

Next goal:

Decomposition of state space

Def. we say $x \rightarrow y$ (x leads to y)
if $f_{xy} > 0$.

Lemma: (i) $x \rightarrow y$ iff $P^n(x, y) > 0$ for some $n \geq 1$.

pf. $f_{xy} = P_x(\exists T_y < \infty) > 0$
 $\Leftrightarrow P^n(x, y) > 0$ for some $n \geq 1$.

(ii) If $x \rightarrow y$ & $y \rightarrow z$, then $x \rightarrow z$.

pf. $x \rightarrow y \Rightarrow \exists n \geq 1$, s.t. $P^n(x, y) > 0$
 $y \rightarrow z \Rightarrow \exists k \geq 1$, s.t. $P^k(y, z) > 0$

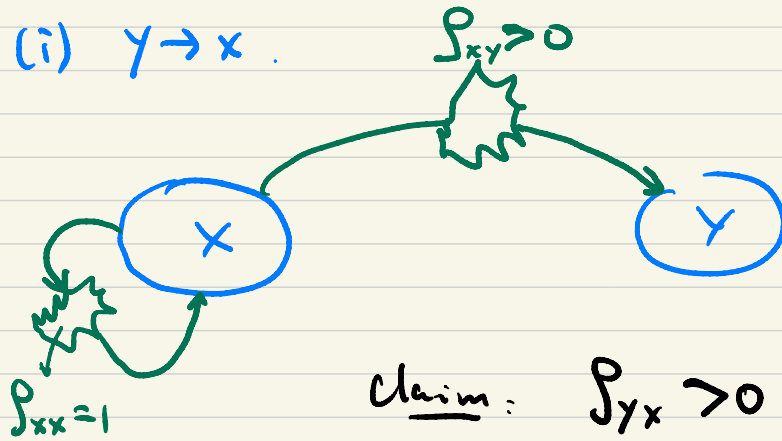
$$\underbrace{P^{n+k}(x, z)}_{= \underbrace{P^n}_{> 0} \cdot \underbrace{P^k}_{> 0}} \geq \underbrace{P^n}_{> 0}(x, y) \underbrace{P^k}_{> 0}(y, z) > 0$$

$\therefore x \rightarrow z$. #

Prop.

$$\left. \begin{array}{l} X \text{ recurrent} \\ (\mathbb{P}_{xx} = 1) \\ X \rightarrow Y \\ (\mathbb{P}_{xy} > 0) \end{array} \right\} \Rightarrow \begin{array}{l} (i) Y \rightarrow X \\ (ii) Y \text{ also recurrent} \\ (iii) \mathbb{P}_{xy} = 1 = \mathbb{P}_{yx} \end{array}$$

pf. (i) $Y \rightarrow X$.



