# THE CIRCULANT OPERATOR IN THE BANACH ALGEBRA OF MATRICES

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Abstract. We study an operator c which maps every n-by-n matrix  $A_n$  to a circulant matrix  $c(A_n)$  that minimizes the Frobenius norm  $||A_n - C_n||_F$  over all n-by-n circulant matrices  $C_n$ . The circulant matrix  $c(A_n)$ , called the optimal circulant preconditioner, has proved to be a good preconditioner for a general class of Toeplitz systems. In this paper, we give different formulations of the operator, discuss its algebraic and geometric properties and compute its operator norms in different Banach algebras of matrices. Using these results, we are able to give an efficient algorithm for finding the super-optimal circulant preconditioner which is defined to be the minimizer of  $||I - C_n^{-1}A_n||_F$  over all nonsingular circulant matrices  $C_n$ .

Abbreviated Title. Circulant Operator

**Key Words.** Toeplitz matrix, circulant matrix, optimal preconditioner, circulant operator, preconditioned conjugate gradient method

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AMS(MOS) Subject Classifications. 65F10, 65F15

# §1 Introduction.

Preconditioned conjugate gradient methods have been used successfully in solving many large matrix problems. Strang [6] first proposed using the method with circulant preconditioners for solving Toeplitz systems. R. Chan and Strang [1] then proved that for Toeplitz systems with generating functions that are positive functions in the Wiener class, the method has a super-linear convergence rate due to the clustering of the eigenvalues of the preconditioned matrices.

Several circulant preconditioners have been proposed since then, see for example T. Chan [4] and Tyrtyshinkov [7]. For any *n*-by-*n* matrix  $A_n$ , the circulant preconditioner proposed in T. Chan [4], called the *optimal circulant preconditioner*, is defined to be the minimizer of  $||C_n - A_n||_F$  over the space of all *n*-by-*n* circulant matrices  $C_n$ . Here  $|| \cdot ||_F$  denotes the Frobenius norm. The circulant preconditioner given in Tyrtyshinkov [7] is defined to be the minimizer of  $||I - C_n^{-1}A_n||_F$  over the space of all nonsingular circulant matrices  $C_n$ , and is called the *super-optimal circulant preconditioner*. From the computational point of view, these optimal circulant preconditioners are better than the one proposed in Strang [6] because they are symmetric positive definite whenever  $A_n$  is. Numerical results in these papers showed that they are very good preconditioners. The analysis of the convergence rates of these preconditioned systems are given in R. Chan [2] and R. Chan et al. [3], and it is proved that for the same class of Toeplitz systems mentioned above, these methods converge at the same rate as the Strang's preconditioned systems.

In this paper, we study these circulant preconditioners from the operator point of view. Let  $(\mathcal{M}_{n \times n}, \|\cdot\|)$  be the Banach algebra of all *n*-by-*n* matrices over the complex field, equipped with a matrix norm  $\|\cdot\|$ . Let  $(\mathcal{C}_{n \times n}, \|\cdot\|)$  be the subalgebra of all circulant matrices. We note that  $\mathcal{C}_{n \times n}$  is an inverse-closed, commutative algebra. Let *c* be an operator defined on  $(\mathcal{M}_{n \times n}, \|\cdot\|)$  such that for any  $A_n$  in  $\mathcal{M}_{n \times n}$ ,  $c(A_n)$  is the minimizer of  $\|A_n - C_n\|_F$  over all  $C_n$  in  $\mathcal{C}_{n \times n}$ . Obviously,  $c(A_n)$  is the optimal circulant preconditioner

proposed in T. Chan [4] and c is an operator from  $(\mathcal{M}_{n \times n}, \|\cdot\|)$  into the subalgebra  $(\mathcal{C}_{n \times n}, \|\cdot\|)$ . We call c the *circulant operator*. In R. Chan et al. [3], we utilized this operator to analyze the convergence rate of Toeplitz systems preconditioned by super-optimal circulant preconditioners.

The purpose of this paper is to discuss some other aspects of this operator. The outline of the paper is as follows. In  $\S 2$ , we introduce other formulations of the operator and prove some of its algebraic and geometric properties. In  $\S 3$ , we compute its operator norms for different Banach algebras of matrices. In  $\S 4$ , we apply these results to derive an algorithm for finding the super-optimal circulant preconditioner. Our algorithm is more efficient than the one proposed in Trytyshinkov [7].

