

Answer to tutorial 1-4 exercises

Selected.

Tutorial 1

Q1: show $(a+b)^2 = a^2 + 2ab + b^2$ & $(-a)^2 = a^2$.

$$\begin{aligned} \text{Ans: } (a+b)^2 &= (a+b) \cdot (a+b) && \text{by def.} \\ &= (a+b) \cdot a + (a+b) \cdot b && \text{by distributivity.} \\ &= a \cdot a + b \cdot a + a \cdot b + b \cdot b && \text{as above.} \\ &= a^2 + b \cdot a + a \cdot b + b^2 && \text{by def.} \\ &= a^2 + a \cdot b + a \cdot b + b^2 && \text{by commutativity.} \\ &= a^2 + 1 \cdot (a \cdot b) + 1 \cdot (a \cdot b) + b^2 && 1 \text{ is multiplicative id.} \\ &= a^2 + (1+1) \cdot (a \cdot b) + b^2 && \text{by reverse of distributivity.} \\ &= a^2 + 2 \cdot (a \cdot b) + b^2 && \text{by def, } 1+1=2. \end{aligned}$$

$$\begin{aligned} (-a)^2 &= (-a) \cdot (-a) && \text{by def} \\ &= ((-1) \cdot a) \cdot ((-1) \cdot a) && \text{by lemma below.} \\ &= ((-1) \cdot a) \cdot (-1) \cdot a && \text{by associativity.} \\ &= ((-1) \cdot ((-1) \cdot a)) \cdot a && \text{by commutativity.} \\ &= (((-1) \cdot (-1)) \cdot a) \cdot a && \text{by associativity.} \\ &= (1 \cdot a) \cdot a && \text{by lemma.} \\ &= a \cdot a = a^2 && \text{by multiplicative id & def.} \end{aligned}$$

lemma: $-a = (-1) \cdot a$ & $(-1) \cdot (-1) = 1$.

subpf: by uniqueness of additive inverse, (i.e. $a+b=0$ and $a+c=0 \Rightarrow b=c$),
one just need to show $(-1) \cdot a + a = 0$.

$$(-1) \cdot a + a = (-1) \cdot a + 1 \cdot a = ((-1) + 1) \cdot a = 0 \cdot a$$

$$0 \cdot a + 0 \cdot a = (0+0) \cdot a = 0 \cdot a$$

$$0 \cdot a = 0 \cdot a + 0 \cdot a + (-1 \cdot 0 \cdot a) = 0 \cdot a + (-1 \cdot 0 \cdot a) = 0$$

$$\text{hence } (-1) \cdot a = -a. \quad \text{hence } (-1) \cdot (-1) = -(-1) = 1$$



Tutorial 2

Q2: show that $1 < 2 < 3 < \dots$ and $(n, n+1) \cap \mathbb{N} = \emptyset \quad \forall n \in \mathbb{N}$.

Ans: one need to formulate it rigorously to get a proof.

which is $P(i)$ proposition: $i < i+1 < i+2 < \dots$
i.e. $Q(j)$ is true where j starts from i

Proposition: if $j \geq i$, then $j+1 > i$.

We verify: $P(1)$:

We verify $Q(1)$: we have $1 \geq 1$,
and $1 > 0$.

hence adding to $2 := 1+1 > 1+0 = 1$
hence $Q(1)$ is true.

Let $j \in \mathbb{N}$. Suppose $j \geq 1$.

Suppose $Q(j)$ is true,

hence $j+1 > 1$ by $Q(j)$.

hence this satisfies requirement. $j+1 \geq i$ of $Q(j+1)$,

hence also $1 > 0$.

adding to $j+2 := (j+1)+1 > i+0 = i$.

hence $Q(j+1)$ is true.

by principle of math induction, $Q(j)$ is true on \mathbb{N} .

hence $P(1)$ is true.

Let $i \in \mathbb{N}$, Suppose $P(i)$ is true,

$i > 0$ we verify $P(i+1)$:

We verify $Q(0)$: obviously $0 < i$, hence $Q(0)$ is true.

