THE CIRCULANT OPERATOR IN THE BANACH ALGEBRA OF MATRICES

Raymond H- Chan XiaoQing Jin and ManChung Yeung Department of Mathematics University of Hong Kong Hong Kong May 1990

Revised July

abstract-t-active study and a padding a circulant mapping an operator completed and μ as a circulant complete matrix $c(A_n)$ that minimizes the Frobenius norm $\|A_n - C_n\|_F$ over all *n*-by-*n* circulant matrices Cn- The circulant matrix c An called the optimal circulant preconditioner has proved to be a good precedence for a general class of Toeplitz systems- \mathcal{L} and \mathcal{L} systems- \mathcal{L} give different formulations of the operator, discuss its algebraic and geometric properties and compute its operator norms in dieres in Banach algebras of matrices-sections of \sim results, we are able to give an efficient algorithm for finding the super-optimal circulant preconditioner which is defined to be the minimizer of $\|I - C_n^{-1}A_n\|_F$ over all nonsingular circulant matrices C_n .

Abbreviated Title- Circulant Operator

aar, oor aan ar rywerd aanversig van den aanversig op verske preconditioner circulant op van a tor, preconditioned conjugate gradient method

AMSMOS Sub ject Classications- F F

§1 Introduction.

Preconditioned conjugate gradient methods have been used successfully in solving many large matrix problems- Strang rst proposed using the method with circulant preconditioners for solving Toeplitz systems- avec the man with strang the change proved that for Toeplitz systems with generating functions that are positive functions in the Wiener class the method has a superlinear convergence rate due to the clustering of the eigenvalues of the preconditioned matrices-

Several circulant preconditioners have been proposed since then, see for example T. Chan and Tyrtyshinkov - For any nbyn matrix An the circulant preconditioner proposed in T- Chan called the optimal circulant preconditioner is de
ned to be the minimizer of $||C_n - A_n||_F$ over the space of all *n*-by-*n* circulant matrices C_n . Here $\|\cdot\|_F$ denotes the Frobenius norm. The circulant preconditioner given in Tyrtyshinkov [7] is defined to be the minimizer of $||I - C_n^{-1}A_n||_F$ over the space of all nonsingular $-$ 10 and is called the super-optimal circulant preconditioner-super-optimal circulant preconditioner-super-optimal circulant preconditioner-super-optimal circulant preconditioner-super-optimal circulant preconditioner-s computational point of view, these optimal circulant preconditoners are better than the one proposed in Strang $[6]$ because they are symmetric positive definite whenever A_n is. Numerical results in these papers showed that they are very good preconditioners- The analysis of the convergence rates of these preconditioned systems are given in R- Chan and R- Chan et al- and it is proved that for the same class of Toeplitz systems mentioned above, these methods converge at the same rate as the Strang's preconditioned systems.

In this paper, we study these circulant preconditioners from the operator point of view. Let $(\mathcal{M}_{n \times n}, \|\cdot\|)$ be the Banach algebra of all n-by-n matrices over the complex field, equipped with a matrix norm $\|\cdot\|$. Let $(\mathcal{C}_{n\times n},\|\cdot\|)$ be the subalgebra of all circulant matrices. We note that $\mathcal{C}_{n \times n}$ is an inverse-closed, commutative algebra. Let c be an operator defined on $({\cal M}_{n\times n},\|\cdot\|)$ such that for any A_n in ${\cal M}_{n\times n},$ $c(A_n)$ is the minimizer of $||A_n - C_n||_F$ over all C_n in $\mathcal{C}_{n \times n}$. Obviously, $c(A_n)$ is the optimal circulant preconditioner

proposed in T. Chan [4] and c is an operator from $(\mathcal{M}_{n\times n},\|\cdot\|)$ into the subalgebra $(C_{n \times n}, \|\cdot\|)$. We call c the circulant operator. In R. Chan et al. [3], we utilized this operator to analyze the convergence rate of Toeplitz systems preconditioned by super optimal circulant preconditioners-

The purpose of this paper is to discuss some other aspects of this operator- The outline of the paper is as follows. In $\S 2$, we introduce other formulations of the operator and prove some of its algebraic and geometric properties. In $\S 3$, we compute its operator norms for different Banach algebras of matrices. In $\S 4$, we apply these results to derive an algorithm for algorithm α and β algorithm is more preconditioners. Our algorithm is more to efficient than the one proposed in Trytyshinkov $[7]$.