#### $\S 2$ The Circulant Operator.

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In this section, we discuss some properties of the circulant operator. For any  $A_n$  in  $\mathcal{M}_{n \times n}$ , let  $\delta(A_n)$  denote the diagonal matrix whose diagonal is equal to the diagonal of the matrix  $A_n$ . We first give two methods for finding  $c(A_n)$ .

**Theorem 1.** Let  $A_n = (a_{ij}) \in \mathcal{M}_{n \times n}$  and  $c(A_n)$  be the minimizer of  $||C_n - A_n||_F$  over all  $C_n \in \mathcal{C}_{n \times n}$ . Then  $c(A_n)$  is uniquely determined by  $A_n$ . Moreover,

(i)  $c(A_n)$  is given by

$$c(A_n) = \sum_{j=0}^{n-1} \left(\frac{1}{n} \sum_{p-q \equiv j \pmod{n}} a_{pq}\right) Q^j , \qquad (1)$$

where Q is the n-by-n circulant matrix

$$Q = \begin{bmatrix} 0 & & & 1 \\ 1 & 0 & & & \\ & 1 & \ddots & & \\ & & \ddots & \ddots & \\ 0 & & & 1 & 0 \end{bmatrix};$$
 (2)

(ii)  $c(A_n)$  is also given by

$$c(A_n) = F^* \delta(F A_n F^*) F, \tag{3}$$

*Proof.* For the proof of (i), see Theorem 2.1 in Tyrtyshnikov [7]. For (ii), we first note that any circulant matrix  $C_n$  can be expressed as  $F^*\Lambda_n F$ , where  $\Lambda_n$  is a diagonal matrix containing the eigenvalues of  $C_n$ , see Davis [5]. Since the Frobenius norm is unitary-invariant, we have

$$||C_n - A_n||_F = ||F^*\Lambda_n F - A_n||_F = ||\Lambda_n - FA_n F^*||_F.$$

Thus the problem of minimizing  $||C_n - A_n||_F$  over  $C_{n \times n}$  is equivalent to the problem of minimizing  $||\Lambda_n - FA_nF^*||_F$  over all diagonal matrices. Since  $\Lambda_n$  can only affect the diagonal entries of  $FA_nF^*$ , we see that the solution for the latter problem is  $\Lambda_n =$  $\delta(FA_nF^*)$ . Hence  $F^*\delta(FA_nF^*)F$  is the minimizer of  $||C_n - A_n||_F$ . It is clear from the argument above that  $\Lambda_n$  and hence  $c(A_n)$  are uniquely determined by  $A_n$ .

We remark that by (1), the *j*-th entry in the first column of  $c(A_n)$  is given by

$$[c(A_n)]_{j0} = \frac{1}{n} \sum_{p-q \equiv j \pmod{n}} a_{pq} = \frac{1}{n} \operatorname{tr} (A_n Q^{-j}), \quad j = 0, 1, \cdots, n-1,$$
(4)

where tr (·) denotes the trace. By (3), the eigenvalues of  $c(A_n)$  are given by  $\delta(FA_nF^*)$ . We also notice that the nonsingularity of  $A_n$  cannot guarantee  $\delta(FA_nF^*)$  to be nonsingular. Hence  $c(A_n)$  may be singular for nonsingular  $A_n$ .

The following Lemma is on the algebraic properties of the circulant operator.

# Lemma 1.

- (i) For all  $A_n$ ,  $B_n \in \mathcal{M}_{n \times n}$  and  $\alpha$ ,  $\beta$  complex scalars,  $c(\alpha A_n + \beta B_n) = \alpha c(A_n) + \beta c(B_n)$ . Moreover, for all  $A_n \in \mathcal{M}_{n \times n}$ ,  $c^2(A_n) = c(c(A)) = c(A_n)$ . Thus c is a linear projection operator.
- (ii) Let  $A_n \in \mathcal{M}_{n \times n}$ , tr  $(c(A_n)) = tr(A_n) = \sum_{j=0}^{n-1} \lambda_j(A_n)$ , where  $\lambda_j(A_n)$  are the eigenvalues of  $A_n$ .