$$\mathbb{P}_{xx} = 1 \iff \mathbb{P}_x(T_x = \infty) = 0$$

claim: $\mathbb{P}_{yx} > 0$

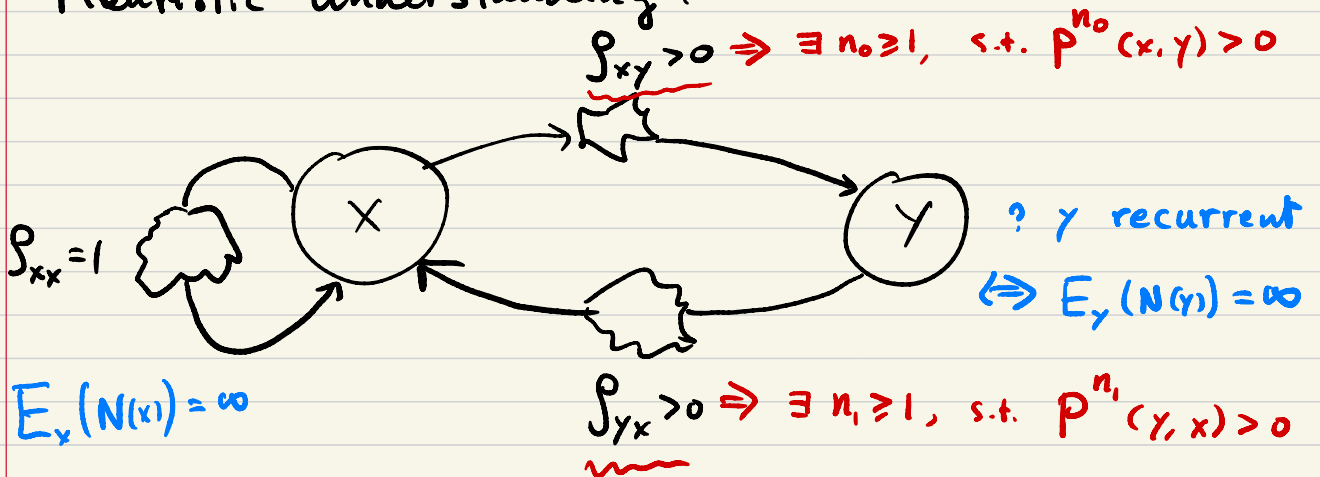
otherwise: $\mathbb{P}_{yx} = 0$

contradiction to " $\mathbb{P}_{xx} = 1$ "

$$\mathbb{P}_{yx} = 0 \iff \mathbb{P}_y(T_x = \infty) = 1$$

(ii) Y is recurrent.

Heuristic understanding:



Mathematical proof :

$$\infty \stackrel{ok}{=} E_y(N(y)) = \sum_{n=1}^{\infty} P^n(y, y)$$

$P^1(x,y) + P^2(x,y) + \dots$
part of the above
 $P_{n_1+1+n_0}^{n_1+n_0+2}(x,y) + P$

(x,y) -entry of
Product of P^{m_1}, P^{m_2} & P^{m_3}
↓
 $P^{m_1+m_2+m_3}(x,y)$

$$\geq \sum_{k=1}^{\infty} P^{n_1+k+n_0}(y, y)$$

$$= (P^{n_1} \cdot P^k \cdot P^{n_0})(y, y)$$

$$= \sum_{x_1, x_2 \in S} P^{m_1}(x, x_1) P^{m_2}(x_1, x_2) P^{m_3}(x_2, y)$$

$$\geq P^{n_1}(y, x) P^k(x, x) P^{n_0}(x, y)$$

$$\geq \underbrace{P^{n_1}(y, x)}_{\substack{\text{finite } (\leq 1) \\ > 0}} \left(\sum_{k=1}^{\infty} P^k(x, x) \right) \underbrace{P^{n_0}(x, y)}_{\substack{\text{finite } (\leq 1) \\ > 0}}$$

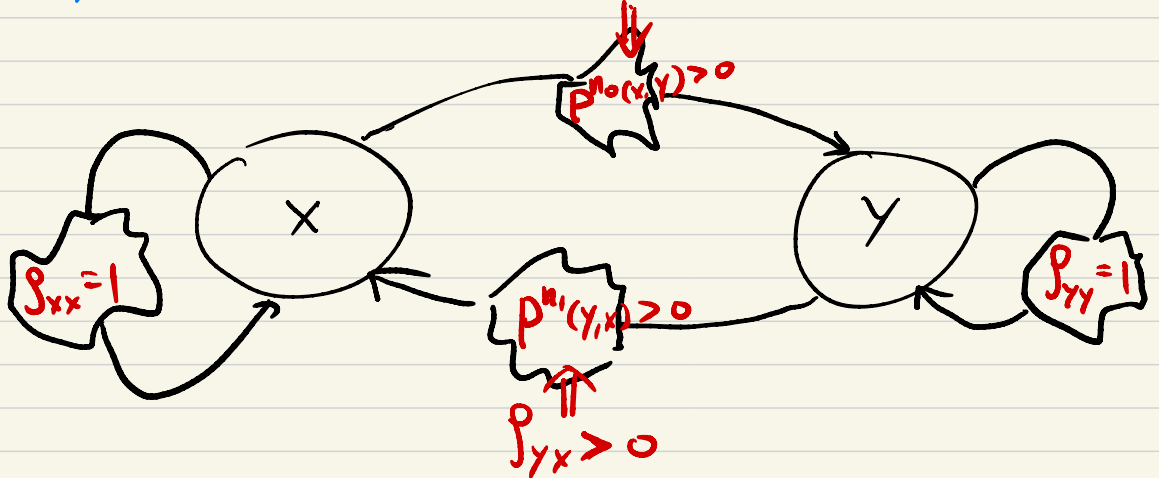
$$= E_x(N(x))$$

$$= \infty \quad (\because x \text{ is recurrent})$$

$$= \infty.$$

$\therefore y$ is recurrent.