$\S 2$ The Circulant Operator.

In this section we discuss some properties of the circulant operator-operator-operator-operator-operator-operator-operator-operator-operator-operator-operator-operator-operator-operator-operator-operator-operator-operator $\mathcal{M}_{n\times n}$, let $\delta(A_n)$ denote the diagonal matrix whose diagonal is equal to the diagonal of $\overline{\mathbf{u}}$ is given to methods for $\overline{\mathbf{v}}$

Theorem 1. Let $A_n = (a_{ij}) \in \mathcal{M}_{n \times n}$ and $c(A_n)$ be the minimizer of $||C_n - A_n||_F$ over all $C_n \in \mathcal{C}_{n \times n}$. Then $c(A_n)$ is uniquely determined by A_n . Moreover,

is given by α and α and α

$$
c(A_n) = \sum_{j=0}^{n-1} \left(\frac{1}{n} \sum_{p-q \equiv j \pmod{n}} a_{pq} Q^j \right), \qquad (1)
$$

where α is the n-distribution of the n-dist

$$
Q = \begin{bmatrix} 0 & & & & 1 \\ 1 & 0 & & & \\ & 1 & \ddots & & \\ & & \ddots & \ddots & \\ 0 & & & 1 & 0 \end{bmatrix};
$$
 (2)

ii contract and also given by an analyzing α

$$
c(A_n) = F^* \delta(F A_n F^*) F,\tag{3}
$$

Proof For the proof of i see Theorem - in Tyrtyshnikov - For ii we rst note that any circulant matrix C_n can be expressed as $F^*\Lambda_nF$, where Λ_n is a diagonal matrix containing the eigenvalues of Constanting the Frobenius of Constantine and the Frobenius of Constantine and th invariant, we have

$$
||C_n - A_n||_F = ||F^* \Lambda_n F - A_n||_F = ||\Lambda_n - F A_n F^*||_F.
$$

Thus the problem of minimizing $||C_n - A_n||_F$ over $\mathcal{C}_{n \times n}$ is equivalent to the problem of minimizing $\|\Lambda_n - FA_nF^*\|_F$ over all diagonal matrices. Since Λ_n can only affect the diagonal entries of FA_nF^* , we see that the solution for the latter problem is $\Lambda_n =$ $\delta (FA_nF^*)$. Hence $F^*\delta (FA_nF^*)F$ is the minimizer of $||C_n-A_n||_F$. It is clear from the \Box argument above that α and α and α and α are uniquely determined by α and β

we remark that the μ $\lambda = \mu$ is given by the second column of c $\lambda = \mu$ is given by λ

$$
[c(A_n)]_{j0} = \frac{1}{n} \sum_{p-q \equiv j \pmod{n}} a_{pq} = \frac{1}{n} \text{ tr } (A_n Q^{-j}), \quad j = 0, 1, \cdots, n-1,
$$
 (4)

where $\operatorname{tr}(\cdot)$ denotes the trace. By (3), the eigenvalues of $c(A_n)$ are given by $\delta (FA_nF^*)$. We also notice that the nonsingularity of A_n cannot guarantee $\delta(FA_nF^*)$ to be nonsingular. $\mathcal{N} = \{n\}$ be singular for non-singular for non-singular for non-singular for $\mathcal{N} = \{n\}$

The following Lemma is on the algebraic properties of the circulant operator-

Lemma 1.

(i) For all A_n , $B_n \in \mathcal{M}_{n \times n}$ and α , β complex scalars, $c(\alpha A_n + \beta B_n) = \alpha c(A_n) + \beta c(B_n)$ Moreover, for all $A_n \in \mathcal{M}_{n \times n}$, $c^2(A_n) = c(c(A)) = c(A_n)$. Thus c is a linear projection operator

(ii) Let
$$
A_n \in \mathcal{M}_{n \times n}
$$
, tr $(c(A_n)) = tr(A_n) = \sum_{j=0}^{n-1} \lambda_j(A_n)$, where $\lambda_j(A_n)$ are the eigenvalues of A_n .

