

Testing for Buffer Overflows by Length-Driven Symbolic Execution

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The Chinese University of Hong Kong*

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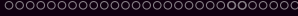
Outline

- 1 Introduction
 - Buffer Overflow
 - Symbolic Execution
- 2 Related work
 - Static Analysis
 - Dynamic Analysis
- 3 Our Approach
 - Motivation
 - Improvement
 - Key Idea
- 4 Q & A



The Essentials

■ The Definition



Two Kinds of Buffer Overflow

■ Stack Overflow



Two Kinds of Buffer Overflow

- Stack Overflow
- Heap Corruption



The power of Buffer Overflow

- An example...



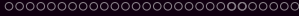
Fighting with Buffer Overflow

- Choice of programming language



Fighting with Buffer Overflow

- Choice of programming language
- Use of safe libraries



Fighting with Buffer Overflow

- Choice of programming language
- Use of safe libraries
- Buffer overflow protection(Dynamic and Static) example...



Fighting with Buffer Overflow

- Choice of programming language
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- Buffer overflow protection(Dynamic and Static) example...
- Pointer protection



Fighting with Buffer Overflow

- Choice of programming language
- Use of safe libraries
- Buffer overflow protection(Dynamic and Static) example...
- Pointer protection
- Address space layout randomization



What We want to do...

- We focus on...



What's Symbolic Execution

■ What's Symbolic Execution



What's Symbolic Execution

- What's Symbolic Execution
- Advantages



What's Symbolic Execution

- What's Symbolic Execution
- Advantages
- Disadvantages



An example

```
typedef struct cell {
  int v;
  struct cell *next;
} cell;
```

```
int f(int v) {
  return 2*v + 1;
}
```

```
int testme(cell *p, int x) {
  if (x > 0)
    if (p != NULL)
      if (f(x) == p->v)
        if (p->next == p)
          abort();
  return 0;
}
```

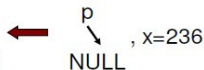
Concrete
Execution

Symbolic
Execution

concrete
state

symbolic
state

constraints



$p=p_0, X=X_0$



An example

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```

p
 ↘
 NULL

, x=236

p=p₀, x=x₀

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Concrete
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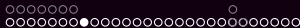
symbolic
state

constraints

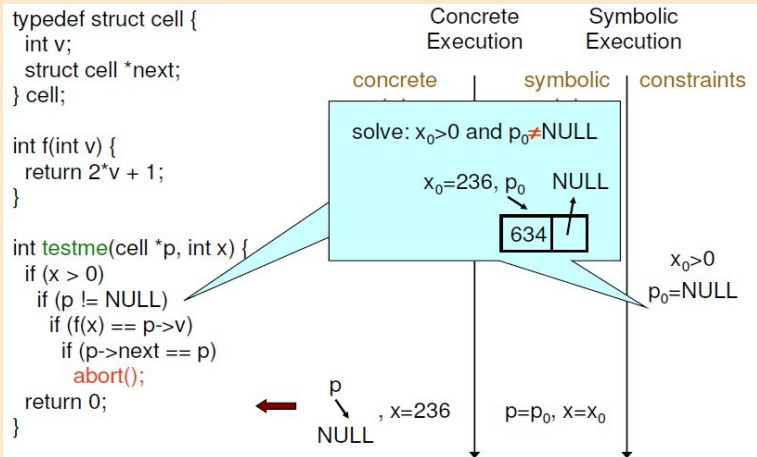
p
 \swarrow
 NULL, $x=236$

$p=p_0, x=x_0$

$x_0 > 0$
 $!(p_0 != \text{NULL})$



An example

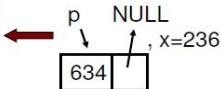


An example

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Concrete
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$p=p_0, x=x_0,$
 $p \rightarrow v = v_0,$
 $p \rightarrow next = n_0$

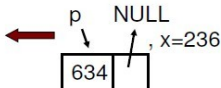
$x_0 > 0$

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Concrete
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$p=p_0, x=x_0,$
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$x_0 > 0$
 $p_0 \neq \text{NULL}$

An example

<pre> typedef struct cell { int v; struct cell *next; } cell; int f(int v) { return 2*v + 1; } int testme(cell *p, int x) { if (x > 0) if (p != NULL) if (f(x) == p->v) if (p->next == p) abort(); return 0; } </pre>	<p>Concrete Execution</p> <p>concrete state</p>	<p>Symbolic Execution</p> <p>symbolic state</p> <p>$p = p_0, x = x_0,$ $p \rightarrow v = v_0,$ $p \rightarrow next = n_0$</p>	<p>constraints</p> <p>$x_0 > 0$ $p_0 \neq NULL$ $2x_0 + 1 \neq v_0$</p>
--	---	---	---

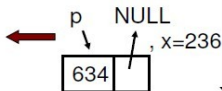


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$x_0 > 0$
 $p_0 \neq \text{NULL}$
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Concrete Execution