(iii) $P_{yx} = 1$. Indeed $P_{xy} > 0$



otherwise, $P_{yx} < 1$,

i.e. $\underbrace{1 - P_{yx}} > 0$

$$= P_y(T_x = \infty)$$

with this (+)-prob. the chain from y
will Never visit $x \Rightarrow \underline{P_x(T_x = \infty) > 0}$

exchang x and y

similarly

$P_{xy} = 1$ hold true.

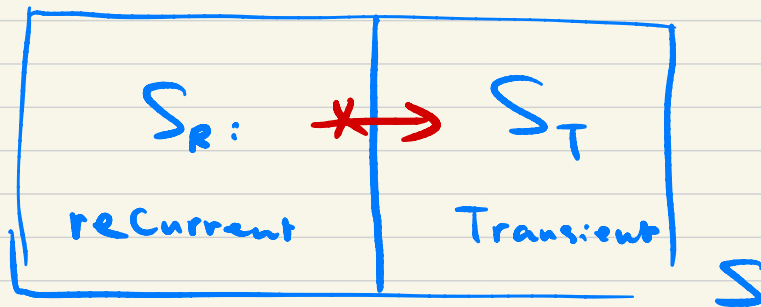
Contradiction to " x is recurrent"

#

$$S = S_R \cup S_T$$

↑
state space

↑
disjoint union



no recurrent state can lead to
any transient state

$S_R = \{\text{all recurrent states}\}$
is "closed". "

Def. $C \subseteq S$ is closed if

$$\underline{P_{xy} = 0}, \quad \forall x \in C, \forall y \notin C$$

($x \not\rightarrow y$)

i.e. no state inside C leads to
any state outside C .

Remarks:

① The following statements are equivalent:

$$(a) C \text{ is closed } (p_{xy} = 0, \forall x \in C, \forall y \notin C)$$

$$(b) P^n(x, y) = 0, \forall x \in C, \forall y \notin C, \forall n \geq 1,$$

$$(c) P(x, y) = 0, \forall x \in C, \forall y \notin C.$$

Pf. direct to see: $(a) \Leftrightarrow (b) \Rightarrow (c)$

it suffices to show: $(c) \Rightarrow (b)$

Fix $x \in C$ and $y \notin C$,

$$\begin{aligned} P^2(x, y) &= \sum_{x_1 \in S} P(x, x_1) P(x_1, y) \\ &= \sum_{\substack{x_1 \in S \\ x_1 \in C}} P(x, x_1) \underbrace{P(x_1, y)}_{\substack{\in C \notin C \\ \cancel{\in C \notin C}}} \\ &\quad + \sum_{\substack{x_1 \in S \\ x_1 \notin C}} \underbrace{P(x, x_1)}_{\substack{\in C \notin C \\ \cancel{\in C \notin C}}} P(x_1, y) \\ &= 0 + 0 = 0 \end{aligned}$$

repeatedly use the same argument

$$\underline{P^n(x, y) = 0, \forall n \geq 2.}$$

② If C is closed, $x \in C$, and $P(x, y) > 0$
then $y \in C$.

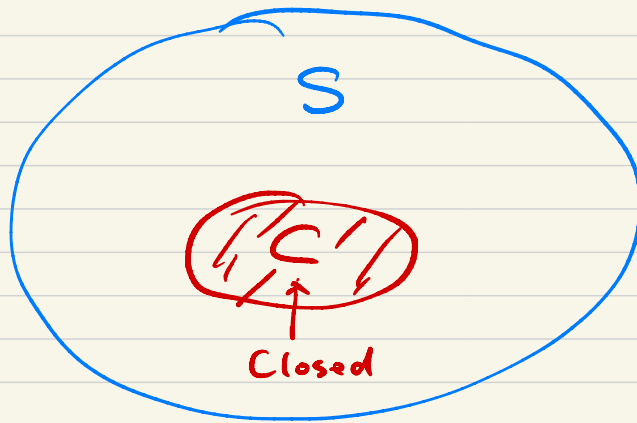
Pf. Contradiction + CC) in ①.

③ If $C \subseteq S$ is closed, then a MC

$$\{X_n\}_{n=0}^{\infty}$$

over S can also be regarded as a

Markov chain restricted to the state space C .



$$X_0, X_1, X_2, \dots$$

$$X_0(\Omega) \in C, \text{ i.e. } X_{0,n} \in C$$

Def. ① A closed set C is irreducible if

$$x \rightarrow y, \quad \forall x, y \in C$$

i.e. any two states in C can lead to each other.

