

Secure ISAC Downlink Beamforming: A SDR Approach with Tightness Guaranteed

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Slides available at www.cse.cuhk.edu.hk/~wykeung/slides/secure-isac.pdf

Motivation

- the fight for bandwidth resources has always been competitive; spectrum congestion remains critical in 5G & beyond as comm. sys. eye at the radar band
- **radar-communication coexistence**: use interference management techniques to allow both radar and comm. systems to service in the same band
- **integrated sensing and communication (ISAC)**: to leverage the assigned resources by designing an unified signal for both radar and comm. purposes
 - advantages: unified hardware means cheap & energy-efficient implementation; allows efficient usage of bandwidth
 - challenges: difficult to employ network layer secrecy measures in dense networks; requires dedicated signal designs that serve both sensing and comm. purposes
- **our goal**: provide a beamformer design framework that serves radar-comm. with physical layer security using the least amount of radiation power

Problem Scenario

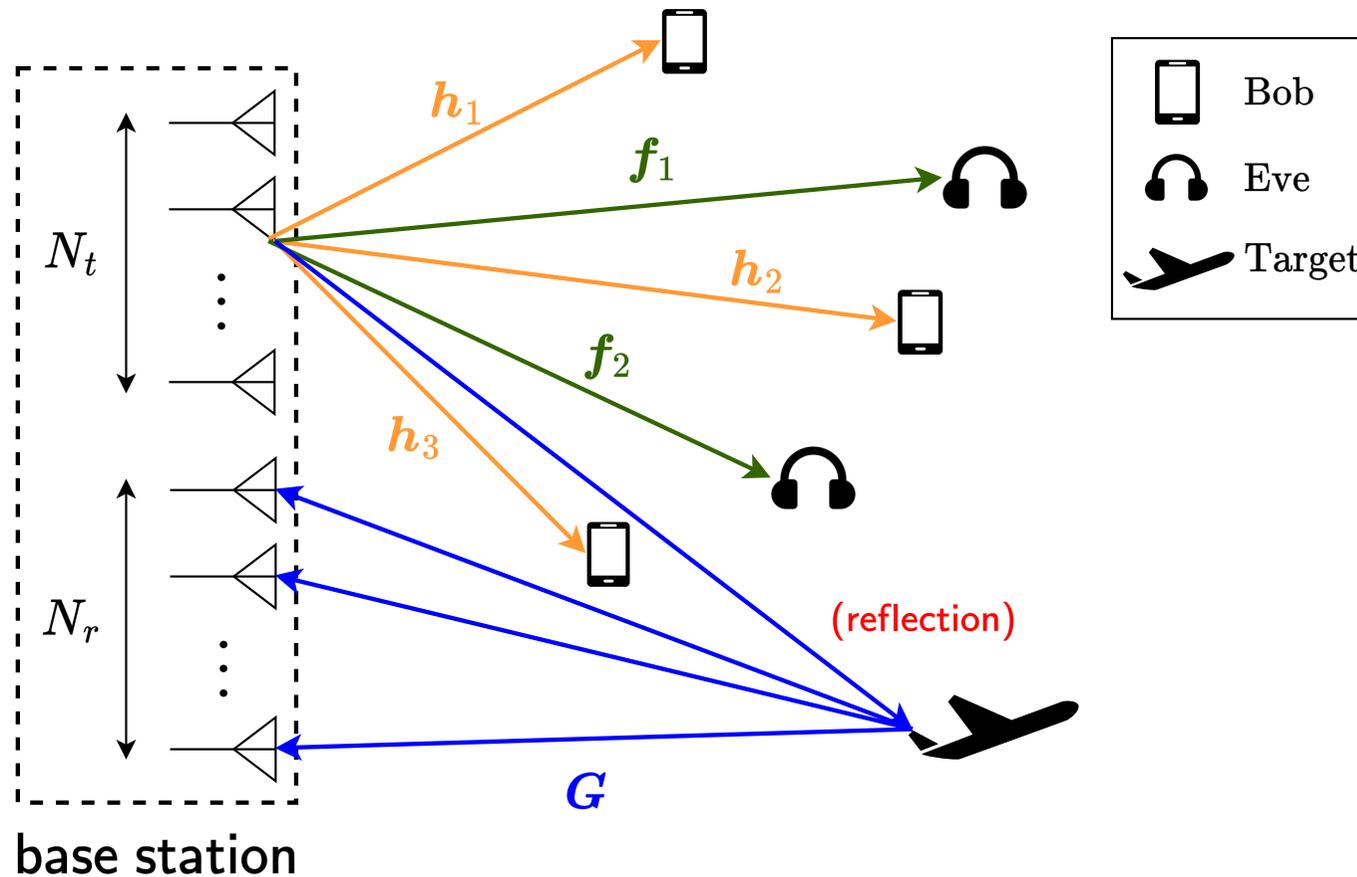


Figure: Illustration of an ISAC beamforming scenario under the presence of eavesdropper.