Let $j \in \mathbb{N}$, suppose $j \geq i+1$, suppose $Q(j)$ is true,

then we verify $Q(j+1)$:

since $j \geq i+1$,

we have $j+1 \geq j \geq i+1$,

now $1 > 0$.

$j+2 := (j+1)+1 > i+0+0 = i+1$

hence $Q(j+1)$ is true.

by principle of math induction, $Q(j)$ is true on \mathbb{N} .

hence $P(j+1)$ is true

by principle of math induction, $P(i)$ is true on \mathbb{N} .

hence proven.

For second part,

Tutorial 1 Q2: (cont'd)

Ans: for second part, you are to show that there is no integer between n and $n+1$.

or you show that
if $j > n \Rightarrow j \geq n+1$.

prove by induction:

$P(0): j > 0$

then by previous prop, you know that j should be in the latter list of number $1 < 2 < 3 < \dots$,

which is meaning j is inside the range of objects which are can attain by adding 1 to the integer 1, i.e. an inductive object.

as the previous prop show, $j \geq 0+1=1$ obviously.

now if $P(i)$ is true,

we show $P(i+1)$:

if $j > i+1$, how $i \in \mathbb{N}$

hence $j > 0$

hence $j > 0+1 > 0$

hence j should be obtained as an inductive object

i.e. $j = j' + 1$ for some $j' \in \mathbb{N}$.

$j' + 1 > i + 1$

now we know $j' > i$

(if not then $j' \leq i$)

hence $j' + 1 \leq i + 1$, contradiction.

hence by $P(i): j' \geq i + 1$

hence ~~$j' \geq i + 1$~~

$j = j' + 1 \geq i + 1 + 1 = i + 2$

here proven by ~~the~~ MI.

now if $j \in (n, n+1) \cap \mathbb{N}$, no \mathbb{N}

then $j > n$ & $j < n+1$.

but by prop \Rightarrow ~~$j \geq n+1$~~ , contradiction.

hence \nexists such j .



Tutorial 1

Q4: $X \neq \emptyset, \quad \sup X = -\inf(-X)$

Ans: ~~$\forall a \in X, a \leq \sup X$~~
~~hence $-a \geq -\sup(X)$~~
~~but $a \in -X$, hence~~

$\forall b \in (-X), \exists a \in X \text{ s.t. } b = -a.$

hence $b = -a.$

$-b = a \in X.$

hence $-b \leq \sup X.$

since \sup is an upper bound

$b \geq -\sup X.$

$\forall b \in (-X).$

hence $-\sup X$ is a lower bound of set $(-X).$

Since $\inf(-X)$ is defined as ~~LEAST~~ GREATEST lower bound of set $(-X),$

hence as lower bound,

$\inf(-X) \geq -\sup(X).$

greatest lower bound

one of lower bound.

$\forall a \in X, \text{ since } X = [-(-X)]$
 $\exists b \in (-X) \text{ s.t. } a = -b.$

hence $a = -b$

$-a = b \in (-X).$

hence $-a \geq \inf(-X)$

since \inf is a lower bound

~~since~~

~~$a \geq -\inf(-X)$~~

$a \leq -\inf(-X).$

$\forall a \in X.$

hence $-\inf(-X)$ is an upper bound of $(X),$

Since $\sup(X)$ is defined as LEAST upper bound,

hence as upper bound,

$\sup X \leq -\inf(-X)$

least upper bound.

one of upper bound.

hence $-\sup X \geq \inf(-X).$

hence $-\sup X = \inf(-X)$

$\sup X = -\inf(-X)$



Tutorial 1

Q5: replace definition of limit of sequence with ε by $5\varepsilon^2$, N by \sqrt{N} , show they are equivalent.

Ans: the two statements are: (fix l).

(*) : $\forall \varepsilon > 0, \exists N \in \mathbb{N}$ s.t. $\forall n \in \mathbb{N}$ s.t. $n \geq N$,
we have $|x_n - l| < \varepsilon$.

