

TreeNet: Deep Point Cloud Embedding for Routing Tree Construction

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Biography

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Routing Tree Construction

Routing Tree Construction: Given a input net $V = \{v_0, V_s\}$, v_0 is the source (red node) and *Vs* is the set of sinks (black node), construct a tree optimizing both wire length and path length.

Examples of routing tree construction. Left: spanning tree; right: Steiner tree.

Wire length (WL) and path length (PL)

Wire length (WL) metric: lightness

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WL ratio with that of minimum spanning tree (MST).

$$
\blacktriangleright \text{ } \textit{lightness} = \frac{w(T)}{w(MST(G))}, w(\cdot) \text{ is the total weight.}
$$

Path length (PL) metric: shallowness by SALT^{*} or normalized path length by PD-II⁺

► Shallowness = max{
$$
\frac{d_T(v_0, v)}{d_G(v_0, v)} | v \in V_s
$$
}, *G* is the connected weighted routing graph.
\n► Normalized path length =
$$
\frac{\sum_{v \in V} d_T(v_0, v)}{\sum_{v \in V} d_G(v_0, v)}
$$
.

∗Gengjie Chen and Evangeline FY Young (2019). "SALT: provably good routing topology by a novel steiner shallow-light tree algorithm". In: *IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems*.

 Ω †Charles J Alpert et al. (2018). "Prim-Dijkstra Revisited: Achieving Superior Timi[ng](#page-2-0)-[dri](#page-4-0)[ve](#page-2-0)[n](#page-3-0) [R](#page-4-0)[ou](#page-1-0)[ti](#page-2-0)[n](#page-6-0)[g](#page-7-0) [T](#page-1-0)[re](#page-2-0)[e](#page-6-0)[s](#page-7-0)[".](#page-0-0) In: *Proc. ISPD*, pp. 10–17.

Best algorithm?

In Neither PD-II nor SALT, two most prominent ones, always dominates the other one in terms of both WL and PL for all nets.

Best parameter?

- Both PD-II and SALT use a parameter to help balance WL and PL.
- Given one WL constraint, what is the best parameter to obtain the best PL?

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Point cloud and its embedding

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Cloud embeddings for tree construction, where point clouds are transformed into unified 2-D Euclidean space. $A \cup B \cup A \cup B \cup A \cup B \cup A \cup B \cup B$

Problem formulation

Given a set of 2-D pins and two routing tree construction algorithms, SALT‡ and PD-II, our objective is to **obtain the embedding** of the given point cloud by TreeNet such that

- 1. the embedding can be used to **select the best algorithm** for the given point cloud;
- 2. the embedding can be used to **estimate the best parameter** ϵ of **SALT** for the given point cloud;
- 3. the embedding can be used to **estimate the best parameter** α of PD-II for the given point cloud.

[‡]Gengjie Chen and Evangeline FY Young (2019). "SALT: provably good routing topology by a novel steiner shallow-light tree algorithm". In: *IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems*.

Charles J Alpert et al. (2018). "Prim-Dijkstra Revisited: Achieving Superior Timing-driven Routing Trees". In: *Proc. ISPD*, pp. 10–17. $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$

Property 1: Down-sampling

Property

Let $d: V \to V'$ be a function for down-sampling, where V' is a proper subset of V . $f(V) \neq f(d(V))$ *holds if there exists* $v \in V - d(V)$ *so that v is not the steiner point in* $f(d(\mathbf{V}))$.

Examples of the down-sampling: (a) The general point cloud without the down-sampling; (b) The general point cloud with the down-sampling; (c) The constructed tree without the down-sampling; (d) The constructed tree with the down-sampling.

Property 2 & 3: Permutation

Property

Let \pmb{V}_s^p be the permutation of the sink set \pmb{V}_s . $f(\{v_0, \pmb{V}_s^p\}) =$ $f(\{v_0, \pmb{V}_s\})$ holds for any $V = \{v_0, V_s\}.$

Property

Let V^p be the permutation of the input net V . $f(V^p) \neq f(V)$ holds if the source in V^p is *different from the source in V.*

Examples of the routing trees with the same node coordinates but di[ffe](#page-7-0)r[en](#page-9-0)[t](#page-7-0) [so](#page-8-0)[u](#page-9-0)[r](#page-6-0)[c](#page-7-0)[e](#page-10-0) QQ $9/21$ (highlighted by red).

Property 4: Inequality of the same *V^s*

Property

For any sink set V_s *with* $|V_s|>1$ *, there exists two different pins,* v_0 *and* v_0^\prime *in the 2-D plane* $\texttt{so that} f(\{ {\sf v}_0, {\sf V}_s\}) \neq f(\{ {\sf v}_0', {\sf V}_s\}).$ Moreover, the inequality holds when we only consider the *topology.*

Examples of the node with the same coordinates and local neighbors but different parent-child relationships. Here root is highlighted in red.