(iii) For all $A_n \in \mathcal{M}_{n \times n}$, we have $c(A_n^*) = c(A_n)^*$.

(iv) Let $A_n \in \mathcal{M}_{n \times n}$ and $C_n \in C_{n \times n}$. Then

$$
c(C_n A_n) = C_n \cdot c(A_n),
$$

$$
c(A_n C_n) = c(A_n) \cdot C_n.
$$

Proof The proofs of $\{P\}$ were $\{P\}$ we observe the formation we omit the model $\{P\}$ we omit the set fact that $\delta(A_n^*) = (\delta(A_n))^*$, one can easily prove (iii). For the proof of (iv), see Theorem \Box in R- Chan et al- -

Next we are going to give some geometric properties of the circulant operator- For all A_n , $B_n \in \mathcal{M}_{n \times n}$, let $\langle A_n, B_n \rangle_F \equiv \frac{1}{n}$ tr $(A_n B_n^*)$. Obviously $\langle A_n, B_n \rangle_F$ is an inner product in $\mathcal{M}_{n\times n}$ and $\langle A_n, A_n \rangle_F = \frac{1}{n} ||A_n||_F^2$. It is easy to show that $\{Q^j | j = 0, \cdots, n-1\}$, where Q is given in (2), is an orthonormal basis of $(C_{n \times n}, \|\cdot\|_F)$. We show below that $A_n - c(A_n)$ is perpendicular to the space $\mathcal{C}_{n \times n}$.

Lemma 2. Let $A_n \in \mathcal{M}_{n \times n}$, then we have

(i) $\langle A_n - c(A_n), C_n \rangle_F = 0$ for all $C_n \in \mathcal{C}_{n \times n}$,

$$
(ii) \ \langle A_n, c(A_n) \rangle_F = \frac{1}{n} ||c(A_n)||_F^2,
$$

 (iii) $||A_n - c(A_n)||_F^2 = ||A_n||_F^2 - ||c(A_n)||_F^2.$

Proof. For (i), since $\{Q^j\}_{i=0}^{n-1}$ is an $j=0 \n=0$ is an orthonormal basis of $\mathcal{C}_{n \times n}$, it suffices to show that $\langle A_n - c(A_n), Q^j \rangle_F = 0$ for $j = 0, \cdots, n-1$. However, by (4) and Lemma 1 (i), we have

$$
\langle A_n - c(A_n), Q^j \rangle_F = \frac{1}{n} \operatorname{tr} \left[(A_n - c(A_n))Q^{-j} \right] = \frac{1}{n} \operatorname{tr} \left(A_n Q^{-j} \right) - \frac{1}{n} \operatorname{tr} \left(c(A_n) Q^{-j} \right)
$$

$$
= [c(A_n)]_{j0} - [c(c(A_n))]_{j0} = [c(A_n)]_{j0} - [c(A_n)]_{j0} = 0.
$$

is a following the contract of the state of the following the state of the stat

$$
||A_n - c(A_n)||_F^2 = n\langle A_n - c(A_n), A_n - c(A_n) \rangle_F = n\langle A_n - c(A_n), A_n \rangle_F
$$

= $n\langle A_n, A_n \rangle_F - n\langle c(A_n), A_n \rangle_F = ||A_n||_F^2 - ||c(A_n)||_F^2$.

$\S 3$ Spectral Properties of the Circulant Operator.