(**) : $\forall \varepsilon > 0, \exists N \in \mathbb{N}$ s.t. $\forall n \in \mathbb{N}$ s.t. $n \geq \sqrt{N}$,
we have $|x_n - l| < 5\varepsilon^2$.

we show $(*) \Rightarrow (**)$:

the reverse side is obviously stronger.

Suppose (x_n) converge to l in definition of (*).
 $\forall \varepsilon > 0$, (this ε is for ε in definition (**).)

For $\varepsilon' = 5\varepsilon^2$,

by definition of (*),

$\exists N' \in \mathbb{N}$ s.t. $\forall m \in \mathbb{N}$ s.t. $m \geq N'$,
we have $|x_m - l| < \varepsilon'$.

take $N = (N')^2$.

now. $\forall n \in \mathbb{N}$ s.t. $n \geq \sqrt{N}$.

then $n \geq \sqrt{N} = N'$,

hence by previous deduction, $|x_n - l| < \varepsilon' = 5\varepsilon^2$.

hence we have (x_n) converge to l
in definition (**). □

Tutorial 2

Ex 2.3.2: $S_2 = \{x \in \mathbb{R} : x > 0\}$. lower bound? upper bound?
~~inf?~~ ~~sup?~~

Ans: ①: $\forall x \in S_2$,
by def, $x > 0$.
hence 0 ~~is~~ is a lower bound to S_2 .
so S_2 has a lower bound.

②: suppose S_2 has an upper bound M ,
then $\forall x \in S_2$, $M \geq x$.
hence $M \geq x > 0$.
hence M is positive.
take $M' = M + 1 > 0$.
hence $M' \in S_2$.
~~but~~ $M \geq M' \in S_2$ by definition,
but $M' = M + 1 > M$ obviously.
contradiction. hence S_2 has no upper bound.

③. we show 0 is the infimum of S_2 .
~~supp~~ by completeness, let $u = \inf S_2$. (since S_2 bounded below by 0).
hence $u \geq 0$.
greater lower bound one of lower bound

now suppose $u > 0$ (i.e. $u \neq 0$).
then consider $u' = \frac{u}{2} > 0$ hence $u' \in S_2$.
~~by~~ by def of inf (as lower bound), ~~but~~
 $u \leq u' \in S_2$.
but $u' = \frac{u}{2} < u$.
contradiction.
hence $u = 0$.

④. Since S_2 has no upper bound,
it cannot have supremum.
since supremum is an upper bound.



Tutorial 2

Ex 2.3.14: S bounded below.
 lower bound w of S is $\inf S$ iff $\forall \epsilon > 0, \exists t \in S$ s.t. $t < w + \epsilon$.

Ans: with w lower bound of S fixed,
 if w is the $\inf S$, then w is the greatest among all lower bound,
 i.e. if $u > w$, then u cannot be a lower bound of S ,
 hence $\forall \epsilon > 0, w + \epsilon > w$
 hence $w + \epsilon$ cannot be a lower bound of S , failing to bound $t \in S$.
 hence $\exists t \in S$ s.t. $t < w + \epsilon$.

Conversely, suppose for contradiction that w is not the $\inf S$.
 hence w is not the greatest among all lower bound,
 hence by completeness, let $w' = \inf S$,
 w' is greatest among all lower bound of S ,
 hence $w' > w$.
~~now let $\epsilon = w' - w > 0$,~~
 now let $\epsilon = w' - w > 0$,
 by assumption, $\exists t \in S$ s.t. $t < w + \epsilon = w + (w' - w) = w'$.
 hence w' fail to bound some $t \in S$.
 hence w' fail to be even a lower bound of S .
 hence contradiction.
 hence w is the \inf of S . \square

Ex 2.4B prove that $\sup_{x,y} h(x,y) = \sup_x (\sup_y h(x,y)) = \sup_y (\sup_x h(x,y))$.

