

CENG 3420

Computer Organization & Design



Lecture 07: Floating Numbers

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Scientific notation: 6.6254×10^{-27}

- A normalized number of certain accuracy (e.g. 6.6254 is called the **mantissa**)
- Scale factors to determine the position of the decimal point (e.g. 10^{-27} indicates position of decimal point and is called the exponent; the **base** is implied)
- **Sign** bit

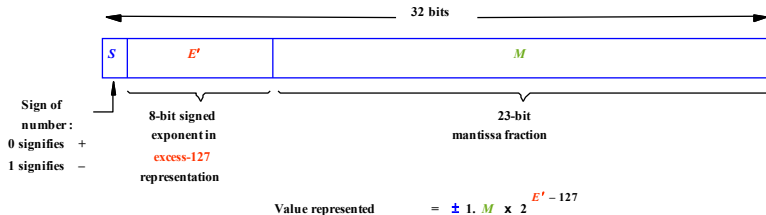


- Floating Point Numbers can have multiple forms, e.g.

$$\begin{aligned}0.232 \times 10^4 &= 2.32 \times 10^3 \\&= 23.2 \times 10^2 \\&= 2320. \times 10^0 \\&= 232000. \times 10^{-2}\end{aligned}$$

- It is desirable for each number to have a unique representation => Normalized Form
- We normalize Mantissa's in the Range $[1..R)$, where R is the Base, e.g.:
 - $[1..2)$ for BINARY
 - $[1..10)$ for DECIMAL

32-bit, float in C / C++ / Java



(a) Single precision



$$\text{Value represented} = +1.001010 \dots 0 \times 2^{-87}$$

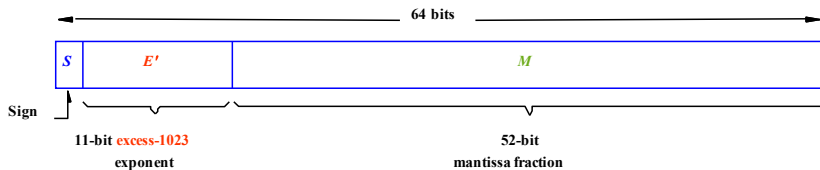
(b) Example of a single-precision number

00101000 → 40

$$40 - 127 = -87$$



64-bit, float in C / C++ / Java



$$\text{Value represented} = \pm 1. M \times 2^{E' - 1023}$$

(c) Double precision



Question:

What is the IEEE single precision number $40C0\ 0000_{16}$ in decimal?



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- Binary: $0100\ 0000\ 1100\ 0000\ 0000\ 0000\ 0000\ 0000$
- Sign: $+$
- Exponent: $129 - 127 = +2$
- Mantissa: $1.100\ 0000\ \dots_2 \rightarrow 1.5_{10} \times 2^{+2}$
- $\rightarrow +110.0000\ \dots_2$
- Decimal Answer = $+6.0_{10}$



Question:

What is -0.5_{10} in IEEE single precision binary floating point format?



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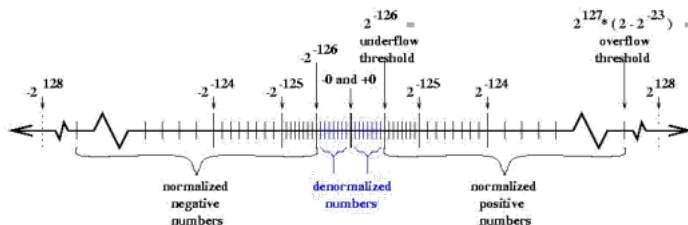
- Binary: $1.0... \times 2^{-1}$ (in binary)
- Exponent: $127 + (-1) = 01111110$
- Sign bit: 1
- Mantissa: 1.000 0000 0000 0000 0000 0000
- Binary representation: 1011 1111 0000 0000 0000 0000 0000 0000

- Normalized $\pm 1.d\dots d \times 2^{\text{exp}}$
- **Denormalized** $\pm 0.d\dots d \times 2^{\text{min_exp}}$ \rightarrow to represent *near-zero* numbers
e.g. $+ 0.0000\dots 0000001 \times 2^{-126}$ for Single Precision