In this section we discuss some spectral properties of the circulant operator- The following theorem was rst proved for the real scalar
eld in Tyrtyshnikov - His proof uses a quation of $\left\{ -\right\}$ and our proof is a contract of motor $\left\{ \cdot\right\}$

 \mathbf{J} is Hermitian then contained contains the containing \mathbf{J}

$$
\lambda_{\min}(A_n) \leq \lambda_{\min}(c(A_n)) \leq \lambda_{\max}(c(A_n)) \leq \lambda_{\max}(A_n) ,
$$

where $\lambda_{\max}(\cdot)$ and $\lambda_{\min}(\cdot)$ denote the targest and the smallest eigenvalues respectively. In particular if μ is positive delivered, where τ positive positive depression of

 \blacksquare it is constant and an isometry and an interval when \blacksquare we know that the eigenvalues of $c(A_n)$ are given by $\delta(FA_nF^*)$. Suppose that $\delta(FA_nF^*)$ = $\lim_{n \to \infty} \alpha_0, \dots, \alpha_{n-1}$ with $\alpha_i - \alpha_{\min}(c(\alpha_n))$ and $\alpha_k - \alpha_{\max}(c(\alpha_n))$. Let e_i and e_k denote the interesting respectively-vectors respectively. The contract μ

$$
\lambda_{\max}(c(A_n)) = \lambda_k = \frac{e_k^* F A_n F^* e_k}{e_k^* e_k} \le \max_{x \neq 0} \frac{x^* F A_n F^* x}{x^* x} = \max_{x \neq 0} \frac{x^* A_n x}{x^* x} = \lambda_{\max}(A_n).
$$

Similarly

$$
\lambda_{\min}(A_n) = \min_{x \neq 0} \frac{x^* A_n x}{x^* x} = \min_{x \neq 0} \frac{x^* F A_n F^* x}{x^* x} \le \frac{e_j^* F A_n F^* e_j}{e_j^* e_j} = \lambda_j = \lambda_{\min}(c(A_n)).
$$

From the inequality above we can easily see that constants above we can easily see that constants μ positive definite. \Box

Lemma 3. For all $A_n \in M_{n \times n}$, $c(A_n A_n^*) - c(A_n)c(A_n^*)$ is a semi-positive definite matrix.

Proof. Let $A_n = (a_{jk})$ and $[F]_{jk} = \frac{1}{\sqrt{n}} \xi_j^k$, where $\xi_j = e^{-\frac{1}{n}}$. If we let

$$
D_n = c(A_n A_n^*) - c(A_n)c(A_n^*) = F^*(\delta(F A_n A_n^* F^*) - \delta(F A_n F^*)\delta(F A_n^* F^*))F,
$$

then for all $\kappa = 0, \cdots, n-1$, we have

$$
[\delta(FA_nA_n^*F^*)]_{kk} = [\delta((FA_n)(FA_n)^*)]_{kk} = \frac{1}{n} \sum_{q=0}^{n-1} (\sum_{p=0}^{n-1} a_{pq} \xi_k^p) (\sum_{p=0}^{n-1} a_{pq} \xi_k^p),
$$

and

$$
[\delta(FA_nF^*)\delta(FA_n^*F^*)]_{kk} = (\frac{1}{n}\sum_{p=0}^{n-1}\sum_{q=0}^{n-1}a_{pq}\xi_k^{p-q})\overline{(\frac{1}{n}\sum_{p=0}^{n-1}\sum_{q=0}^{n-1}a_{pq}\xi_k^{p-q})}.
$$

Hence the k-th eigenvalue of D_n is given by

$$
\lambda_k(D_n) = \frac{1}{n} \sum_{q=0}^{n-1} |\sum_{p=0}^{n-1} a_{pq} \xi_k^p|^2 - |\frac{1}{n} \sum_{p=0}^{n-1} \sum_{q=0}^{n-1} a_{pq} \xi_k^{p-q}|^2.
$$

Since

$$
|\frac{1}{n}\sum_{p=0}^{n-1}\sum_{q=0}^{n-1}a_{pq}\xi_k^{p-q}| \leq \frac{1}{n}\sum_{q=0}^{n-1}|\sum_{p=0}^{n-1}a_{pq}\xi_k^p||\xi_k^{-q}| = \frac{1}{n}\sum_{q=0}^{n-1}|\sum_{p=0}^{n-1}a_{pq}\xi_k^p|,
$$

we have

$$
\lambda_k(D_n) \geq \frac{1}{n} \sum_{q=0}^{n-1} |\sum_{p=0}^{n-1} a_{pq} \xi_k^p|^2 - \left(\frac{1}{n} \sum_{q=0}^{n-1} |\sum_{p=0}^{n-1} a_{pq} \xi_k^p| \right)^2.
$$