Ans. we only show one side, $\sup_{x,y} = \sup_x (\sup_y)$.
 $\forall x \in X, \forall y \in Y, h(x,y) \leq \sup_y h(x,y)$ (fixing x here, vary over y).
 The RHS become a function of x only, denote it as $F(x)$,
 hence $h(x,y) \leq F(x) \forall x,y$.
 vary over x .
 We have $h(x,y) \leq \sup_x F(x)$.
 hence $\sup_x F(x) = \sup_x (\sup_y h(x,y))$ is an upper bound of all $h(x,y)$,
 hence $\sup_{x,y} h(x,y) \leq \sup_x (\sup_y h(x,y))$.

Tutorial 2 (cont'd)

Ex 2.4.12

Ans: Conversely, consider the set of all values, denoted in short hand $h(X, Y)$.

now if we fix $a \in X$, $\forall y \in Y$
then $h(a, y) \in h(X, Y)$.

collect these into set, written $h(a, Y) \subseteq h(X, Y)$. are sets
hence (sup of subset \leq sup of set).

$$\sup(h(a, Y)) \leq \sup(h(X, Y)) \quad \leftarrow \text{sup of set.}$$

but the LHS

is then a function of a , denoted $F(a)$,

$$\text{hence } F(a) \leq \sup h(X, Y). \quad \forall a \in X.$$

hence $\sup h(X, Y)$ is an upper bound to all such $F(a)$,

$$\text{hence } \sup h(X, Y) \geq \sup_a F(a).$$

$$= \sup_a (\sup_b h(a, b))$$

hence conclude the equation. \square

Ex. 2.5.8: show $\bigcap J_n = \emptyset$ ($J_n = (0, \frac{1}{n})$, $n \in \mathbb{N}_x$).

Ans: Suppose $a \in \bigcap J_n$,
then $a \in J_n \quad \forall n \in \mathbb{N}_x$

hence ~~sup~~ $a > 0$, hence $0 < a < \frac{1}{n} \quad \forall n \in \mathbb{N}_x$.

by Archimedean principle, $\exists m \in \mathbb{N}$ st. $\frac{1}{m} < a$, $m > 0$.

hence $a > \frac{1}{m}$, hence $a \notin J_m$.

Contradiction.

hence $\bigcap J_n = \emptyset$. \square

Tutorial 2

Ex 3.1.12 show $\lim_{n \rightarrow \infty} (\sqrt{n^2+1} - n) = 0$.

Ans: before doing pf, do rough work here.

want $|\sqrt{n^2+1} - n - 0| < \epsilon$.

$\sqrt{n^2+1} - n < \epsilon$.

$(n^2+1) - 2n\sqrt{n^2+1} + n^2 < \epsilon^2$.

$2n^2+1 - 2n\sqrt{n^2+1} < \epsilon^2$.

$(2n^2+1) - \epsilon^2 < 2n\sqrt{n^2+1}$.

$(2n^2+1)^2 - 2\epsilon^2(2n^2+1) + \epsilon^4 < 4n^2(n^2+1)$.

$4n^4 + 4n^2 + 1 - 2\epsilon^2(2n^2+1) + \epsilon^4 < 4n^4 + 4n^2$.

$1 - 2\epsilon^2(2n^2+1) + \epsilon^4 < 0$.

$\epsilon^4 + 1 < 2\epsilon^2(2n^2+1)$.

Since n expected to be large, assume RHS > 4
take ϵ small \rightarrow obviously this is true.

so need $2\epsilon^2(2n^2) > 4$

$n > \frac{1}{\epsilon}$ is a good choice.

also mind the case if ϵ is not small. (say $\epsilon \geq 1$).

then we need $2\epsilon^2(2n^2) > \epsilon^4$

$n > \epsilon$.

proceed with it

Roughwork area

$\forall \epsilon > 0$, take $N = \max \{ \lceil \frac{1}{\epsilon} \rceil, \lceil \epsilon \rceil \} + 1$

$\forall n \geq N$, divide into 2 cases:

(no need take ceiling function just we archimedean principle to take integer $\rightarrow \frac{1}{\epsilon}, \epsilon$.)

case 1: $\epsilon \geq 1$, then $\epsilon^4 + 1 < \epsilon^2 \cdot n^2 + \epsilon^2$.