(iii) For all  $A_n \in \mathcal{M}_{n \times n}$ , we have  $c(A_n^*) = c(A_n)^*$ .

(iv) Let  $A_n \in \mathcal{M}_{n \times n}$  and  $C_n \in C_{n \times n}$ . Then

$$c(C_n A_n) = C_n \cdot c(A_n),$$
$$c(A_n C_n) = c(A_n) \cdot C_n.$$

*Proof.* The proofs of (i) and (ii) are trivial, therefore we omit them. By using (3) and the fact that  $\delta(A_n^*) = (\delta(A_n))^*$ , one can easily prove (iii). For the proof of (iv), see Theorem 2 in R. Chan et al. [3].

Next we are going to give some geometric properties of the circulant operator. For all  $A_n$ ,  $B_n \in \mathcal{M}_{n \times n}$ , let  $\langle A_n, B_n \rangle_F \equiv \frac{1}{n}$  tr  $(A_n B_n^*)$ . Obviously  $\langle A_n, B_n \rangle_F$  is an inner product in  $\mathcal{M}_{n \times n}$  and  $\langle A_n, A_n \rangle_F = \frac{1}{n} ||A_n||_F^2$ . It is easy to show that  $\{Q^j | j = 0, \dots, n-1\}$ , where Q is given in (2), is an orthonormal basis of  $(\mathcal{C}_{n \times n}, || \cdot ||_F)$ . We show below that  $A_n - c(A_n)$  is perpendicular to the space  $\mathcal{C}_{n \times n}$ .

**Lemma 2.** Let  $A_n \in \mathcal{M}_{n \times n}$ , then we have

(i)  $\langle A_n - c(A_n), C_n \rangle_F = 0$  for all  $C_n \in \mathcal{C}_{n \times n}$ ,

(*ii*) 
$$\langle A_n, c(A_n) \rangle_F = \frac{1}{n} \| c(A_n) \|_F^2$$

(*iii*)  $||A_n - c(A_n)||_F^2 = ||A_n||_F^2 - ||c(A_n)||_F^2$ .

*Proof.* For (i), since  $\{Q^j\}_{j=0}^{n-1}$  is an orthonormal basis of  $\mathcal{C}_{n \times n}$ , it suffices to show that  $\langle A_n - c(A_n), Q^j \rangle_F = 0$  for  $j = 0, \dots, n-1$ . However, by (4) and Lemma 1 (i), we have

$$\langle A_n - c(A_n), Q^j \rangle_F = \frac{1}{n} \operatorname{tr} \left[ (A_n - c(A_n))Q^{-j} \right] = \frac{1}{n} \operatorname{tr} \left( A_n Q^{-j} \right) - \frac{1}{n} \operatorname{tr} \left( c(A_n)Q^{-j} \right)$$
  
=  $[c(A_n)]_{j0} - [c(c(A_n))]_{j0} = [c(A_n)]_{j0} - [c(A_n)]_{j0} = 0.$ 

Now (ii) follows directly from (i). For (iii), we have, by parts (i) and (ii) above,

$$\|A_n - c(A_n)\|_F^2 = n \langle A_n - c(A_n), A_n - c(A_n) \rangle_F = n \langle A_n - c(A_n), A_n \rangle_F$$
$$= n \langle A_n, A_n \rangle_F - n \langle c(A_n), A_n \rangle_F = \|A_n\|_F^2 - \|c(A_n)\|_F^2.$$

### §3 Spectral Properties of the Circulant Operator.

In this section, we discuss some spectral properties of the circulant operator. The following theorem was first proved for the real scalar field in Tyrtyshnikov [7]. His proof uses equation (1) and our proof here uses equation (3).

**Theorem 2.** If  $A_n$  is Hermitian, then  $c(A_n)$  is Hermitian. Moreover, we have

$$\lambda_{\min}(A_n) \leq \lambda_{\min}(c(A_n)) \leq \lambda_{\max}(c(A_n)) \leq \lambda_{\max}(A_n)$$

where  $\lambda_{\max}(\cdot)$  and  $\lambda_{\min}(\cdot)$  denote the largest and the smallest eigenvalues respectively. In particular, if  $A_n$  is positive definite, then  $c(A_n)$  is also positive definite.