- the base station (BS) has N_t transmit antenna with N_r sensing antenna, servicing in a cell over a block-fading channel
- BS communicates with K legitimate users (Bobs) with J eavesdroppers (Eves) overhearing the channel; BS also aims to estimate an unknown response matrix G by radar

Signal Model

- $\mathbf{x}_t \in \mathbb{C}^{N_t}$ is the signal vector; under some standard assumptions, it can be written as

$$\mathbf{x}_t = \sum_{k=1}^K \mathbf{w}_k s_{k,t} + \mathbf{b}_t$$

where

- $s_{k,t} \in \mathcal{S}$ is the information carrying symbol, where \mathcal{S} is a finite alphabet
 - $\mathbf{w}_k \in \mathbb{C}^{N_t}$ is the beamformer for the k -th Bob
 - \mathbf{b}_t is a non-message carrying signal to serve the radar and jamming purposes
- conventionally $\mathbf{b}_t \sim \mathcal{CN}(\mathbf{0}, \mathbf{\Sigma})$ is called as the artificial noise (AN) in physical layer security; we denote its spatial covariance matrix as $\mathbf{\Sigma} = \mathbf{W}_{K+1}$ for convenience
 - received signal of the k -th Bob, the i -th Eve, and the ℓ -th sensing antenna are:

$$y_{k,t} = \mathbf{h}_k^H \mathbf{x}_t + v_{k,t}, \quad k = 1, \dots, K, \quad (\text{Bob's rx signal})$$

$$r_{i,t} = \mathbf{f}_i^H \mathbf{x}_t + v_{i,t}, \quad i = 1, \dots, J, \quad (\text{Eves' rx signal})$$

$$z_{\ell,t} = \mathbf{g}_{\ell}^H \mathbf{x}_t + v_{\ell,t}, \quad \ell = 1, \dots, N_r, \quad (\text{Radar rx signal})$$

where v 's are the respective background noise with zero-mean and σ^2 -variance

- assumption: the BS knows \mathbf{h}_k 's and \mathbf{f}_i 's, but not \mathbf{g}_{ℓ} 's

Quality of Service: Our Choices

$$\begin{array}{ll} \text{minimize} & \text{signal power} \\ \{\mathbf{W}_k\}_{k=1}^{K+1} & \end{array} \quad (2)$$

$$\text{subject to} \quad \text{comm. metric} \geq \gamma \quad (C)$$

$$\text{security metric} \leq \rho \quad (J)$$

$$\text{sensing error metric} \leq \beta \quad (S)$$

$$\mathbf{W}_{K+1} \succeq \mathbf{0}, \quad \mathbf{W}_k = \mathbf{w}_k \mathbf{w}_k^H \quad \forall k = 1, \dots, K$$

- signal power: $\mathbb{E}(\mathbf{x}_t^H \mathbf{x}_t) = \sum_{k=1}^K \mathbf{w}_k^H \mathbf{w}_k \mathbb{E}(|s_{k,t}|^2) + \mathbb{E}(\mathbf{b}_t^H \mathbf{b}_t) = \sum_{k=1}^{K+1} \text{tr}(\mathbf{W}_k)$
- comm./ security metric: signal-to-interference-plus-noise ratio, e.g.

$$\text{SINR}_k(\{\mathbf{W}_i\}_{i=1}^{K+1}) = \frac{|\mathbf{h}_k^H \mathbf{w}_k|^2}{\sum_{j=1, j \neq k}^K |\mathbf{h}_k^H \mathbf{w}_j|^2 + \mathbf{h}_k^H \mathbf{W}_{K+1} \mathbf{h}_k + \sigma^2}$$

is defined for Bobs; replace \mathbf{h}_k 's with \mathbf{f}_i 's for that of Eves