② A Markov chain $\{X_n\}_{n=0}^{\infty}$ is irreducible if the state space S is irreducible.

Continue : $S = S_R \cup S_T$

\uparrow closed set \uparrow disjoint

1° $S_R \neq \emptyset$,

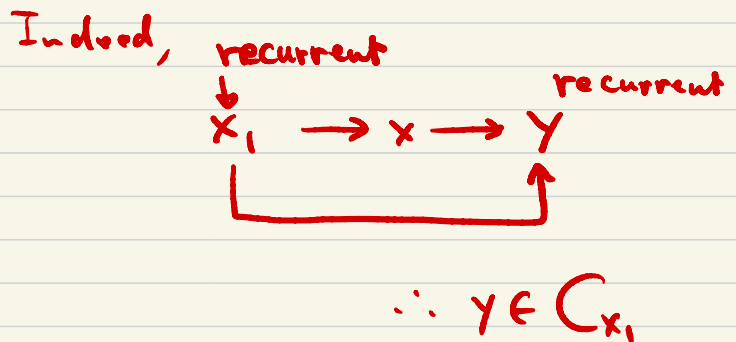
then $\exists x_1$ recurrent $\in S_R$

$$C_{x_1} \stackrel{\text{def.}}{=} \left\{ x \in S_R : \underbrace{x_1 \rightarrow x}_{\text{recurrent}} \right\} \neq \emptyset$$

(contains x_1)

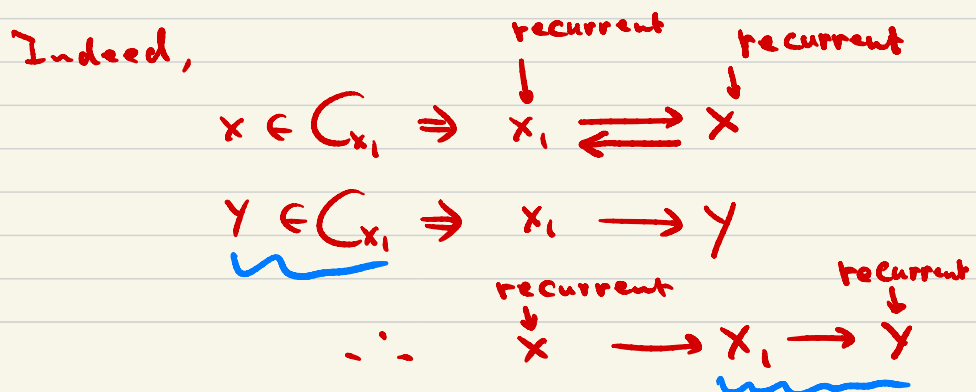
Claim : C_{x_1} is closed and irreducible

Pf. : 1° "Closed" : $\left. \begin{matrix} x \in C_{x_1} \\ x \rightarrow y \end{matrix} \right\} \Rightarrow y \in C_{x_1}$



2° "Irreducible" :

$$\left. \begin{matrix} x \in C_{x_1} \\ y \in C_{x_1} \end{matrix} \right\} \Rightarrow x \rightarrow y$$





$x \rightarrow y$. #

Rk: C_{x_1} is the "largest" closed & irreducible set containing x_1 .

2° $S_R \setminus \underline{C_{x_1}} = \emptyset$ or $\neq \emptyset$

If $S_R \setminus C_{x_1} \neq \emptyset$

then $\exists x_2 \in S_R \setminus C_{x_1}$ (i.e. $x_2 \in S_R, x_2 \notin C_{x_1}$)

define again, $C_{x_2} \stackrel{\text{def.}}{=} \{x \in S_R : x_2 \rightarrow x\}$

is closed and irreducible.

claim: $C_{x_2} \cap C_{x_1} = \emptyset$.

pf: Exercise. (Use contradiction)

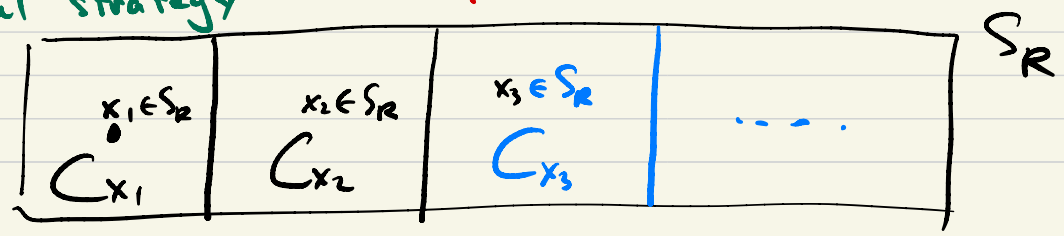
\Downarrow
 $C_{x_1} = C_{x_2}$
 \Downarrow
 Contradiction to $x_2 \notin C_{x_1}$.