*Proof.* By Lemma 1 (iii), it is clear that  $c(A_n)$  is Hermitian when  $A_n$  is Hermitian. By (3), we know that the eigenvalues of  $c(A_n)$  are given by  $\delta(FA_nF^*)$ . Suppose that  $\delta(FA_nF^*) =$ diag $(\lambda_0, \dots, \lambda_{n-1})$  with  $\lambda_j = \lambda_{\min}(c(A_n))$  and  $\lambda_k = \lambda_{\max}(c(A_n))$ . Let  $e_j$  and  $e_k$  denote the *j*-th and the *k*-th unit vectors respectively. Since  $A_n$  is Hermitian, we have

$$\lambda_{\max}(c(A_n)) = \lambda_k = \frac{e_k^* F A_n F^* e_k}{e_k^* e_k} \le \max_{x \neq 0} \frac{x^* F A_n F^* x}{x^* x} = \max_{x \neq 0} \frac{x^* A_n x}{x^* x} = \lambda_{\max}(A_n).$$

Similarly,

$$\lambda_{\min}(A_n) = \min_{x \neq 0} \frac{x^* A_n x}{x^* x} = \min_{x \neq 0} \frac{x^* F A_n F^* x}{x^* x} \le \frac{e_j^* F A_n F^* e_j}{e_j^* e_j} = \lambda_j = \lambda_{\min}(c(A_n)).$$

From the inequality above, we can easily see that  $c(A_n)$  is positive definite when  $A_n$  is positive definite.

**Lemma 3.** For all  $A_n \in \mathcal{M}_{n \times n}$ ,  $c(A_n A_n^*) - c(A_n)c(A_n^*)$  is a semi-positive definite matrix.

*Proof.* Let  $A_n = (a_{jk})$  and  $[F]_{jk} = \frac{1}{\sqrt{n}} \xi_j^k$ , where  $\xi_j = e^{-\frac{2\pi i j}{n}}$ . If we let

$$D_n = c(A_n A_n^*) - c(A_n)c(A_n^*) = F^*(\delta(FA_n A_n^* F^*) - \delta(FA_n F^*)\delta(FA_n^* F^*))F,$$

then for all  $k = 0, \cdots, n - 1$ , we have

$$[\delta(FA_nA_n^*F^*)]_{kk} = [\delta((FA_n)(FA_n)^*)]_{kk} = \frac{1}{n} \sum_{q=0}^{n-1} (\sum_{p=0}^{n-1} a_{pq}\xi_k^p) (\sum_{p=0}^{n-1} a_{pq}\xi_k^p),$$

 $\operatorname{and}$ 

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$$[\delta(FA_nF^*)\delta(FA_n^*F^*)]_{kk} = (\frac{1}{n}\sum_{p=0}^{n-1}\sum_{q=0}^{n-1}a_{pq}\xi_k^{p-q})(\frac{1}{n}\sum_{p=0}^{n-1}\sum_{q=0}^{n-1}a_{pq}\xi_k^{p-q}).$$

Hence the k-th eigenvalue of  $D_n$  is given by

$$\lambda_k(D_n) = \frac{1}{n} \sum_{q=0}^{n-1} |\sum_{p=0}^{n-1} a_{pq} \xi_k^p|^2 - |\frac{1}{n} \sum_{p=0}^{n-1} \sum_{q=0}^{n-1} a_{pq} \xi_k^{p-q}|^2.$$

Since

$$\left|\frac{1}{n}\sum_{p=0}^{n-1}\sum_{q=0}^{n-1}a_{pq}\xi_{k}^{p-q}\right| \leq \frac{1}{n}\sum_{q=0}^{n-1}\left|\sum_{p=0}^{n-1}a_{pq}\xi_{k}^{p}\right| |\xi_{k}^{-q}| = \frac{1}{n}\sum_{q=0}^{n-1}\left|\sum_{p=0}^{n-1}a_{pq}\xi_{k}^{p}\right|,$$

we have

$$\lambda_k(D_n) \ge \frac{1}{n} \sum_{q=0}^{n-1} |\sum_{p=0}^{n-1} a_{pq} \xi_k^p|^2 - \left(\frac{1}{n} \sum_{q=0}^{n-1} |\sum_{p=0}^{n-1} a_{pq} \xi_k^p|\right)^2.$$

Let  $d_{qk} = \frac{1}{n} \sum_{p=0}^{n-1} a_{pq} \xi_k^p$ , then by Cauchy-Schwartz inequality, we have

$$\lambda_k(D_n) \ge n \sum_{q=0}^{n-1} d_{qk}^2 - (\sum_{q=0}^{n-1} d_{qk})^2 \ge 0, \quad k = 0, \cdots, n-1.$$

Thus  $D_n$  is semi-positive definite.