- radar sensing metric: Cramér-Rao bound of the target response matrix estimate [LLL⁺22]:

$$\text{CRB}(\{\mathbf{W}_k\}_{k=1}^{K+1}) \propto \text{tr}[(\sum_{k=1}^{K+1} \mathbf{W}_k)^{-1}]$$

- the QoS requirements are symbolized by γ, ρ and β

Quick Review on Semidefinite Relaxation (SDR)

Problem: complex-valued quadratically constrained quadratic programme (QCQP)

$$\begin{aligned} & \underset{\mathbf{w}_1, \dots, \mathbf{w}_K \in \mathbb{C}^N}{\text{minimize}} && \sum_{k=1}^K \mathbf{w}_k^H \mathbf{w}_k \\ & \text{subject to} && \sum_{k=1}^K \mathbf{w}_k^H \mathbf{C}_i \mathbf{w}_k \geq b_i, \quad \mathbf{C}_i \in \mathbb{H}^N, b_i \in \mathbb{R} \quad i = 1, \dots, I \end{aligned}$$

where the constraint matrices \mathbf{C}_i 's are Hermitian; The problem is non-convex.

SDR: leverage the equivalence of $\mathbf{W}_k = \mathbf{w}_k \mathbf{w}_k^H \iff \mathbf{W}_k \succeq \mathbf{0}, \text{rank}(\mathbf{W}_k) = 1$; dropping the rank one constraint directs us to a convex programme:

$$\begin{aligned} & \underset{\mathbf{W}_1, \dots, \mathbf{W}_K \in \mathbb{H}^N}{\text{minimize}} && \sum_{k=1}^K \text{tr}(\mathbf{W}_k) \\ & \text{subject to} && \sum_{k=1}^K \text{tr}(\mathbf{C}_i \mathbf{W}_k) \geq b_i, \quad \mathbf{C}_i \in \mathbb{H}^N, b_i \in \mathbb{R} \quad i = 1, \dots, I \\ & && \mathbf{W}_1, \dots, \mathbf{W}_K \succeq \mathbf{0}, \quad \text{rank}(\mathbf{W}_1) = \dots = \text{rank}(\mathbf{W}_K) = 1 \end{aligned}$$

This allows a comfortable convex approx. scheme for QCQP, and is suitable to many beamforming problems; but several issues remain:

- approx. means the solutions \mathbf{W}_k^* *may or may not* be of rank one
- the relaxation is said to be **tight** if $\text{rank}(\mathbf{W}_k^*) = 1$, or, equivalently $\mathbf{W}_k^* = \mathbf{w}_k \mathbf{w}_k^H$, for $k = 1, \dots, K$, i.e. \mathbf{w}_k serves as a feasible solution to the QCQP

Existing SDR-based Beamforming Formulations

- one of the most **classical formulations**:

$$\begin{aligned} & \text{minimize} && \text{signal power} \\ & \{\mathbf{W}_k = \mathbf{w}_k \mathbf{w}_k^H\}_{k=1}^K \\ & \text{subject to} && \text{Bobs' SINR} \geq \gamma \end{aligned}$$

has its SDR shown to be tight by the uplink-downlink duality [B001], or the Shapiro-Barvinok-Pataki SDP rank-reduction result [HP10]

- another design for **AN-assisted physical layer security** was studied in [LM16]:

$$\begin{aligned} & \text{minimize} && \text{signal power} \\ & \{\mathbf{W}_k = \mathbf{w}_k \mathbf{w}_k^H\}_{k=1}^K, \mathbf{W}_{K+1} \succeq \mathbf{0} \\ & \text{subject to} && \text{Bobs' SINR} \geq \gamma, \quad \text{Eves' SINR} \leq \rho \end{aligned}$$

wherein its SDR is tight

- an **ISAC variation** has been studied in [LLL⁺22], viz.:

$$\begin{aligned} & \text{minimize} && \text{CRB of the sensing error} \\ & \{\mathbf{W}_k = \mathbf{w}_k \mathbf{w}_k^H\}_{k=1}^K, \mathbf{W}_{K+1} \succeq \mathbf{0} \\ & \text{subject to} && \text{Bobs' SINR} \geq \gamma, \quad \text{signal power} \leq P \end{aligned}$$

in which a rank one solution can be obtained by a cheap post-processing technique

