CSCI3160: Special Exercise Set 6

Prepared by Yufei Tao

Problem 1. Define function f(x) — where $x \ge 0$ is an integer — as follows:

- f(0) = 0
- f(1) = 1
- f(x) = f(x-1) + f(x-2).

Give an algorithm to calculate f(n) in O(n) time (you can assume that f(x) fits in a word for all $x \leq n$).

Problem 2. Let A be an array of n integers. Consider the following recursive function which is defined for any i, j satisfying $1 \le i \le j \le n$:

$$f(i,j) = \begin{cases} 0 & \text{if } i = j \\ A[i] \cdot A[j] + \min_{k=i+1}^{j-1} f(i,k) + f(k,j) & \text{if } i \neq j \end{cases}$$

Design an algorithm to calculate f(1, n) in $O(n^3)$ time.

Problem 3. In Lecture Notes 8, we defined function f(i, j) based on strings x = ABC and y = BDCA. Calculate f(i, j) for all possible i and j.

Problem 4. In the rod-cutting problem, suppose that n = 5 and the price array P is (2, 6, 7, 9, 10). What is the maximum revenue achievable?

Problem 5 (Textbook Problem 15.1-3). Consider a modification of the rod-cutting problem in which, in addition to a price P[i] for each length $i \in [1, n]$, each cut incurs a fixed cost of c. The revenue associated with a solution is now the sum of the prices of the segments minus the total cost of making the cuts. Give a dynamic-programming algorithm to solve this modified problem in $O(n^2)$ time.