Let $d_{qk} = \frac{1}{n} \left| \sum_{p=0}^{n-1} a_{pq} \xi_k^p \right|$ $a_{pq}\xi_k^p$, then by Cauchy-Schwartz inequality, we have

$$
\lambda_k(D_n) \ge n \sum_{q=0}^{n-1} d_{qk}^2 - (\sum_{q=0}^{n-1} d_{qk})^2 \ge 0, \quad k = 0, \cdots, n-1.
$$

Thus D_n is semi-positive definite. \Box

Theorem 3. For all $n \geq 1$, we have

(i)
$$
||c||_1 \equiv \sup_{\|A_n\|_1=1} ||c(A_n)||_1 = 1,
$$

(ii) $||c||_{\infty} \equiv \sup ||c(A_n)||$ $\sup_{\|A_n\|_{\infty}=1} \|c(A_n)\|_{\infty}=1,$

(iii)
$$
||c||_F \equiv \sup_{\|A_n\|_F=1} ||c(A_n)||_F = 1,
$$

(iv)
$$
||c||_2 \equiv \sup_{\|A_n\|_2=1} ||c(A_n)||_2 = 1
$$
.

Proof. To prove (i), we first note that if $A_n = I$, then $||c(A_n)||_1 = ||I||_1 = 1$. For general A_n in $\mathcal{M}_{n \times n}$, we have by (1)

$$
||c(A_n)||_1 = \sum_{j=0}^{n-1} \left| \frac{1}{n} \sum_{\substack{p-q \equiv j \pmod{n}} a_{pq}} \right| \le \frac{1}{n} \sum_{j=0}^{n-1} \sum_{\substack{p-q \equiv j \pmod{n}} |a_{pq}|}
$$

= $\frac{1}{n} \sum_{i=0}^{n-1} \sum_{k=0}^{n-1} |a_{ik}| \le \frac{1}{n} \cdot n \cdot ||A_n||_1.$

Hence $||c||_1 = 1$ for all n. By a similar argument, we can prove (ii).

To prove (iii), we notice that if $A_n = \frac{1}{n}I$, then $||c(A_n)||_F = \frac{1}{n}||I||_F = 1$. For general A_n in $\mathcal{M}_{n \times n}$, by Lemma 2 (iii), we have

$$
||c(A_n)||_F^2 = ||A_n||_F^2 - ||A_n - c(A_n)||_F^2 \le ||A_n||_F^2.
$$

Thus $||c(A_n)||_F \le ||A_n||_F$. Hence $||c||_F = 1$ for all n.

To prove $\{1, 1, 2, \ldots\}$. The second intervals in the second contract of the second

$$
||c(A_n)||_2^2 = \lambda_{\max}(c(A_n)^*c(A_n)) = \lambda_{\max}(c(A_n^*)c(A_n))
$$

$$
\leq \lambda_{\max}(c(A_n^*A_n)) \leq \lambda_{\max}(A_n^*A_n) = ||A_n||_2^2,
$$

for all A_n in $\mathcal{M}_{n \times n}$. Since $||c(I)||_2 = ||I||_2 = 1$, $||c||_2 = 1$. \Box

$\S 4$ The Super-optimal Circulant Preconditioner.

In this section, we apply the results in previous sections to analyze the super-optimal circulant preconditioner proposed in Tyrtyshnikov [7]. For A_n in $\mathcal{M}_{n \times n}$, the preconditioner is defined to be the minimizer of $||I - C_n^{-1}A_n||_F$ over all nonsingular $C_n \in C_{n \times n}$. First we generalize Thoerem - in Tyrtyshnikov from the real eld to the complex eld- and with Theorem II with Theorem II with the complete $\{ - \}$ with a second complete $\{ + \}$

Theorem 4. Let $A_n \in M_{n \times n}$ be such that both A_n and $c(A_n)$ are nonsingular. Then the super-optimal circulant preconditioner for A_n exists and is equal to $c(A_nA_n^*)c(A_n^*)^{-1}$.