$\leq 2\epsilon^2 \cdot 2n^2 + 2\epsilon^2 = 2\epsilon^2(2n^2+1)$

$\epsilon^4 + 1 < 2\epsilon^2(2n^2+1)$.

$1 - 2\epsilon^2(2n^2+1) + \epsilon^4 < 0$.

$4n^4 + 4n^2 + 1 - 2\epsilon^2(2n^2+1) + \epsilon^4 < 4n^4 + 4n^2$.

$(2n^2+1)^2 - 2\epsilon^2(2n^2+1) + \epsilon^4 < 4n^2(n^2+1)$.

$((2n^2+1) - \epsilon^2)^2 < (2n\sqrt{n^2+1})^2$.

since $2n^2+1 > \epsilon^2$.

$2n^2+1 - \epsilon^2 < 2n\sqrt{n^2+1}$.

$(n^2+1) - 2n\sqrt{n^2+1} + n^2 < \epsilon^2$.

$(\sqrt{n^2+1} - n)^2 < \epsilon^2$

$\rightarrow \sqrt{n^2+1} - n < \epsilon$
 $|\sqrt{n^2+1} - n - 0| < \epsilon$.

since $\frac{n^2+1}{n^2+1} > n^2$
 $\frac{n^2+1}{n^2+1} > n$.

proceed back ward of above rough work.

(Cont'd)

Tutorial 2

Ex 3.1.12

Case 2: $\epsilon < 1, n > \frac{1}{\epsilon}$

Ans:

hence $2\epsilon^2(2n^2+1) > 4+2\epsilon^2 > 1+2\epsilon^4 > 1+\epsilon^4$

hence $\epsilon^4+1 < 2\epsilon^2(2n^2+1)$

$1-2\epsilon^2(2n^2+1)+\epsilon^4 < 0$

$(4n^4+4n^2+1)-2\epsilon^2(2n^2+1)+\epsilon^4 < 4n^4+4n^2$

$(2n^2+1)^2-2\epsilon^2(2n^2+1)+\epsilon^4 < 4n^2(n^2+1)$

$((2n^2+1)-\epsilon^2)^2 < (2n\sqrt{n^2+1})^2$

$2n^2+1-\epsilon^2 < 2n\sqrt{n^2+1}$

since $1 > \epsilon^2$

$(n^2+1)-2n\sqrt{n^2+1}+n^2 < \epsilon^2$

$(\sqrt{n^2+1}-n)^2 < \epsilon^2$

$\sqrt{n^2+1}-n < \epsilon$

since $n^2+1 > n^2$
 $\sqrt{n^2+1} > n$

$|\sqrt{n^2+1}-n-0| < \epsilon$

hence true in both case. □

~~Ex 3.1.16: show $\lim_{n \rightarrow \infty} \frac{n^2}{n!} = 0$~~

~~Ans: Do roughwork first:~~

~~wait $\frac{n^2}{n!} = 0 < \epsilon$ Roughwork~~
 ~~$\frac{n^2}{n!} < \epsilon$~~
 ~~$\frac{n}{(n-1)!} < \epsilon$ wait $n > 1$~~
~~hence $\frac{1}{(n-1)!} < \epsilon$ since $2(n-1) > n$~~
~~stronger $\frac{1}{2(n-1)!} < \epsilon$ since $2 > \frac{n}{n-1}$~~

Tutorial 3

HW 2.4. show $\sup(A+B) = \sup A + \sup B$ if either side exists.

Ans: ~~let~~ let $a = \sup A$, $b = \sup B$, $e := \sup(A+B)$.

$\forall x \in (A+B)$, $\exists c \in A, d \in B$ s.t. $x = c+d$.

hence $x = c+d$ ~~$\leq a+b$~~

$\leq a+d$ since $a \geq c \forall c \in A$.