Property 5: Graph construction methods

Property

Let Gball, Gknn and Gbbox be the graph constructed from V by ball query, k nearest neighbor and bounding box respectively. The minimum spanning tree, T may not be the subgraph of G_{ball} *or* G_{nn} *, but always the subgraph of* G_{bbox} *.*

Comparison among ball query (a) k-nn (b) and k-bbox (c) grouping methods $(k = 2$ in this example). The orange regions represent the query ball in (a) and bo[un](#page-9-0)d[in](#page-11-0)[g](#page-9-0) [b](#page-10-0)[o](#page-11-0)[xe](#page-6-0)[s](#page-7-0) [in](#page-11-0)[\(](#page-7-0)[c\)](#page-10-0)[.](#page-11-0) Ω $_{11/21}$ The centroid is highlighted by black and the root is by red.

TreeConv

- \triangleright Sampling selects a set of centroids from the original point cloud
	- Omited considering Property [1.](#page-7-1)
	- Each node is selected as the centroid.
- \triangleright Grouping selects a set of neighbors for each centroid.
	- Selecting *k* nearest *bbox-neighbors* of *uⁱ* as the neighbors.
	- Grouping **returns a list of neighbors** $\boldsymbol{E}_i \in \mathbb{R}^k$ **for each centroid** u_i **.**
- \triangleright Encoding is to encode the new centroid feature using the original one and the local feature aggregated from the neighbors of the centroid.
	- $v'_{ic} = \max_{j \in E_i} \sigma(\boldsymbol{\theta}_c \cdot \text{CONCAT}(\boldsymbol{v}_i, \boldsymbol{v}_i \boldsymbol{v}_j, \boldsymbol{v}_i \boldsymbol{v}_r))$
	- followed by a Squeeze-and-Excitation (SE) block¶

[¶]Jie Hu, Li Shen, and Gang Sun (2018). "Squeeze-and-excitation networks". In: *Proceedings of the IEEE conference on computer vision and pattern recognition*, pp. 7132–7141. $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$ $(1,1)$

Illustration of TreeConv. Brighter blocks indicate Grouping and darker blocks indicate Encoding. メロトメ 伊 トメ ミトメ ミト

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TreeConv vs. existing methods.

‖Charles R Qi et al. (2017). "Pointnet: Deep learning on point sets for 3d classification and segmentation". In: *Proc. CVPR*, pp. 652–660.

∗∗Charles Ruizhongtai Qi et al. (2017). "PointNet++: Deep hierarchical feature learning on point sets in a metric space". In: *Advances in Neural Information Processing Systems*, pp. 5099–5108.

††Yangyan Li et al. (2018). "PointCNN: Convolution on x-transformed points". In: *Advances in Neural Information Processing Systems*, pp. 820–830.

‡‡Yue Wang et al. (2019). "Dynamic graph CNN for learning on point clouds". In: *ACM Transactions on Graphics* 38.5, pp. 1–12. $(1 + 4\sqrt{3}) + 4\sqrt{3} + 4\sqrt{3} + 4\sqrt{3} + \sqrt{3} + 4\sqrt{3} + 4\$

Illustration of TreeNet Architecture for the cloud embedding.

$$
\blacktriangleright \text{ Normalization: } \tilde{v}_i = \frac{v_i - v_r}{d_{max}}.
$$

Algorithm selction & parameter predition

Algorithm selction

 $y =$ softmax $(W_3\sigma(W_2\sigma(W_1H_c+b_1)+b_2)),$

Parameter predition

- ▶ 20 valid parameter $\epsilon_i, i \in \{1, ..., 20\}$ candidates for SALT
- Following similar structure with algorithm selection to obtain the output $y \in \mathbb{R}^{20}$.
- Given the output *y*, the predicted parameter ϵ is calculated by an element-wise summation and can be formulated as $\epsilon = \sum_{i=1}^{20} \epsilon_i \cdot y_i.$
- \blacktriangleright The predicted parameter guides the routing tree construction by a simple heuristic rule

The workflow of our framework. Dotted arrows represent that TreeNet generates cloud embeddings and use them to select the algorithm or to predict parameters. The yellow blocks are executed in our framework while the purple blocks are executed by the selected algorithms.

Comparison to existing methods

Comparison to SALT & PD-II (shallowness & normalized PL)

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Runtime

Runtime comparison with SALT and SALT[∗] .

Runtime breakdown of our framework.

- **We formalize special properties** of the point cloud for the routing tree construction;
- I We design **TreeNet**, a novel deep net architecture to obtain the cloud embedding for the tree construction;
- I We propose an **adaptive flow** for the routing tree construction, which uses the cloud embedding to **select the best approach** and **predict the best parameter**;
- Experiments on widely used benchmarks demonstrate the effectiveness of our embedding representation, compared with all other deep learning models;
- Experiments also show that our methods outperform other state-of-the-art routing tree construction methods in terms of both quality and runtime.