**Theorem 3.** For all  $n \ge 1$ , we have

(i) 
$$||c||_1 \equiv \sup_{||A_n||_1=1} ||c(A_n)||_1 = 1,$$

- (ii)  $||c||_{\infty} \equiv \sup_{||A_n||_{\infty}=1} ||c(A_n)||_{\infty} = 1,$
- (iii)  $||c||_F \equiv \sup_{||A_n||_F=1} ||c(A_n)||_F = 1,$
- (iv)  $||c||_2 \equiv \sup_{||A_n||_2=1} ||c(A_n)||_2 = 1$ .

*Proof.* To prove (i), we first note that if  $A_n = I$ , then  $||c(A_n)||_1 = ||I||_1 = 1$ . For general  $A_n$  in  $\mathcal{M}_{n \times n}$ , we have by (1)

$$\begin{aligned} \|c(A_n)\|_1 &= \sum_{j=0}^{n-1} \left| \frac{1}{n} \sum_{p-q \equiv j \pmod{n}} a_{pq} \right| \le \frac{1}{n} \sum_{j=0}^{n-1} \sum_{p-q \equiv j \pmod{n}} |a_{pq}| \\ &= \frac{1}{n} \sum_{i=0}^{n-1} \sum_{k=0}^{n-1} |a_{ik}| \le \frac{1}{n} \cdot n \cdot \|A_n\|_1. \end{aligned}$$

Hence  $||c||_1 = 1$  for all *n*. By a similar argument, we can prove (ii).

To prove (iii), we notice that if  $A_n = \frac{1}{n}I$ , then  $||c(A_n)||_F = \frac{1}{n}||I||_F = 1$ . For general  $A_n$  in  $\mathcal{M}_{n \times n}$ , by Lemma 2 (iii), we have

$$||c(A_n)||_F^2 = ||A_n||_F^2 - ||A_n - c(A_n)||_F^2 \le ||A_n||_F^2.$$

Thus  $||c(A_n)||_F \leq ||A_n||_F$ . Hence  $||c||_F = 1$  for all n.

To prove (iv), by Lemma 1 (iii), Lemma 3 and Theorem 2, we have

$$\|c(A_n)\|_{2}^{2} = \lambda_{\max}(c(A_n)^{*}c(A_n)) = \lambda_{\max}(c(A_n^{*})c(A_n))$$
$$\leq \lambda_{\max}(c(A_n^{*}A_n)) \leq \lambda_{\max}(A_n^{*}A_n) = \|A_n\|_{2}^{2},$$

for all  $A_n$  in  $\mathcal{M}_{n \times n}$ . Since  $||c(I)||_2 = ||I||_2 = 1$ ,  $||c||_2 = 1$ .

# §4 The Super-optimal Circulant Preconditioner.

In this section, we apply the results in previous sections to analyze the super-optimal circulant preconditioner proposed in Tyrtyshnikov [7]. For  $A_n$  in  $\mathcal{M}_{n \times n}$ , the preconditioner is defined to be the minimizer of  $||I - C_n^{-1}A_n||_F$  over all nonsingular  $C_n \in \mathcal{C}_{n \times n}$ . First, we generalize Theorem 4.1 in Tyrtyshnikov [7] from the real field to the complex field. As with Theorem 2, his proof uses equation (1) and ours uses equation (3).

**Theorem 4.** Let  $A_n \in \mathcal{M}_{n \times n}$  be such that both  $A_n$  and  $c(A_n)$  are nonsingular. Then the super-optimal circulant preconditioner for  $A_n$  exists and is equal to  $c(A_nA_n^*)c(A_n^*)^{-1}$ .