3° Continue the process:

$S_R \setminus \left(\bigcup_{i=1}^2 C_{x_i} \right) = \emptyset$ or $\neq \emptyset$

↑ disjoint union of closed & irreducible sets of recurrent

general strategy



Finally,

Thm. Assume $S_R \neq \emptyset$, then

$$S_R = \bigcup_{i=1}^k C_{x_i} \quad (k < \infty \text{ or } k = \infty)$$

↑ disjoint union ↑ closed & irreducible set of recurrent states.

Thm. 1°. If C is irreducible & closed set,
then

either $C \subset S_R$ or $C \subset S_T$.

2°. If C is a finite, irreducible closed set,
then

$$C \subset S_R.$$

Pf. : 1°. Contradiction

2°. Consequence of the fact

“ A MC over the finite state space must contain at least one recurrent state.”

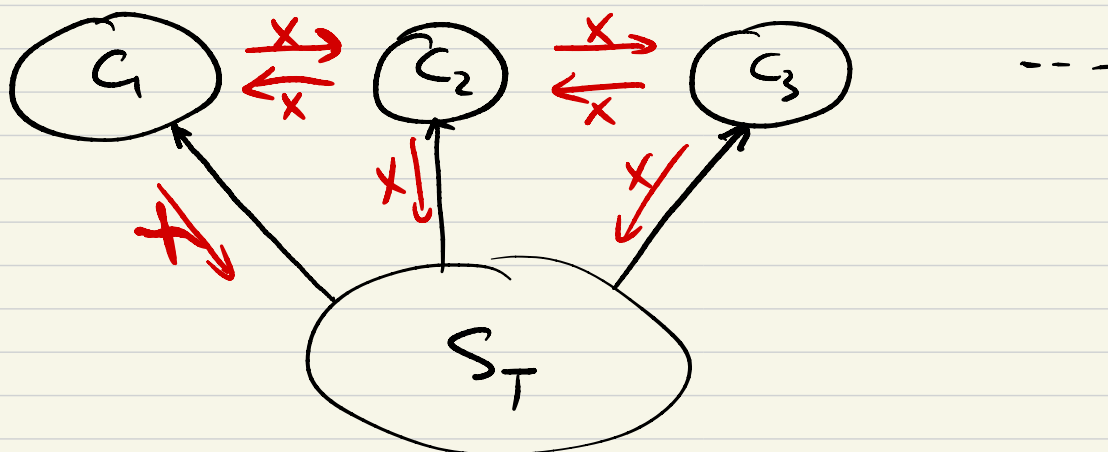
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$$S = S_T \cup S_R$$

↑
State Space

$$= S_T \cup \left(\bigcup_{i=1}^k C_i \right)$$

$k \leq \infty$
 disjoint ↑ closed & irreducible of recurrent states



Reordering states by the decomposition

for instance $S = S_T \cup C_1 \cup C_2$
(case $k=2$)

\tilde{P} : def.