$\leq a+b$ since $b \geq d \forall d \in B$.

hence $a+b$ constant is an upper bound to any $x \in (A+B)$, hence to whole set $(A+B)$,

hence $a+b \geq \sup(A+B) = e$

one of upper bound. ~~least~~ upper bound.

Conversely,

$\forall c \in A, d \in B$,

note that $(A+B) = \bigcup_{c \in A} (c+B)$. cosets.

hence fix $c \in A$, $c+B \subseteq A+B$.

hence $\sup(c+B) \leq \sup(A+B) = e$.

$c + \sup B \leq e$.

$c + b \leq e \quad \forall c \in A$.

hence $c \leq e - b$

$e - b$ is hence an upper bound to any $c \in A$ constant. hence.

$e - b \geq \sup A = a$.
one of upper bound. least upper bound.

hence $e \geq a+b$.

hence done. \square

Ex 3.2.5 show they are not convergent.

(2^n) , $((-1)^n n^2)$.

Ans: (a): $\forall l \in \mathbb{R}$, (we show $\exists \epsilon > 0$ s.t. $\forall N \in \mathbb{N}$, $\exists n \geq N$ s.t. $|x_n - l| \geq \epsilon$)

take $\epsilon = 1$, $\forall N \in \mathbb{N}$,

take by Archimedean prop. N s.t. $\forall n \geq N$, ~~$|2^n - l| \geq 1$~~ ~~$|(-1)^n n^2 - l| \geq 1$~~

hence ~~2^n~~ $N > |l| + 2$.

(Cont'd)

tutorial 3

Ex 3.2.5

Ans: (a): with $n \geq N$, $n > |l| + 1$.

$$\text{hence } 2^n > 2^{|l|+1} = 2 \cdot 2^{|l|} = 2^{|l|} + 2^{|l|} \\ \geq 2^{|l|} + 1 \geq |l| + 1 \geq l + 1.$$

$$\text{hence } |2^n - l| \geq 1.$$

hence is divergent.

(b): $\forall l \in \mathbb{R}$,

take $\epsilon = 1$, $\forall N \in \mathbb{N}$,

take $n \geq N$ and $n \geq |l| + 2$ and n is even.

hence ~~$(-1)^n n^2 \geq (l+1)^n$~~

$$(-1)^n n^2 = n^2$$

$$\geq (|l|+2)^2 = |l|^2 + 4|l| + 4 > |l| + 1.$$

hence ~~$(-1)^n n^2 \geq (l+1)$~~

$$\geq l + 1$$

$$|(-1)^n n^2 - (l)| \geq 1.$$

hence divergent.

Ex 3.2.11

determine limit of $(3\sqrt[n]{n})^{\frac{1}{2n}}$, $(n+1)^{\frac{1}{n(n+1)}}$

Ans: (a): do rough work first.

$$\text{want } |(3\sqrt[n]{n})^{\frac{1}{2n}} - 1| < \epsilon.$$

$$(3\sqrt[n]{n})^{\frac{1}{2n}} < 1 + \epsilon.$$

$$3\sqrt[n]{n} < (1 + \epsilon)^{2n}.$$

show
Steyer
statement.

$$3\sqrt[n]{n} < 1 + \sqrt[n]{n} \epsilon + \frac{\sqrt[n]{n}(\sqrt[n]{n}-1)}{2} \epsilon^2$$

by
binomial series

$$\text{take } n \text{ s.t. } \frac{\sqrt[n]{n}-1}{2} \epsilon^2 > 3.$$

if $\epsilon \geq 1$,
then need

$$\frac{\sqrt[n]{n}-1}{2} > 3.$$

$$n > 49.$$

$\forall \epsilon > 0$, take N s.t. $N > \left(\frac{6}{\epsilon^2} + 1\right)^2$, hence $\frac{\sqrt[n]{n}-1}{2} \epsilon^2 > 3$.

and $N > 49$.

$\forall n \geq N$.