Proof. Instead of minimizing  $||I - C_n^{-1}A_n||_F$ , we consider the problem of minimizing  $||I - \widehat{C}_n A_n||_F$  over all nonsingular  $\widehat{C}_n$  in  $\mathcal{C}_{n \times n}$ . Letting  $\widehat{C}_n = F^* \Lambda_n F$ , we have

$$\|I - \hat{C}_n A_n\|_F = \|I - F^* \Lambda_n F A_n\|_F = \|I - \Lambda_n F A_n F^*\|_F$$
$$= \operatorname{tr} \left(I - \Lambda_n F A_n F^* - F A_n^* F^* \Lambda_n^* + \Lambda_n F A_n A_n^* F^* \Lambda_n^*\right)$$
$$= \operatorname{tr} \left(I - \Lambda_n \delta(F A_n F^*) - \delta(F A_n^* F^*) \Lambda_n^* + \Lambda_n \delta(F A_n A_n^* F^*) \Lambda_n^*\right).$$

Let  $\Lambda_n$ ,  $\delta(FA_nF^*)$  and  $\delta(FA_nA_n^*F^*)$  be given by  $\operatorname{diag}(\lambda_0, \dots, \lambda_{n-1})$ ,  $\operatorname{diag}(u_0, \dots, u_{n-1})$ and  $\operatorname{diag}(w_0, \dots, w_{n-1})$  respectively. We have

$$\min \|I - \widehat{C}_n A_n\|_F = \min \left\{ \operatorname{tr} \left[I - \Lambda_n \delta(F A_n F^*) - \delta(F A_n^* F^*) \Lambda_n^* + \Lambda_n \delta(F A_n A_n^* F^*) \Lambda_n^*\right] \right\}$$
$$= \min_{\{\lambda_0, \cdots, \lambda_{n-1}\}} \sum_{k=0}^{n-1} (1 - \lambda_k u_k - \overline{u}_k \overline{\lambda}_k + \lambda_k w_k \overline{\lambda}_k).$$

Notice that by (3) and Lemma 3,  $w_k \ge u_k \bar{u}_k$  for all  $k = 0, \dots, n-1$ . Hence for all complex scalars  $\lambda_k$ ,  $k = 0, \dots, n-1$ , the terms  $1 - \lambda_k u_k - \bar{u}_k \overline{\lambda}_k + \lambda_k w_k \overline{\lambda}_k$  are nonnegative. Differentiating them with respect to the real and imaginary parts of  $\lambda_k$  and setting the derivatives to zero, we get

$$\lambda_k = rac{\overline{u}_k}{w_k}, \qquad k = 0, \cdots, n-1.$$

Since  $A_n$  and  $c(A_n)$  are nonsingular, both  $w_k$  and  $u_k$  are nonzero. Hence  $\lambda_k$  are also nonzero. Thus the minimizer of  $||I - \widehat{C}_n A_n||_F$  is nonsingular and is given by

$$\hat{C}_n = F^* \Lambda_n F = F^* \delta(FA_n^*F^*) [\delta(FA_n A_n^*F^*)]^{-1} F$$
$$= (F^* \delta(FA_n^*F^*)F) (F^* \delta(FA_n A_n^*F^*)F)^{-1} = c(A_n^*) c(A_n A_n^*)^{-1}.$$

Therefore the super-optimal circulant preconditioner is given by  $\hat{C}_n^{-1} = c(A_n A_n^*)c(A_n^*)^{-1}$ .

We remark that by Theorem 2, if  $A_n$  is Hermitian positive definite, then  $c(A_n)$  is nonsingular. Hence the super-optimal circulant preconditioner is defined for all Hermitian positive definite matrices.