Format	# bits	# significant bits	macheps	# exponent bits	exponent range
Single	32	1+23	2^{-24} ($\sim 10^{-7}$)	8	$2^{-126} - 2^{+127}$ ($\sim 10^{\pm 38}$)
Double	64	1+52	2^{-53} ($\sim 10^{-16}$)	11	$2^{-1022} - 2^{+1023}$ ($\sim 10^{\pm 308}$)
Double Extended	≥ 80	≥ 64	$\leq 2^{-64}$ ($\sim 10^{-19}$)	≥ 15	$2^{-16382} - 2^{+16383}$ ($\sim 10^{\pm 4932}$)

(Double Extended is 80 bits on all Intel machines)
macheps = Machine Epsilon = $2 - (\# \text{ significand bits})$

$$\epsilon_{\text{mach}}$$





Exponents of all 0's and all 1's have special meaning

- $E=0, M=0$ represents 0 (sign bit still used so there is ± 0)
- $E=0, M \neq 0$ is a denormalized number $\pm 0.M \times 2^{-127}$ (smaller than the smallest normalized number)
- $E=\text{All } 1\text{'s}, M=0$ represents $\pm \text{Infinity}$, depending on Sign
- $E=\text{All } 1\text{'s}, M \neq 0$ represents NaN



+, -, x, /, sqrt, remainder, as well as conversion to and from integer are correctly rounded

- As if computed with infinite precision and then rounded
- Transcendental functions (that cannot be computed in a finite number of steps e.g., sine, cosine, logarithmic, , e, etc.) may not be correctly rounded

Exceptions and Status Flags

- Invalid Operation, Overflow, Division by zero, Underflow, Inexact

Floating point numbers can be treated as “integer bit-patterns” for comparisons

- If Exponent is all zeroes, it represents a denormalized, very small and near (or equal to) zero number
- If Exponent is all ones, it represents a very large number and is considered infinity (see next slide.)

Dual Zeroes: +0 (0x00000000) and -0 (0x80000000): they are treated as the same



Infinity is like the mathematical one

- $\text{Finite} / \text{Infinity} \rightarrow 0$
- $\text{Infinity} \times \text{Infinity} \rightarrow \text{Infinity}$
- $\text{Non-zero} / 0 \rightarrow \text{Infinity}$
- $\text{Infinity}^{\{\text{Finite or Infinity}\}} \rightarrow \text{Infinity}$

NaN (Not-a-Number) is produced whenever a limiting value cannot be determined:

- $\text{Infinity} - \text{Infinity} \rightarrow \text{NaN}$
- $\text{Infinity} / \text{Infinity} \rightarrow \text{NaN}$
- $0 / 0 \rightarrow \text{NaN}$
- $\text{Infinity} \times 0 \rightarrow \text{NaN}$
- If x is a NaN, $x \neq x$
- Many systems just store the result quietly as a NaN (all 1's in exponent), some systems will signal or raise an exception



- E.g. Find 1st root of a quadratic equation
 - $r = (-b + \sqrt{b*b - 4*a*c}) / (2*a)$

Sparc processor, Solaris, gcc 3.3 (ANSI C),

Expected Answer	0.00023025562642476431
double	0.00023025562638524986
float	0.00024670246057212353

- Problem is that if c is near zero,

$$\sqrt{b*b - 4*a*c} \approx b$$

- Rule of thumb: use the highest precision which does not give up too much speed



- $(a - b)$ is inaccurate when $a \approx b$
- Decimal Examples
 - Using 2 significant digits to compute mean of 5.1 and 5.2 using the formula $(a+b) / 2$:
 $a + b = 10$ (with 2 sig. digits, 10.3 can only be stored as 10)
 $10 / 2 = 5.0$ (the computed mean is less than both numbers!!!)
 - Using 8 significant digits to compute sum of three numbers:
 $(11111113 + (-11111111)) + 7.5111111 = 9.5111111$
 $11111113 + ((-11111111) + 7.5111111) = 10.000000$
- Catastrophic cancellation occurs when

$$\left| \frac{[\text{round}(x) \bullet \text{round}(y)] - \text{round}(x \bullet y)}{\text{round}(x \bullet y)} \right| \gg \epsilon_{mach}$$