Proposed Formulation

$$\begin{aligned}
 & \underset{\{\mathbf{W}_k\}_{k=1}^{K+1}}{\text{minimize}} && \sum_{k=1}^K \text{tr}(\mathbf{W}_k) + \alpha \text{tr}(\mathbf{W}_{K+1}) \\
 & \text{subject to} && \frac{1}{\gamma} \text{tr}(\mathbf{W}_k \mathbf{H}_k) \geq \sum_{j=1, j \neq k}^{K+1} \text{tr}(\mathbf{W}_j \mathbf{H}_k) + \sigma^2 \quad (\text{Bob's SINR} \geq \gamma) && \text{(C)} \\
 & && \frac{1}{\rho} \text{tr}(\mathbf{W}_k \mathbf{F}_i) \leq \sum_{j=1, j \neq k}^{K+1} \text{tr}(\mathbf{W}_j \mathbf{F}_i) + \sigma^2 \quad (\text{Eve's SINR} \leq \rho) && \text{(J)} \\
 & && \text{tr}[(\sum_{k=1}^{K+1} \mathbf{W}_k)^{-1}] \leq T\beta/\sigma^2 N_r \quad (\text{CRB} \leq \beta) && \text{(S)} \\
 & && \mathbf{W}_{K+1} \succeq \mathbf{0}, \mathbf{W}_k = \mathbf{w}_k \mathbf{w}_k^H, \mathbf{W}_k \succeq \mathbf{0} \\
 & && \mathbf{H}_k = \mathbf{h}_k \mathbf{h}_k^H, \mathbf{F}_i = \mathbf{f}_i \mathbf{f}_i^H, \quad k = 1, \dots, K, i = 1, \dots, J.
 \end{aligned}$$

- major differences from (2):
 - * the rank one constraints on $\{\mathbf{W}_k\}_{k=1}^K$ are relaxed to PSD constraints
 - * a linear weighting $0 < \alpha < 1$ is introduced on the AN cov. matrix \mathbf{W}_{K+1}
- our study reveals that the above problem solves the non-cvx problem (2) almost exactly!

Main Result

$$\begin{aligned} & \underset{\{\mathbf{W}_k \in \mathbb{H}^{N_t}\}_{k=1}^K}{\text{minimize}} && \sum_{k=1}^K \text{tr}(\mathbf{W}_k) + \alpha \text{tr}(\mathbf{W}_{K+1}) && (\mathcal{R}) \\ & \text{subject to} && \text{constraints (C, J, S)} \\ & && \mathbf{W}_1, \dots, \mathbf{W}_{K+1} \succeq \mathbf{0} \end{aligned}$$

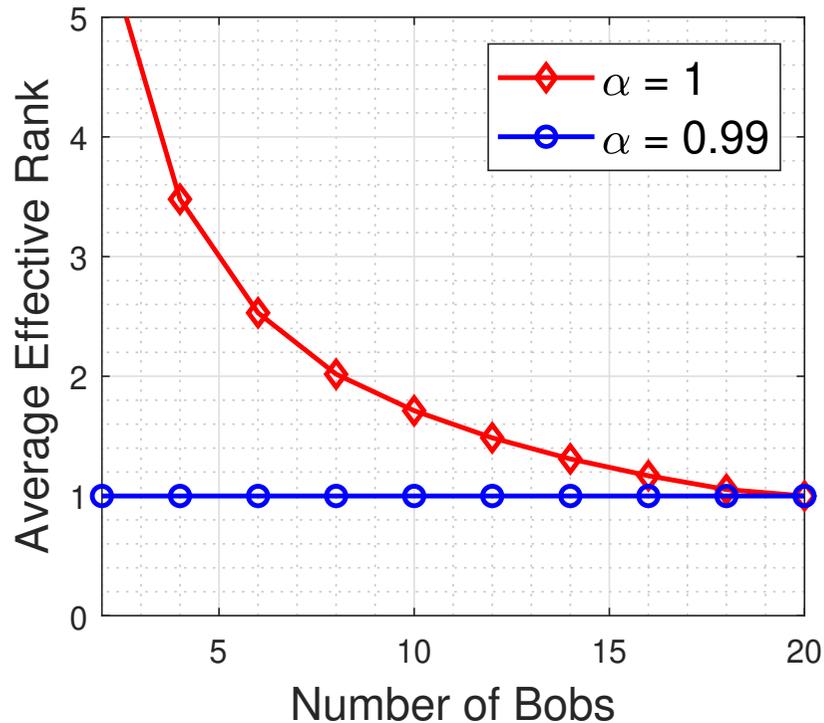
- **Theorem.** Suppose that the above problem and its dual have optimal solutions, and it attains zero duality gap. If the coefficient $\alpha \in (0, 1)$, then every optimum to (\mathcal{R}) satisfies