When the system  $A_n x = b$  is solved by preconditioned conjugate gradient method with the super-optimal circulant preconditioner  $c(A_n A_n^*)c(A_n^*)^{-1}$ , then in each iteration, we have to compute a matrix-vector multiplication of the form  $c(A_n^*)c(A_n A_n^*)^{-1}y$ . We now derive an algorithm for finding  $c(A_n^*)c(A_n A_n^*)^{-1}$ . We begin by considering a general  $A_n$  that has no special structure. We first note that  $c(A_n^*)c(A_n A_n^*)^{-1} = \widehat{C}_n$  is circulant. Hence it is determined by its first column, which is given by

$$\hat{C}_{n}e_{0} = c(A_{n}^{*})[c(A_{n}A_{n}^{*})]^{-1}e_{0} = F^{*}\delta(FA_{n}^{*}F^{*})[\delta(FA_{n}A_{n}^{*}F^{*})]^{-1}Fe_{0}$$
$$= F^{*}\delta(FA_{n}^{*}F^{*})[\delta(FA_{n}A_{n}^{*}F^{*})]^{-1}\mathbf{1}.$$
(5)

Here **1** is the vector of all ones. To compute  $\delta(FA_n^*F^*)$ , it is clear from (1) that the first column  $c(A_n^*)e_0$  of  $c(A_n^*)$  can be computed in  $n^2$  additions and n multiplications. Since by (3),  $\delta(FA_n^*F^*)\mathbf{1} = Fc(A_n^*)e_0$ , one FFT is required to obtain  $\delta(FA_n^*F^*)$ . To compute  $\delta(FA_nA_n^*F^*) = \delta((FA_n)(FA_n)^*)$ , we first need n FFTs to get  $FA_n$ , then another  $n^2$ additions and  $n^2$  multiplications to obtain the diagonal entries of  $\delta((FA_n)(FA_n)^*)$ . Now  $\delta(FA_n^*F^*)[\delta(FA_nA_n^*F^*)]^{-1}$  can be obtained by n multiplications. By (5), one additional FFT is required to get  $\widehat{C}_n e_0$ . Thus for an arbitrary n-by-n matrix  $A_n$ ,  $\widehat{C}_n$  can be computed within  $2n^2$  additions,  $2n + n^2$  multiplications and (n + 2) FFTs.

We remark that from the computational point of view, we do not require the explicit form of  $\widehat{C}_n$ , we only need its eigenvalues and they are given by the diagonal entries of  $\delta(FA_n^*F^*)[\delta(FA_nA_n^*F^*)]^{-1}$ . In fact, given any vector y,  $\widehat{C}_n y$  can be computed by

$$\widehat{C}_n y = F^* \delta(FA_n^*F^*) [\delta(FA_n A_n^*F^*)]^{-1} F y.$$

Hence the last FFT in the above algorithm can usually be saved.

Next we study how a Toeplitz structure can be exploited to accelerate the computation of  $\hat{C}_n$ . The algorithm presented here is more efficient than the one proposed in Trytyshinov [7] where a Toeplitz matrix is partitioned into the sum of low and upper triangular Toeplitz matrices. Here we will partition a Toeplitz matrix into the sum of a circulant matrix  $C_n$ and a skew-circulant matrix  $S_n$ . Let  $A = (a_{ij}) = (a_{i-j})$  be Toeplitz, define  $C_n$  and  $S_n$  by

$$C_{n} = \frac{1}{2} \begin{bmatrix} a_{0} & a_{-1} + a_{n-1} & & a_{-(n-1)} + a_{1} \\ a_{1} + a_{-(n-1)} & a_{0} & & & \\ & & \ddots & \ddots & \\ & & & \ddots & \ddots & \\ a_{n-1} + a_{-1} & & & & a_{0} \end{bmatrix}$$

and

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$$S_n = \frac{1}{2} \begin{bmatrix} a_0 & a_{-1} - a_{n-1} & a_{-(n-1)} - a_1 \\ -(a_{-(n-1)} - a_1) & a_0 & & \\ & \ddots & \ddots & \\ & & \ddots & \ddots & \\ -(a_{-1} - a_{n-1}) & & & a_0 \end{bmatrix}.$$

Clearly  $C_n$  is circulant and  $S_n$  is skew-circulant. Moreover, we have  $A_n = C_n + S_n$  and that  $C_n e_0$  and  $S_n e_0$  can be computed by 2n multiplications and 2n additions. We remark that  $2 \cdot C_n$  is the circulant preconditioner proposed in R. Chan [2].

