ .

Eigenvalues:  $\lambda_1 = 3 + \sqrt{5}$ ,  $\lambda_2 = 3 - \sqrt{5}$ ,  $\lambda_3 = -2$ .

Want to find  $3 + \sqrt{5}$ .

Let  $\vec{x}^{(0)} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$ ,  $\mu_0 = 2.00$

Then:  $\vec{x}^{(1)} \approx \begin{pmatrix} -0.57927 \\ -0.57348 \\ -0.57927 \end{pmatrix}$  with  $\mu_1 = 5.3355$

Converges very fast!  $\mu_3 = 5.281 \approx 3 + \sqrt{5}$ !

Remark:

- RQI works for SPD  $A$
- May or may not work for other  $A$ .

What if  $\lambda_1$  has multiplicity  $> 1$ ?

Consider the case when  $A$  is diagonalizable.

Let  $\{\vec{x}_1, \vec{x}_2, \dots, \vec{x}_n\}$  be the basis of eigenvectors with eigenvalues equal to  $\lambda_1, \lambda_2, \dots, \lambda_n$ .

Assume that:  $\lambda_1 = \lambda_2 = \dots = \lambda_i > |\lambda_{i+1}| \geq \dots \geq |\lambda_n|$ .

Let  $\vec{x}^{(0)} = a_1 \vec{x}_1 + a_2 \vec{x}_2 + \dots + a_n \vec{x}_n$  (with  $a_1 \neq 0$ )

Easy to check:  $\vec{x}^{(k)} = \frac{\lambda_1^k (a_1 \vec{x}_1 + \dots + a_i \vec{x}_i + (\frac{\lambda_{i+1}}{\lambda_1})^k \vec{x}_{i+1} + \dots + (\frac{\lambda_n}{\lambda_1})^k \vec{x}_n)}{\left\| \lambda_1^k (a_1 \vec{x}_1 + \dots + a_i \vec{x}_i + (\frac{\lambda_{i+1}}{\lambda_1})^k \vec{x}_{i+1} + \dots + (\frac{\lambda_n}{\lambda_1})^k \vec{x}_n) \right\|_\infty}$

$$\rightarrow \frac{a_1 \vec{x}_1 + \dots + a_i \vec{x}_i}{\|a_1 \vec{x}_1 + \dots + a_i \vec{x}_i\|_\infty} \text{ as } k \rightarrow \infty.$$

↑  
Eigenvector of  $\lambda_1$ .

Also,  $\|A \vec{x}^{(k)}\|_\infty \rightarrow \left\| \frac{A(a_1 \vec{x}_1 + \dots + a_i \vec{x}_i)}{a_1 \vec{x}_1 + \dots + a_i \vec{x}_i} \right\|_\infty = |\lambda_1| \text{ as } k \rightarrow \infty$

Remark: The condition on multiplicity ( $= 1$ ) can be relaxed.

## Method 2: QR method

Preliminary: QR factorization

Definition:  $Q \in M_{n \times n}(\mathbb{R})$  is orthogonal if  $Q^T Q = I_n$

Remark: -  $Q^{-1} = Q^T$

- Columns of  $Q$  forms orthonormal set.

## QR factorization

Let  $A$  be a non-singular  $n \times n$  matrix. There exists an orthogonal matrix  $Q$  and upper triangular matrix  $R$  such that:

$$A = QR$$



Preview: (Move next time)

QR method to find eigenvalues

Algorithm: (QR algorithm)

Input :  $A \in M_{n \times n}(\mathbb{R})$

Step 1: Let  $A^{(0)} = A$ . Compute QR factorization of  $A^{(0)} = Q_0 R_0$ .

Let  $A^{(1)} = R_0 Q_0$ .

Step 2: Assume  $A^{(1)}, \dots, A^{(k)}$  are computed. Let  $A^{(k)} = Q_k R_k$ .

be the QR factorization of  $A^{(k)}$ . Let  $A^{(k+1)} = R_k Q_k$ .