$$\text{rank}(\mathbf{W}_k^*) \leq 1, \quad k = 1, \dots, K.$$

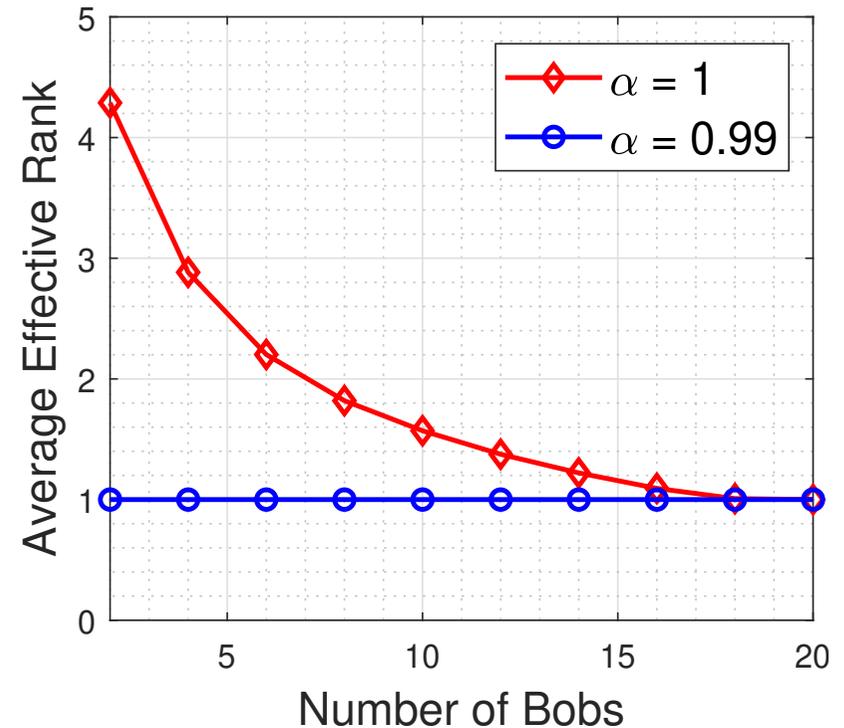
In other words, our SDR is *tight*.

- **Idea of the proof:** check the KKT conditions of (\mathcal{R}) , and the skeleton mainly follows from the Proof of Theorem 1 in [LM16]; also we exploit the SINR constraint's structure
- **Implication:** despite the original formulation (2) is very challenging to solve, our formulation (\mathcal{R}) is capable of solving (2) *almost exactly*, i.e. when α is *very close to one*, and the solution to (\mathcal{R}) must admit rank one

