

## CSCI2100: Quiz 2

Name:

Student ID:

**Note:** A multiple-choice question has only one correct answer unless otherwise stated.

**Problem 1 (10%).** After applying the following operations to an empty stack:

push(35), push(36), push(43), push(8), pop, pop, push(51), pop

what is the content of the stack? Answer: [       ]

- A. 35, 36
- B. 36, 43
- C. 35, 8
- D. 35, 51

**Answer:** A

**Problem 2 (10%).** After applying the following operations to an empty queue:

enqueue(35), enqueue(36), enqueue(43), enqueue(8), dequeue, dequeue, enqueue(51), dequeue

what is the content of the queue? Answer: [       ]

- A. 8, 51
- B. 36, 43
- C. 35, 8
- D. 35, 51

**Answer:** A

**Problem 3 (10%).** Identify the operations below that can be performed in  $O(1)$  cost (including  $O(1)$  expected cost). There is more than one correct choice; no marks are given unless you can identify all of them. Answer: [       ]

- A. Push an element into a stack.
- B. Insert an element into a linked list.
- C. Dequeue an element from a queue.
- D. Determine whether the value 10 is in a hash table.

**Answer:** ABCD

**Problem 4 (10%).** Which of the following are true? There is more than one correct choice; no marks are given unless you can identify all of them. Answer: [       ]

- A. Consider a data structure that supports a certain operation in  $O(1)$  amortized time. Then, any sequence of  $n$  such operations requires  $O(n)$  worst case time, regardless of the value of  $n$ .
- B. Consider a data structure that supports a certain operation in  $O(1)$  amortized time. But still, it is possible for the structure to take  $O(n)$  time to process *one* operation, where  $n$  is the number of operations that have already been processed.
- C. There is a hash function that can guarantee  $O(1)$  expected dictionary search on all input sets.
- D. In a tree, the number of internal nodes cannot exceed that of leaf nodes.

**Answer:** AB

**Problem 5 (10%).** Consider  $S = \{1, 3, 4, 5, 7, 8, 9, 10, 11, 12, 13, 15\}$ . We build a hash table on  $S$  with hash function  $h(k) = 1 + k \bmod 4$ . If we look up an element  $q$  in the hash table, which of the following value of  $q$  has the lowest look up cost? Answer: [            ]  
A. 24    B. 23    C. 22    D. 21

**Answer:** C

**Problem 6 (20%).** Consider the hash function  $h(k) = 1 + k \bmod 4$ . Give a set  $S$  of 10 integers to meet both conditions below:

- If we build a hash table on  $S$  using  $h(k)$ , then all the integers of  $S$  fall in the same bucket (recall that a *bucket* contains all the elements of  $S$  having the same hash value).
- The aforementioned bucket is the one we probe in order to look for integer 35.

**Answer:**  $S = \{3, 7, 11, 15, 19, 23, 27, 31, 35, 39\}$ .

**Problem 7 (30%).** Let  $S$  be a set of  $n \geq 2$  distinct integers where  $n$  is a power of 2. The set  $S$  is given in an array that has not been sorted. Give an algorithm to find the  $\log_2 n$  largest integers of  $S$  in  $O(n \log \log n)$  time. For example, if  $S = \{30, 50, 10, 90, 80, 20, 70, 60\}$ , then you should output 70, 80, 90.

Note: by using  $k$ -selection, we can solve the problem in  $O(n)$  expected time. However, here our  $O(n \log \log n)$  time bound needs to hold deterministically. If you want to use a deterministic  $k$ -selection algorithm, you must describe the algorithm in full (because it has not been covered in this course). This problem admits an elegant solution that does not require  $k$ -selection.

**Answer 1:** Initialize an empty priority queue  $H$  (min-heap). Process each element  $e \in S$  in turn as follows. First, insert  $e$  to  $H$ . Then, check if  $H$  has more than  $\log_2 n$  elements; if so, perform a delete-min (after which  $H$  will have exactly  $\log_2 n$  elements). After all the elements of  $S$  have been processed, report the elements of  $H$  in ascending order with  $\log_2 n$  delete-mins.

**Answer 2:** Construct a max-heap from  $S$  in  $O(n)$  time. Then, perform  $\log_2 n$  delete-max operations and returned the elements found. The total cost is  $O(n + \log^2 n) = O(n)$ .