

- A (mathematical) statement is a sentence,
or a number of carefully worded inter-related sentences,
(with mathematical content,) for which it is meaningful to say
it is true or it is false.
- Mathematical statements are 'predicates with no ambiguity'.
- A predicate with variables $x, y, z, \dots$ is
a statement 'modulo' the ambiguity of possibly one or several variables,
so that, upon the specification of the variables by
substituting concrete objects into the variables,
a statement will be yielded.

Examples of statements.

(Some of them are true statements.
Some of them are false statements.)

(a) $1 + 1 = 2$.

(b) $1 + 1 > 3$.

(c) $\sqrt{2}$ is an irrational number.

(d) 6 is divisible by 8.

(e) There exists some natural number x
such that x is divisible by 0.

(f) For any real number x , for any real number y ,
if $x < y$ then $x^2 < y^2$.

(g) For any real number x , there exists some rational number y
such that $|x - y|^2 < |x - y| + 1$.

(h) There exists some natural number y such that for any real number x , $y \leq |x|$.

Examples of predicates with some variables.

(Once the variables concerned are replaced by concrete objects, the respective predicates become statements.)

(a) $x + 1 = 2.$

(a') $x + y = 2.$

(a'') $x + 1 = y.$

(b) $x + y > 3.$

(b') $x + y > z.$

(c) $\sqrt{x}$ is an irrational number.

(d) x is divisible by $y.$

(e) x is divisible by 0.

(e') There exists some natural number x such that x is divisible by $y.$

(f) For any real number y , if $x < y$ then $x^2 < y^2.$

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(g) There exists some rational number y such that $|x - y|^2 < |x - y| + 1.$

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(g'') For any real number x , there exists some rational number y such that $|x - y|^2 < |x - y| + z.$

The same statement may be presented as one 'very condensed' sentence or several inter-related sentences. The logical relation amongst the various 'components' of the statement is indicated by the logical connectives and quantifiers.

Illustration.

These are the same statement:

(a) For any real number x , for any real number y , if $x < y$ then $x^2 < y^2$.

(b) Let x, y be real numbers. Suppose $x < y$. Then $x^2 < y^2$.

(c) Let x, y be (arbitrary) objects.

Suppose x, y are real numbers and $x < y$. Then $x^2 < y^2$.

(d) Let x, y be (arbitrary) objects.

Suppose x is a real number and y is a real number and $x < y$. Then $x^2 < y^2$.

Words circled in red are where the presence of logical connectives are signified.

Words circled in blue are where the presence of quantifiers are signified.

- Predicate calculus is a study of the properties of the logical connectives and the quantifiers.
- We want to know how the truth-hood and falsity of a 'compound statement/predicate' (formed by joining several 'shorter' statement/predicates with logical connectives and appending them with quantifiers) will be decided by the truth-hood and falsity of those 'shorter' statements/predicates in accordance with the logical connectives and quantifiers.
- Through the above, we look for some special types of statements, such as, tautology, contradiction, laws of logical equivalence, and rules of inference.

Five logical connectives

not	$\sim$	negation
and	$\wedge$	conjunction
or	$\vee$	disjunction
if ... then	$\rightarrow$	conditional
iff	$\leftrightarrow$	biconditional

Examples of usage

P : '2 is less than 3'

Q : '3 + 4 equals 8'.

$\sim P$: '2 is not less than 3'.
negation of P

$P \wedge Q$: '2 is less than 3 and
3 + 4 equals 8'.
conjunction of P, Q

$P \vee Q$: '2 is less than 3 or
3 + 4 equals 8'.
disjunction of P, Q

$P \rightarrow Q$: 'If 2 is less than 3
then 3 + 4 equals 8'.
Conditional from assumption P
to conclusion Q

$P \leftrightarrow Q$: '2 is less than 3 iff
3 + 4 equals 8'.
biconditional from P to Q

Compound statements
formed by joining P, Q with one logical connective.

- Every statement can be 'broken up' into (one or) several statements (or 'atoms') when the logical connectives in the statement are 'deleted'.
- 'Shorter' statements can be 'joint together' to give a 'longer' statement with the logical connectives. The resultant statement is called a compound statement.

Truth tables.

These are the simplest tools for studying the behaviour of the logical connectives.

P	$\sim P$
T	F
F	T

'T' stands for 'true'
'F' stands for 'false'

P	Q	$P \wedge Q$
T	T	T
T	F	F
F	T	F
F	F	F

P	Q	$P \vee Q$
T	T	T
T	F	T
F	T	T
F	F	F

Logical equivalence can be understood through truth tables.

After some work, we see that:

P	Q	$P \rightarrow Q$	$(\sim P) \vee Q$
T	T	T	T
T	F	F	F
F	T	T	T
F	F	T	T

This is one of several key features of the 'conditional' that deserve attention. They are to be found out.

Identical columns
Hence ' $P \rightarrow Q$ ', ' $(\sim P) \vee Q$ ' are logically equivalent.

P	Q	$P \leftrightarrow Q$	$(P \rightarrow Q) \wedge (Q \rightarrow P)$
T	T	T	T
T	F	F	F
F	T	F	F
F	F	T	T

Identical columns

Tautology: a (compound) statement which is always true (irrespective of the truth values of the 'atoms').

Examples

$$P \vee (\sim P)$$

$$P \rightarrow P$$

$$(P \wedge Q) \rightarrow P$$

$$P \rightarrow (P \vee Q)$$

Contradiction: a (compound) statement which is always false (irrespective of the truth values of the 'atoms').

$$P \wedge (\sim P)$$

$$P \leftrightarrow (\sim P)$$

Contingent statement: a statement which is neither a tautology nor a contradiction.

$$(P \vee Q) \rightarrow P$$

$$P \rightarrow (\sim P)$$

Rules of Inference

Some of them
are important
and have
names. They
are to be
found out.

Statements of the form

'blah-blah-blah' → 'bleh-bleh-bleh'
probably a
compound statement

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which are tautologies.

Logical equivalence understood
as tautologies.

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Statements of the form
'blah-blah-blah' ↔ 'bleh-bleh-bleh'
which are tautologies.

('blah-blah-blah', 'bleh-bleh-bleh' are logically equivalent.)

Examples of important
rules of inference.

$$(P \wedge Q) \rightarrow P$$

$$P \rightarrow (P \vee Q)$$

$$[(P \rightarrow Q) \wedge (Q \rightarrow R)] \rightarrow (P \rightarrow R)$$

Modus
ponens.

$$[(P \rightarrow Q) \wedge P] \rightarrow Q$$

$$[(P \rightarrow Q) \wedge (\neg Q)] \rightarrow (\neg P)$$

Examples of important
logical equivalence:

$$(P \rightarrow Q) \leftrightarrow [(\neg P) \vee Q]$$

$$[(P \wedge Q) \vee R] \leftrightarrow [(P \vee R) \wedge (Q \vee R)]$$

$$[\neg(P \wedge Q)] \leftrightarrow (\neg P) \vee (\neg Q)$$

$$[\neg(\neg P)] \leftrightarrow P$$