Experiment: Rank One Guarantee



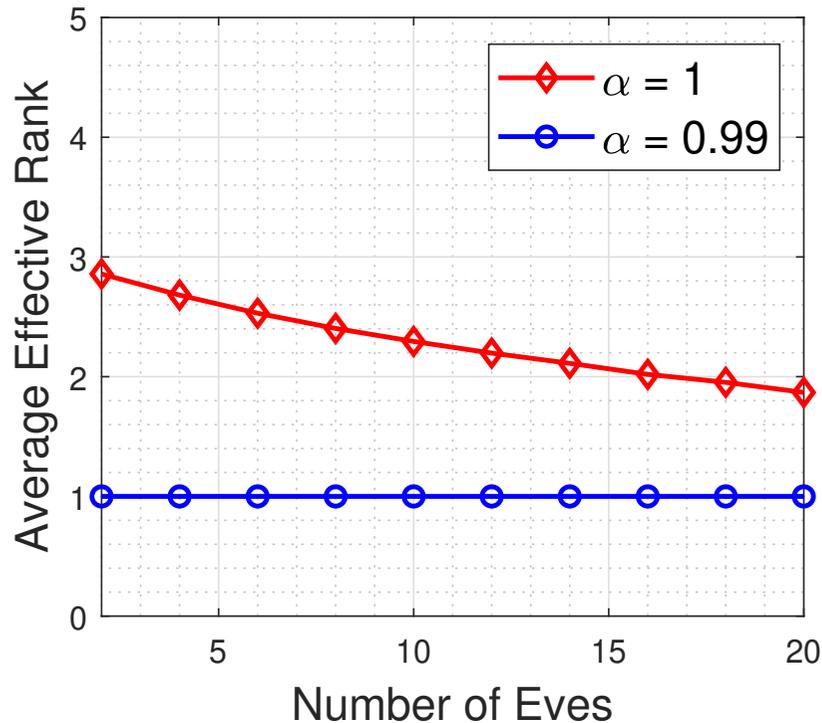
(a) When there are 6 Eves



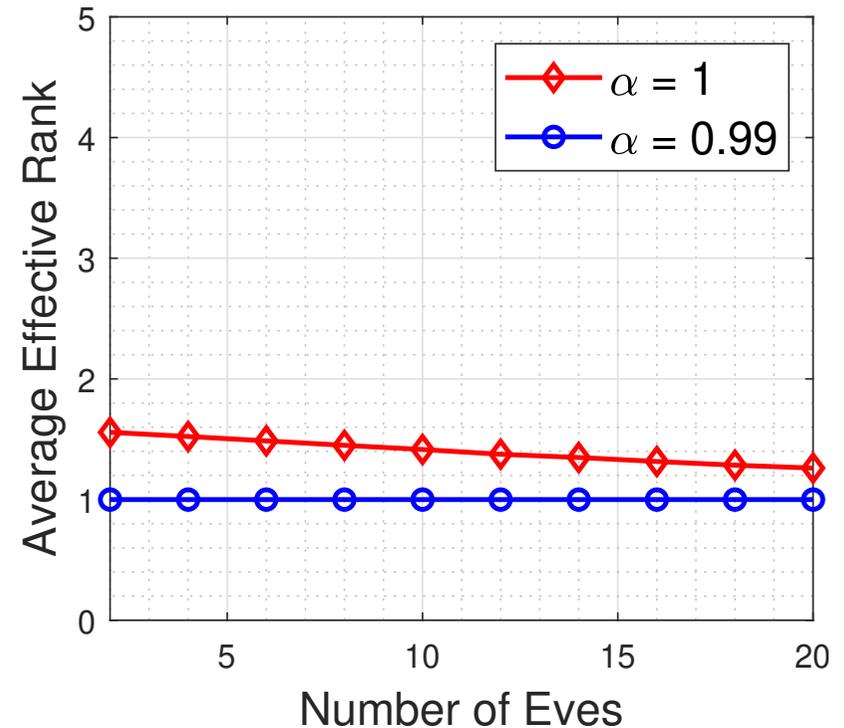
(b) When there are 12 Eves

Average effective rank of the solution to (\mathcal{R}) under different number of Bobs and Eves. $N_t = 20$, $N_r = 30$, block length $T = 800$, $\sigma^2 = 0\text{dB}$; performance bounds are $\gamma = 10\text{dB}$, $\beta = \rho = -10\text{dB}$