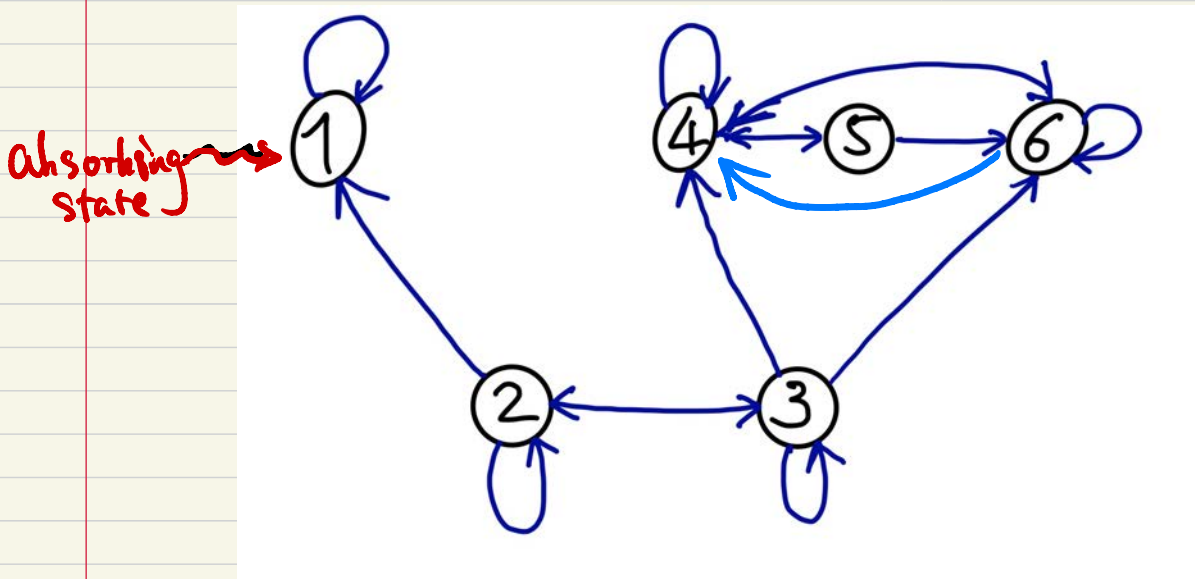
↑ Markov matrix

	C_1	C_2	S_T
C_1	P_1	O	O
C_2	O	P_2	O
S_T	$*$	$*$	Q

Canonical forms of the Markov matrix P .

An example: Find the canonical Markov matrix for

$$P = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 & 6 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{matrix} & \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ \frac{1}{4} & \frac{1}{2} & \frac{1}{4} & 0 & 0 & 0 \\ 0 & \frac{1}{5} & \frac{2}{5} & \frac{1}{5} & 0 & \frac{1}{5} \\ 0 & 0 & 0 & \frac{1}{6} & \frac{1}{3} & \frac{1}{2} \\ 0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 0 & 0 & \frac{1}{4} & 0 & \frac{3}{4} \end{bmatrix} \end{matrix}$$



1^o. $\{4, 5, 6\}$ is closed (\because no state in this set leads to a state outside it.)

$4 \rightarrow 5 \rightarrow 6 \rightarrow 4$; \therefore $\{4, 5, 6\}$ is irreducible
leads to

\therefore $\{4, 5, 6\}$ is an irreducible closed set

\therefore it is finite

$\therefore C_2 \stackrel{\text{def.}}{=} \{4, 5, 6\}$ is an irreducible closed set of recurrent state

2° $C_1 = \{1\}$: irreducible closed set of recurrent states.

3° $S_T = \{2, 3\}$

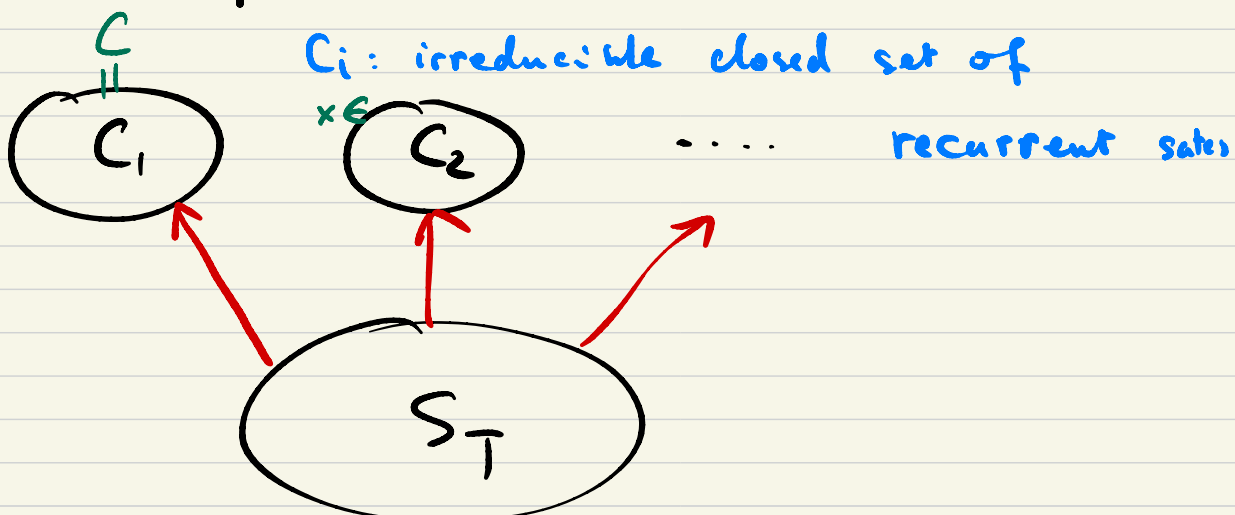
$$S = S_T \cup C_1 \cup C_2$$

$\begin{matrix} \parallel \\ \{2, 3\} \\ \parallel \\ e \quad f \end{matrix}$
 $\begin{matrix} \parallel \\ \{1\} \\ \parallel \\ a \end{matrix}$
 $\begin{matrix} \parallel \\ \{4, 5, 6\} \\ \parallel \\ b \quad c \quad d \end{matrix}$