Proof. Instead of minimizing $||I - C_n^{-1}A_n||_F$, we consider the problem of minimizing $||I - C_n^{-1}A_n||_F$ $C_nA_n||_F$ over all nonsingular C_n in $\mathcal{C}_{n\times n}$. Letting $C_n = F^*\Lambda_n F$, we have

$$
||I - \hat{C}_n A_n||_F = ||I - F^* \Lambda_n F A_n||_F = ||I - \Lambda_n F A_n F^*||_F
$$

= tr (I - \Lambda_n F A_n F^* - F A_n^* F^* \Lambda_n^* + \Lambda_n F A_n A_n^* F^* \Lambda_n^*)
= tr (I - \Lambda_n \delta (F A_n F^*) - \delta (F A_n^* F^*)\Lambda_n^* + \Lambda_n \delta (F A_n A_n^* F^*)\Lambda_n^*).

Let Λ_n , $\delta(FA_nF^*)$ and $\delta(FA_nA_n^*F^*)$ be given by $\text{diag}(\lambda_0,\cdots,\lambda_{n-1}), \text{diag}(u_0,\cdots,u_{n-1})$ and diag (w_0, \dots, w_{n-1}) respectively. We have

$$
\min ||I - \widehat{C}_n A_n||_F = \min \{ \text{ tr } [I - \Lambda_n \delta (FA_n F^*) - \delta (FA_n^* F^*) \Lambda_n^* + \Lambda_n \delta (FA_n A_n^* F^*) \Lambda_n^*] \}
$$

$$
= \min_{\{\lambda_0, \cdots, \lambda_{n-1}\}} \sum_{k=0}^{n-1} (1 - \lambda_k u_k - \overline{u}_k \overline{\lambda}_k + \lambda_k w_k \overline{\lambda}_k).
$$

Notice that by (3) and Lemma 3, $w_k \geq u_k \bar{u}_k$ for all $k = 0, \dots, n-1$. Hence for all complex scalars λ_k , $\kappa = 0, \dots, n-1$, the terms $1 - \lambda_k u_k - u_k \lambda_k + \lambda_k w_k \lambda_k$ are nonnegative- Dierentiating them with respect to the real and imaginary parts of k and setting the derivatives to zero, we get

$$
\lambda_k = \frac{\overline{u}_k}{w_k}, \qquad k = 0, \cdots, n-1.
$$

Since An and c An are nonsingular both wk and uk are nonzero- Hence k are also nonzero. Thus the minimizer of $||I - C_n A_n||_F$ is nonsingular and is given by

$$
\widehat{C}_n = F^* \Lambda_n F = F^* \delta(F A_n^* F^*) [\delta(F A_n A_n^* F^*)]^{-1} F
$$

=
$$
(F^* \delta(F A_n^* F^*) F)(F^* \delta(F A_n A_n^* F^*) F)^{-1} = c(A_n^*) c(A_n A_n^*)^{-1}.
$$

Therefore the super-optimal circulant preconditioner is given by $C_n^{-1} = c(A_n A_n^*) c(A_n^*)^{-1}$.

 \Box

which there is the contract that α is the positive decomposition of α is the contract of α is α nonsingular- Hence the superoptimal circulant preconditioner is defined and the superstanding positive definite matrices.

When the system $A_n x = b$ is solved by preconditioned conjugate gradient method with the super-optimal circulant preconditioner $c(A_nA_n^*)(c(A_n^*)^{-1})$, then in each iteration, we have to compute a matrix-vector multiplication of the form $c(A_n^*)c(A_nA_n^*)^{-1}y$. We now derive an algorithm for finding $c(A_n^*)c(A_nA_n^*)^{-1}$. We begin by considering a general A_n that has no special structure. We first note that $c(A_n^*)c(A_nA_n^*)^{-1} = C_n$ is circulant.