We will compute the first column of  $\widehat{C}_n$  by (5). We first compute  $\delta(FA_nA_n^*F^*)$ . Since  $A_n = C_n + S_n = F^*\Lambda_c F + S_n$ , where  $\Lambda_c$  is the diagonal matrix containing the eigenvalues of  $C_n$ , we have

$$FA_nF^* = \Lambda_c + FS_nF^*. \tag{6}$$

Hence

$$\delta(FA_n A_n^* F^*) = \delta((FA_n F^*)(FA_n F^*)^*) = \delta((\Lambda_c + FS_n F^*)(\Lambda_c^* + FS_n^* F^*))$$
$$= \Lambda_c \Lambda_c^* + \delta(FS_n F^*) \Lambda_c^* + \Lambda_c \delta(FS_n^* F^*) + \delta(FS_n S_n^* F^*).$$
(7)

We now consider the terms in the right hand side of (7) one by one.

- (i) For the first term in (7), we first compute  $\Lambda_c$  by using  $\Lambda_c \mathbf{1} = \Lambda_c F e_0 = F C_n e_0$ . That requires one FFT. Then  $\Lambda_c \Lambda_c^*$  can be computed in *n* multiplications.
- (ii) For  $\delta(FS_nF^*)\Lambda_c^*$ , we know that by (3),

$$\delta(FS_nF^*)\mathbf{1} = \delta(FS_nF^*)Fe_0 = Fc(S_n)e_0.$$

Since  $S_n$  is skew-circulant,  $c(S_n)e_0$  can be computed in 3n multiplications and n additions. Then  $\delta(FS_nF^*)\Lambda_c^*$  can be obtained by an additional n multiplications and one FFT.

(iii) The third term in (7) is just the conjugate transpose of the second term in (7). Hence it can be computed without any work at all.

$$\delta(FS_nS_n^*F^*)\mathbf{1} = \delta(FS_nS_n^*F^*)Fe_0 = Fc(S_nS_n^*)e_0.$$
(8)

Thus the main work is to compute  $c(S_n S_n^*)e_0$ . We first find  $S_n S_n^*$ . We note that for all skew-circulant matrices, and in particular for  $S_n$ , they can be written as

$$S_n = \Theta^* F^* \Lambda_s F \Theta, \tag{9}$$

where  $\Theta = \text{diag}(1, e^{\frac{\pi}{n}i}, \dots, e^{\frac{(n-1)\pi}{n}i})$  and  $\Lambda_s$  is the diagonal matrix containing the eigenvalues of  $S_n$ , see for instance, Davis [5]. Because  $\Lambda_s \mathbf{1} = \Lambda_s F \Theta e_0 = F \Theta S_n e_0$ ,  $\Lambda_s$  can be computed in *n* multiplications and one FFT. Since  $S_n S_n^*$  is still skew-circulant, it is determined by its first column  $S_n S_n^* e_0$ . By (9),

$$S_n S_n^* e_0 = \Theta^* F^* \Lambda_s \Lambda_s^* F \Theta e_0 = \Theta^* F^* \Lambda_s \Lambda_s^* \mathbf{1},$$

which can be computed by using one FFT and 2n multiplications. Once we know  $S_n S_n^* e_0$ ,  $c(S_n S_n^*) e_0$  can be computed by using another 3n multiplications and n additions. Finally by (8), one additional FFT is required to get  $\delta(FS_n S_n^*F^*)$ .

By adding the four terms in (7), we see that  $\delta(FA_nA_n^*F^*)$  can be obtained by using 11*n* multiplications, 5*n* additions and 5 FFTs. We note that by (6),

$$\delta(FA_n^*F^*) = \delta(FA_nF^*)^* = [\Lambda_c + \delta(FS_nF^*)]^*,$$

where  $\Lambda_c$  and  $\delta(FS_nF^*)$  are already computed in part (i) and (ii) above. Thus  $\delta(FA_n^*F)$ can be computed in n additions. By (5), we see that  $\hat{C}_n e_0$  can be computed by an additional n multiplications and one FFT. Finally, by recalling that  $C_n e_0$  and  $S_n e_0$  are computed in 2n additions and 2n multiplications, we see that  $\hat{C}_n$  can be obtained in totally 8n additions, 14n multiplications and 6 FFTs. As remarked above, the last FFT can be saved because we only need to know the eigenvalues of  $\hat{C}_n$  but not its explicit form. Comparing with the algorithm proposed in Tyrtyshnikov [7] which requires 9 FFTs and O(n) operations, we see that our method is more efficient.

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