Find the Canonical form

	C_1	C_2			S_T	
	$a=1$	$b=4$	$c=5$	$d=6$	$e=2$	$f=3$
C_1	$a=1$	1	0	0	0	0
C_2	$b=4$	0	$\frac{1}{6}$	$\frac{1}{3}$	$\frac{1}{2}$	0
	$c=5$	0	$\frac{1}{2}$	0	$\frac{1}{2}$	0
	$d=6$	0	$\frac{1}{4}$	0	$\frac{3}{4}$	0
S_T	$e=2$	$\frac{1}{4}$	0	0	0	$\frac{1}{2}$
	$f=3$	0	$\frac{1}{5}$	0	$\frac{1}{5}$	$\frac{2}{5}$

The last computational issue:



Q.: Compute for $C = \text{some } C_i$

irreducible
closed set
of recurrent states,

$$f_C(x) \stackrel{\text{def.}}{=} P_x(T_C < \infty), \quad x \in S$$

Recall: $T_C \stackrel{\text{def.}}{=} \min \{ n \geq 1 : X_n \in C \}$

is the 1st positive time
the chain hits C .

is the prob. that the chain from x
enters into C in finite time

Observe:

$$P_x(T_C < \infty)$$

$$f_C(x) =$$

1

if $x \in C = C_i$

0

if $x \in S_R$ is still
recurrent, but
 $x \notin C = C_i$

?

if $x \in S_T$ is
a transient.

different from

$$f_{xy} = P_x(T_y < \infty)$$

We need to determine:

$$f_C(x) = P_x(T_C < \infty), \quad x \in S_T$$

↑
irreducible closed set of recurrent states

Claim: $f_C(x) = P_x(1 \leq T_C < \infty)$

$x \in S_T$

$$\stackrel{\uparrow}{=} P_x(T_C=1) + P_x(2 \leq T_C < \infty)$$

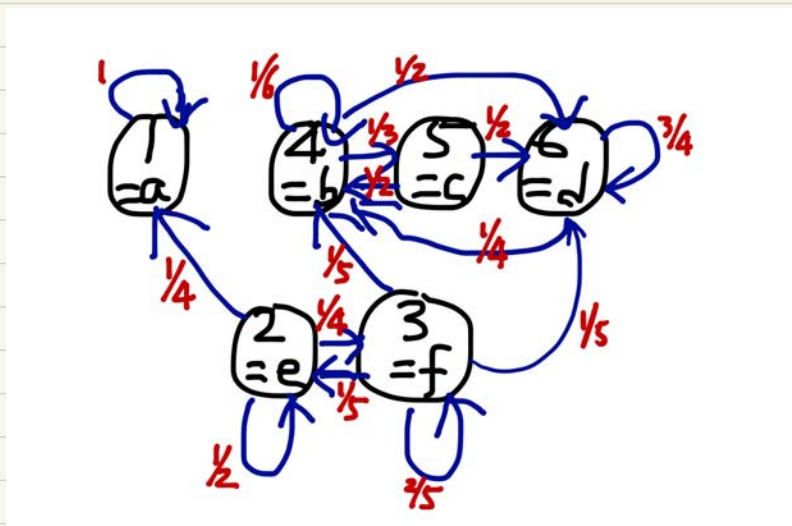
use the two-steps strategy

$$= \sum_{Y \in C} P(x, Y) + \sum_{Y \in S_T} P(x, Y) \underbrace{P_C(Y)}_{\parallel P_Y(1 \leq T_C < \infty)}$$

this is a linear system of equations for the unknowns $\{P_C(x)\}_{x \in S_T}$

Thm: If S_T (set of all transient states) is finite, then the linear system admits one and only one solution.