Hence it is determined by its first column, which is given by

$$
\widehat{C}_n e_0 = c(A_n^*)[c(A_n A_n^*)]^{-1} e_0 = F^* \delta(F A_n^* F^*)[\delta(F A_n A_n^* F^*)]^{-1} F e_0
$$

= $F^* \delta(F A_n^* F^*)[\delta(F A_n A_n^* F^*)]^{-1}$ 1. (5)

Here 1 is the vector of all ones. To compute $\delta(F A_n^* F^*)$, it is clear from (1) that the first column $c(A_n^*)e_0$ of $c(A_n^*)$ can be computed in n^2 additions and n multiplications. Since by (3), $\delta (FA_n^*F^*)\mathbf{1} = Fc(A_n^*)e_0$, one FFT is required to obtain $\delta (FA_n^*F^*)$. To compute $\delta(FA_nA_n^*F^*) = \delta((FA_n)(FA_n)^*)$, we first need n FFTs to get FA_n , then another n^2 additions and n^2 multiplications to obtain the diagonal entries of $\delta((FA_n)(FA_n)^*)$. Now $\delta(FA_n^*F^*)[\delta(FA_nA_n^*F^*)]^{-1}$ can be obtained by n multiplications. By (5), one additional FFT is required to get $C_n\varepsilon_0$. Thus for an arbitrary n-by-n matrix A_n , C_n can be computed within $2n^2$ additions, $2n + n^2$ multiplications and $(n + 2)$ ff is.

We remark that from the computational point of view, we do not require the explicit form of \cup_n , we only need its eigenvalues and they are given by the diagonal entries of $\delta(FA_n^*F^*)[\delta(FA_nA_n^*F^*)]^{-1}$. In fact, given any vector y, C_ny can be computed by

$$
\widehat{C}_n y = F^* \delta(F A_n^* F^*) [\delta(F A_n A_n^* F^*)]^{-1} F y.
$$

Hence the last FFT in the above algorithm can usually be saved.

Next we study howaToeplitz structure can be exploited to accelerate the computation or C_{n} . The algorithm presented here is more emerent than the one proposed in Trytyshinov where a Toeplitz matrix is partitioned into the sum of low and upper triangular Toeplitz matrices-set into the will partition at Toeplitz matrix into the sum of a circulant matrix \mathbf{v}_h and a snowledge matrix φ_{h} . $\exists \varphi$, if $\|\varphi_{l}\|$ and $\|\varphi_{l}\|$ by $\exists \varphi$

$$
C_n = \frac{1}{2} \begin{bmatrix} a_0 & a_{-1} + a_{n-1} & a_{-(n-1)} + a_1 \\ a_1 + a_{-(n-1)} & a_0 & \cdots & \cdots \\ \vdots & \vdots & \ddots & \vdots \\ a_{n-1} + a_{-1} & a_0 & \cdots \end{bmatrix},
$$

and

12

$$
S_n = \frac{1}{2} \begin{bmatrix} a_0 & a_{-1} - a_{n-1} & a_{-(n-1)} - a_1 \\ a_0 & a_0 & \ddots & \ddots \\ \vdots & \vdots & \ddots & \vdots \\ a_{-1} - a_{n-1} & a_{n-1} & a_0 \end{bmatrix}.
$$

 $\begin{array}{ccc} \text{C} & \text{$ that Cne and Sne can be computed by non-cannot be computed by n multiplications and n additions-of-computed by $\det Z \cdot C_n$ is the circulant preconditioner proposed in R. Chan $|Z|$.

We will compute the first column of C_n by (5). We first compute $\delta (FA_nA_n^*F^*)$. Since $A_n = C_n + S_n = F^* \Lambda_c F + S_n$, where Λ_c is the diagonal matrix containing the eigenvalues of C_n , we have

$$
FA_nF^* = \Lambda_c + FS_nF^*.
$$
\n⁽⁶⁾

Hence

$$
\delta(FA_n A_n^* F^*) = \delta((FA_n F^*)(FA_n F^*)^*) = \delta((\Lambda_c + FS_n F^*)(\Lambda_c^* + FS_n^* F^*))
$$

$$
= \Lambda_c \Lambda_c^* + \delta(FS_n F^*)\Lambda_c^* + \Lambda_c \delta(FS_n^* F^*) + \delta(FS_n S_n^* F^*).
$$
 (7)