(Cont'd)

Interval 3

Ex 3.2.11

Ans. (a):

Case 1: $\epsilon \geq 1$,

hence with $N > 49$, $n \geq N$.

hence $\frac{\sqrt{2n}-1}{2} > 3$.

hence $3\sqrt{n} < \frac{\sqrt{2n}(\sqrt{2n}-1)}{2} \epsilon^2 < (1+\sqrt{2n} \epsilon) + \frac{\sqrt{2n}(\sqrt{2n}-1)}{2} \epsilon^2 < (1+\epsilon)^{\sqrt{2n}}$
 $3\sqrt{n} < (1+\epsilon)^{\frac{\sqrt{2n}}{2}}$

$(3\sqrt{n})^{\frac{1}{\sqrt{2n}}} < \epsilon + 1$

$| (3\sqrt{n})^{\frac{1}{\sqrt{2n}}} - 1 | < \epsilon$. done.

Case 2: $\epsilon < 1$.

$n \geq N > \left(\frac{6}{\epsilon^2} + 1\right)^2$ hence $\frac{\sqrt{2n}-1}{2} \epsilon^2 > 3$.

hence $3\sqrt{n} < \frac{\sqrt{2n}(\sqrt{2n}-1)}{2} \epsilon^2 < 1 + \sqrt{2n} \epsilon + \frac{\sqrt{2n}(\sqrt{2n}-1)}{2} \epsilon^2 < (1+\epsilon)^{\sqrt{2n}}$
 $3\sqrt{n} < (1+\epsilon)^{\frac{\sqrt{2n}}{2}} < (1+\epsilon)^{\sqrt{2n}}$

proceed as above.

(b):

$(n+1)^{\frac{1}{g(n+1)}} = e^{\frac{\log(n+1)}{g(n+1)}} = 1$ □

hence $\forall \epsilon > 0$, take $N=1$,

then $\forall n \geq N$,

$| (n+1)^{\frac{1}{g(n+1)}} - 1 | = | 1 - 1 | = 0 < \epsilon$.

here done. □

Tutorial 4

Ex 3.3.2 $x_1 > 1, x_{n+1} = 2 - \frac{1}{x_n} \quad \forall n \in \mathbb{N}$

show it is bounded & monotone & limit.

Ans.

Suppose it is bounded & monotone,
by monotone convergence thm, \exists limit l ,

\rightarrow taking limit on both side

$$x_{n+1} = 2 - \frac{1}{x_n}$$

$$l = 2 - \frac{1}{l}$$

$$l^2 - 2l + 1 = 0$$

$$l = 1$$

(Justification:

$X^+ := (x_{n+1})$ as tail of sequence $X = (x_n)$.

X^+ & X has same limit.

and limit of $X^+ = (x_{n+1}) = (2 - \frac{1}{x_n})$

$$= \lim(2) - \lim \frac{1}{x_n}$$

$$= 2 - \frac{1}{\lim x_n}$$

by limit thm.

here we have $\lim X^+ = 2 - \frac{1}{\lim X}$

assuming $\lim X > 0$.

$$\lim X$$

$$l = 2 - \frac{1}{l}$$

show bounded:

inductively $x_i \in (1, 2)$.

$P(1): x_1 > 1$

$P(i)$ true

$\rightarrow P(i+1): x_{i+1} = 2 - \frac{1}{x_i}$
 $> 2 - 1 = 1$

and $x_{i+1} = 2 - \frac{1}{x_i} < 2$ clearly.

show monotone:

inductively $P(1)$:

$$x_2 - x_1 = 2 - \frac{1}{x_1} - x_1 = \frac{2x_1 - 1 - x_1^2}{x_1}$$

$$= -\frac{(x_1 - 1)^2}{x_1} < 0$$

here $x_2 < x_1$

$P(i)$ true \rightarrow

$P(i+1)$:

$$x_{i+1} - x_i$$

$$= 2 - \frac{1}{x_i} - x_i$$

$$= -\frac{1}{x_i} (x_i^2 - 2x_i + 1) = -\frac{(x_i - 1)^2}{x_i} < 0$$

here $x_{i+1} < x_i$

hence decreasing.