pf. omitted. #



$$C_1 = \{1\}$$

$$C_2 = \{4, 5, 6\}$$

$$S_T = \{2, 3\}$$

Compute

$$P_{C_2}(x), x \in S_T$$

i.e. look for

$$\text{two unknowns} \begin{cases} x \stackrel{\text{def.}}{=} P_{C_2}(2) = P_2(T_{C_2} < \infty) \\ y \stackrel{\text{def.}}{=} P_{C_2}(3) = P_3(T_{C_2} < \infty) \end{cases}$$

$$x \stackrel{2 \rightarrow C_2}{=} \underbrace{0 + 0 + 0}_{\text{one-step transition}} + \underbrace{\frac{1}{2} \cdot x + \frac{1}{4} \cdot y}_{\text{two-steps transition}}$$

Similarly,

$$y \stackrel{3 \rightarrow C_2}{=} \left(\frac{1}{5} + 0 + \frac{1}{5} \right) + \left(\frac{1}{5} \cdot x + \frac{2}{5} \cdot y \right)$$

$$\therefore \begin{cases} x = \frac{1}{2}x + \frac{1}{4}y \\ y = \frac{2}{5} + \left(\frac{1}{5}x + \frac{2}{5}y \right) \end{cases}$$

$$\Rightarrow \dots \Rightarrow x = \frac{2}{5}, \quad y = \frac{4}{5} \neq$$

Remarks:

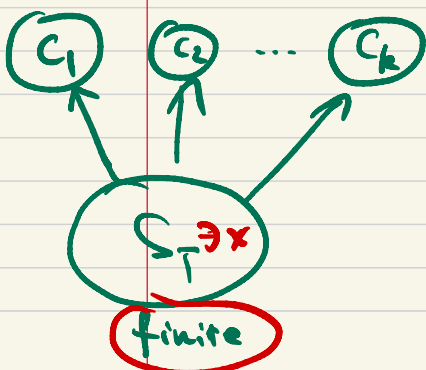
(i) If S_T is finite, then for any $x \in S_T$,

$$\sum_{i=1}^k P_{C_i}(x) = 1$$

where

$$S_R = \bigcup_{i=1}^{k < \infty \text{ or } = \infty} C_i$$

disjoint
irred + closed of recurrent



The chain from x hits some C_i in finite time "for sure"

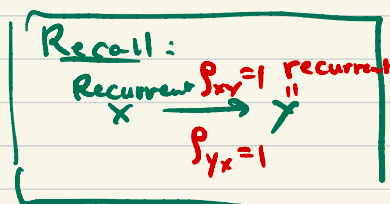
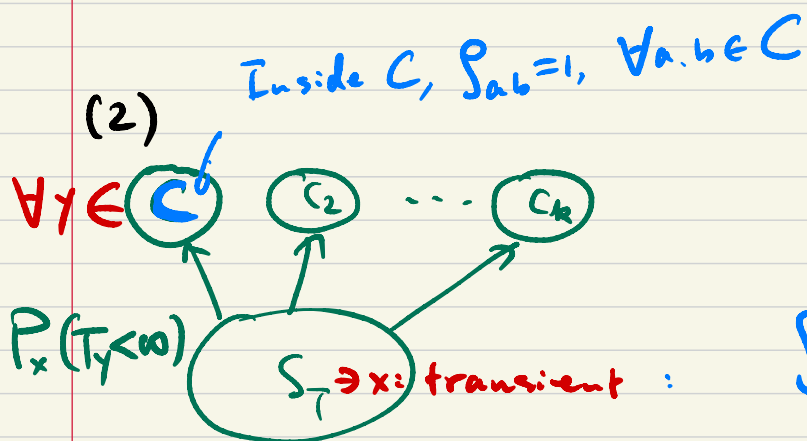
(i.e. with prob = 1)

Heuristic argument:

The total times the chain from x visit any transient state have to be finite for sure, and after then the chain must hit some C_i .

$$\begin{aligned}
 \text{LHS} &= \sum_{i=1}^k f_{C_i}(x) \\
 &= \sum_{i=1}^k P_x(1 \leq T_{C_i} < \infty) \\
 &= \dots = P_x(\underbrace{T_{S_P} < \infty}) = 1 \quad \#
 \end{aligned}$$

We only have finite transient states. #



$$f_C(x) = P_x(T_C < \infty)$$

Claim: $f_C(x) = P_{xy}, \forall y \in C$.

Pf. $\{1 \leq T_C < \infty\} \stackrel{=} {=} \{1 \leq T_y < \infty\} \quad \#$