We now consider the terms in the right hand side of \mathcal{N}

- is the computer in the form in the computer computer α is the computer computer of α and α is the computer of α requires one FFT. Then $\Lambda_c \Lambda_c^*$ can be computed in n multiplications.
- (ii) For $\delta(FS_nF^*)\Lambda_c^*$, we know that by (3),

$$
\delta(FS_nF^*)\mathbf{1} = \delta(FS_nF^*)Fe_0 = Fc(S_n)e_0.
$$

Since Sn is skewcirculant computed in the computed in the computed in multiplications and n multiplications and additions. Then $\delta(FS_nF^*)\Lambda_c^*$ can be obtained by an additional n multiplications and one FFT-

in the third term is the conjugate term in a second term in the second term in the second term in the second t it can be computed without any work at all(iv) Finally for $\delta(FS_nS_n^*F^*)$, we have by (3) again,

$$
\delta(FS_n S_n^* F^*) \mathbf{1} = \delta(FS_n S_n^* F^*) F e_0 = Fc(S_n S_n^*) e_0.
$$
\n(8)

Thus the main work is to compute $c(S_nS_n^*)e_0$. We first find $S_nS_n^*$. We note that for all skew-circulant matrices, and in particular for S_n , they can be written as

$$
S_n = \Theta^* F^* \Lambda_s F \Theta,
$$
\n⁽⁹⁾

where $\Theta = \text{diag}(1, e^{\frac{1}{n}i}, \dots, e^{\frac{1}{n}i})$ and Λ_s is the diagonal matrix containing the eigenvalues of Sne see for instance Davis - Sne see for instance Davis - Sne see for instance Davis - Sne see can be computed in n multiplications and one FFT. Since $S_nS_n^*$ is still skew-circulant, it is determined by its first column $S_nS_n^*e_0$. By (9)

$$
S_n S_n^* e_0 = \Theta^* F^* \Lambda_s \Lambda_s^* F \Theta e_0 = \Theta^* F^* \Lambda_s \Lambda_s^* \mathbf{1},
$$

which can be computed by using one FFT and n multiplications-by using one FFT and n multiplications- $S_n S_n^* e_0$, $c(S_n S_n^*) e_0$ can be computed by using another 3n multiplications and n additions. Finally by (8), one additional FFT is required to get $\delta(FS_nS_n^*F^*)$.

By adding the four terms in (7), we see that $\delta (FA_nA_n^*F^*)$ can be obtained by using n multiplications n additions and FTS-contractions and FTS-contrac

$$
\delta(F A_n^* F^*) = \delta(F A_n F^*)^* = [\Lambda_c + \delta(F S_n F^*)]^*,
$$

where Λ_c and $\delta(FS_nF^*)$ are already computed in part (i) and (ii) above. Thus $\delta(FA_n^*F)$ can be computed in *n* additions. By (0), we see that $C_n c_0$ can be computed by an additional n multiplications and one FMT-ditributions and ρ recalling that C_1 computed in $2n$ additions and $2n$ multiplications, we see that C_n can be obtained in totally n additions n multiplications and FFTs- As remarked above the last FFT can be saved because we only need to know the eigenvalues of C_n but not its explicit form. Comparing with the algorithm proposed in Tyrtyshnikov [7] which requires 9 FFTs and O n operations we see that our method is more ecient-

- R- Chan and G- Strang Toeplitz Equations by Conjugate Gradients with Circulant Preconditioner is the statistic of the statistic order in the statistic property of the statistic order in the statistic order in
- re av Channel Circulant Fictioners for Hermitian Teoplitz Systems Systems Systems Statement Systems Systems Sys Anal- Appl- pp- -
- R- Chan X- Jin and M- Yeung The Spectra of Super-optimal Circulant Preconditioned To appear - To
- T- Chan An Optimal Circulant Preconditioner for Toeplitz Systems SIAM J- Scistatistic product in the statistic product of the computation of the c
- P- Davis Circulant Matrices John Wiley Sons Inc- New York -
- ist Grootstaan, aar arreferent for Applicitie Calculations Community Mathematical Calculations of the Calculation , _ _ _ _ *,* , *_ _ _ _ _ _ _* _ _ . _ .
- E- Tyrtyshnikov Optimal and Super-optimal Circulant Preconditioners To appear -