CSCI2100: Regular Exercise Set 12

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Problem 1. Let $G = (V, E)$ be a directed graph. Suppose that we perform BFS starting from a source vertex s, and obtain a BFS-tree T. For any vertex $v \in V$, denote by $l(v)$ the level of v in the BFS-tree. Prove that BFS en-queues the vertices v of V in non-descending order of $l(v)$.

Solution. Take any vertices u, v such that $l(u) > l(v)$. Let $v_1, v_2, ..., v_{l(v)}$ be the vertices on the path from the root to v in T; note that $v_1 = s$ and $v_{l(v)} = v$. Let $u_1, u_2, ..., u_{l(v)}$ be the last $l(v)$ vertices on the path from the root to u in T; note that $u_1 \neq s$ and $u_{l(v)} = u$. It thus follows that v_1 is en-queued before u_1. Remember that BFS en-queues v_2 when de-queuing v_1, and similarly, en-queues u_2 when de-queuing u_1. By the FIFO property of the queue, we know that v_2 is en-queued before u_2. By the same reasoning, v_3 is en-queued before u_3, v_4 before u_4, etc. This means that v is before u.

Problem 2. Let $G = (V, E)$ be a directed graph. Suppose that we perform BFS starting from a source vertex s, and obtain a BFS-tree T. For any vertex $v \in V$, prove that the path from s to v in T is a shortest path from s to v in G.

Solution. We will instead prove the following claim: *all the vertices with shortest path distance l from s are at level l* (recall that the root is at level 0). This will establish the conclusion in Problem 3 because the path from s to a level-l node v in T has length l.

We will prove the claim by induction on l. The base case where $l = 0$ is obviously true.

Assuming that the claim holds for all $l \leq k - 1$ ($k \geq 1$), next we prove that the claim is also true for $l = k$. Let v be a vertex with shortest path distance k from s. Consider all the shortest paths from s to v. From every such shortest path, take the vertex immediately before v (i.e., the predecessor of v in that path), and put that vertex into a set. Let S be the set of all such "predecessors of v" collected. Let u be the vertex in S that is the earliest one entering the queue. We know that the shortest path distance from s to u is $k - 1$. It thus follows from the inductive assumption that u is at level $k - 1$ of T.

Consider the moment when u is removed from the queue in BFS. We will argue that the color of v must be white. Hence, BFS makes v a child of u, thus making v at level k.

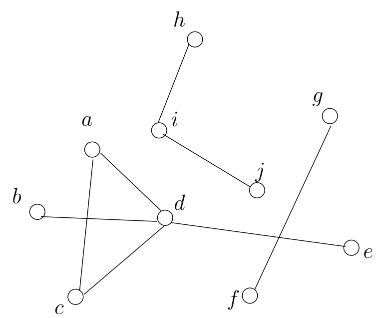
Suppose for contradiction that the color of v is gray or black. This means that v has been put into the queue when another vertex u' was de-queued earlier. From the conclusion of Problem 2 and the definition of u, we know that $l(u') < l(u)$ (otherwise, $l(u') = l(u) = k - 1$; by the inductive assumption, this means that the shortest path distance from s to u' is $k - 1$, which further implies $u' \in S$, contradicting the definition of u). From the inductive assumption, this means that the shortest path distance of u' from s that is less than $k - 1$, implying that the shortest path distance from s to v is less than k, thus giving a contradiction.

Problem 3. Let $G = (V, E)$ be an undirected graph. We will denote an edge between vertices u, v as $\{u, v\}$. Next, we define the single source shortest path (SSSP) problem on G. Define a *path* from s to t as a sequence of edges $\{v_1, v_2\}, \{v_2, v_3\}, ..., \{v_t, v_{t+1}\}$, where $t \geq 1$, $v_1 = s$, and $v_{t+1} = t$. The *length* of the path equals t. Then, the SSSP problem gives a source vertex s, and asks to find shortest paths from s to all the other vertices in G. Adapt BFS to solve this problem in $O(|V| + |E|)$ time. Once again, you need to produce a BFS tree where, for each vertex $v \in V$, the path from the root to v gives a shortest path from s to v.

Solution. Same as BFS, except that when a vertex v is de-queued, we inspect all its neighbors (as opposed to its out-neighbors as in the directed version).

Problem 4 (Connected Components). Let $G = (V, E)$ be an undirected graph. A *connected component* (CC) of G includes a set $S \subseteq V$ of vertices such that

- For any vertices $u, v \in S$, there is a path from u to v, and a path from v to u.
- (Maximality) It is not possible to add any vertex into S while still ensuring the previous property.



For example, in the above graph, $\{a, b, c, d, e\}$ is a CC, but $\{a, b, c, d\}$ is not, and neither is $\{g, f, e\}$.

Prove: Let S_1, S_2 be two CCs. Then, they must be disjoint, i.e., $S_1 \cap S_2 = \emptyset$.

Solution. Suppose that a vertex v is in $S_1 \cap S_2$. Then, for any vertex $u_1 \in S_1$ and $u_2 \in S_2$, we know:

- There is a path from u_1 to u_2 by way of v.
- There is a path from u_2 to u_1 by way of v.

This violates the fact that S_1 and S_2 must be maximal.

Problem 5. Let $G = (V, E)$ be an undirected graph. Describe an algorithm to divide V into a set of CCs. For example, in the example of Problem 5, your algorithm should return 3 CCs: $\{a, b, c, d, e, \}, \{g, f\}$, and $\{h, i, j\}$.

Solution. Run BFS starting from an arbitrary vertex in V. All the vertices in the BFS-tree constitute the first CC. Then, start another BFS from an arbitrary vertex that is still white. All the vertices in this BFS-tree constitute another CC. Repeat this until V has no more white vertices.