Experiment: Rank One Guarantee



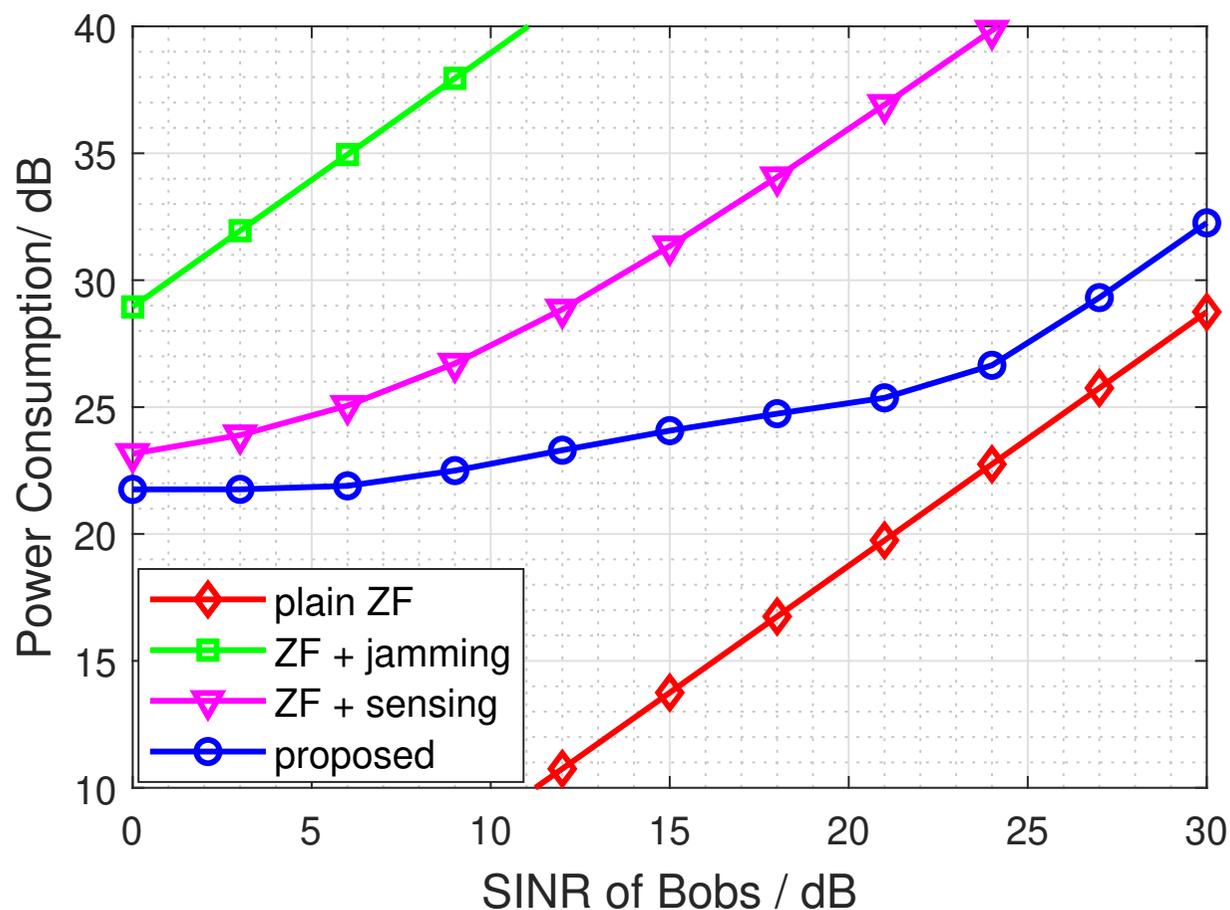
(c) When there are 6 Bobs



(d) When there are 12 Bobs

Average effective rank of the solution to (\mathcal{R}) under different number of Bobs and Eves. $N_t = 20$, $N_r = 30$, block length $T = 800$, $\sigma^2 = 0\text{dB}$; performance bounds are $\gamma = 10\text{dB}$, $\beta = \rho = -10\text{dB}$

Experiment: Power Consumption Versus Bobs' SINR γ



Average power consumption against the SINR requirement of Bobs γ . $N_t = 20$, $N_r = 30$, block length $T = 800$, $J = K = 12$, $\beta = \rho = -15$ dB; proposed method (\mathcal{R}) is run with $\alpha = 0.99$.

Conclusions

- we have reformulated the secure ISAC beamforming problem to achieve a tight SDR
- both theoretical analysis and numerical simulations agree with our argument

That's all. Thank you! Questions?

Key References

- [BO01] M. Bengtsson and B. Ottersten, *Optimal and suboptimal transmit beamforming*, Handbook of Antennas in Wireless Communications (2001).
- [HP10] Y. Huang and D. P. Palomar, *Rank-constrained separable semidefinite programming with applications to optimal beam-forming*, IEEE Trans. on Sig. Proc. **58** (2010), no. 2, 664–678.
- [LLL⁺22] F. Liu, Y.F. L., A. Li, et al., *Cramér-rao bound optimization for joint radar-communication beamforming*, IEEE Trans. on Sig. Proc. **70** (2022), 240–253.
- [LM16] Q. Li and W.K. Ma, *A new low-rank solution result for a semidefinite program problem subclass with applications to transmit beamforming optimization*, 2016 IEEE ICASSP, 2016, pp. 3446–3450.