Symbolic Execution

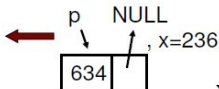
concrete

symbolic

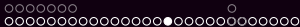
constraints

solve: $x_0 > 0$ and $p_0 \neq \text{NULL}$
and $2x_0 + 1 = v_0$

$x_0 > 0$
 $p_0 \neq \text{NULL}$
 $2x_0 + 1 \neq v_0$



$p=p_0, x=x_0,$
 $p \rightarrow v = v_0,$
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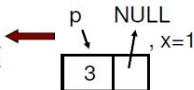


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Concrete Execution

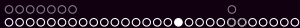
Symbolic Execution

concrete state

symbolic state

constraints

$p=p_0, x=x_0,$
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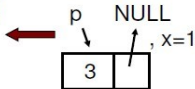


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concrete state

Symbolic Execution

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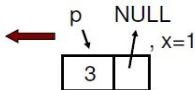
$x_0 > 0$

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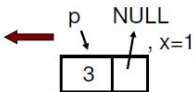
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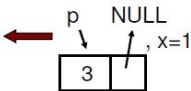
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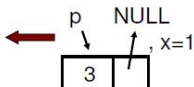
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Concrete
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concrete
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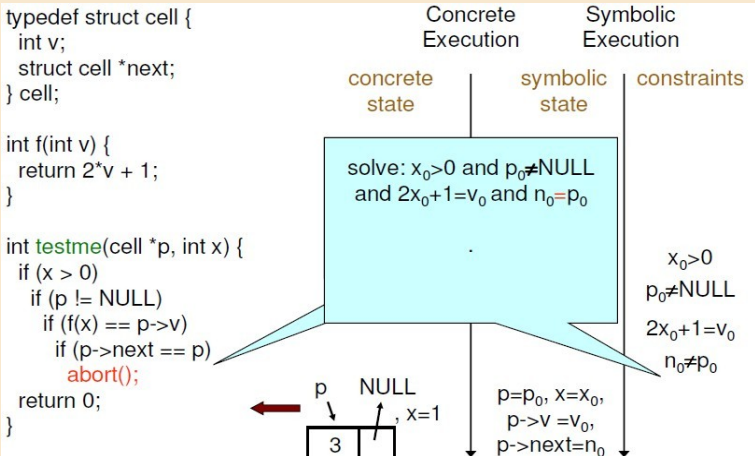
symbolic
state

constraints

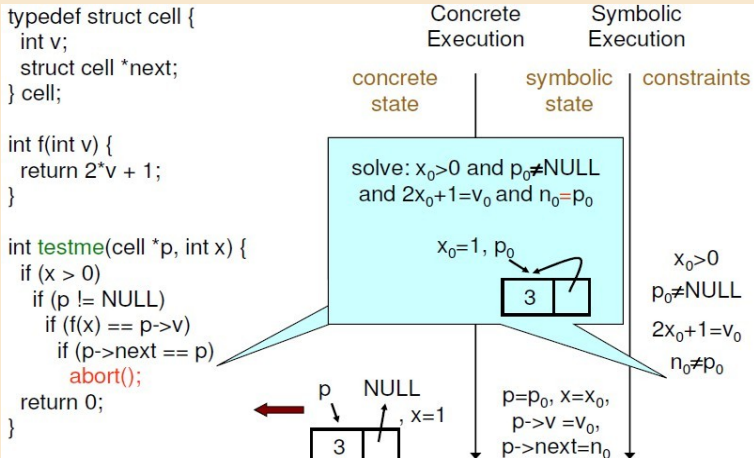
$p=p_0, x=x_0,$
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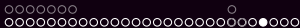
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An example



An example



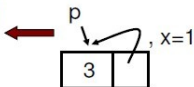


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Concrete
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Symbolic
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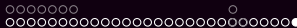
concrete
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$x_0 > 0$

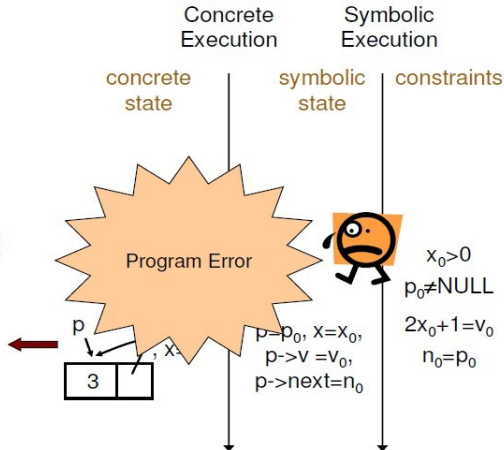


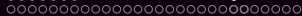
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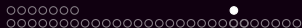
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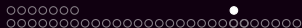
Static Analysis

■ Previous Work



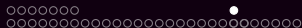
Static Analysis

- Previous Work
 - LCLint (USENIX Security 01)



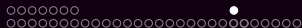
Static Analysis

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 - CSSV (PLDI 03)



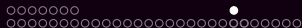
Static Analysis

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 - ESC/Java (PLDI 02)



Static Analysis

- Previous Work
 - LCLint (USENIX Security 01)
 - CSSV (PLDI 03)
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 - ARCHER (FSE 03)



Static Analysis

- Previous Work
 - LCLint (USENIX Security 01)
 - CSSV (PLDI 03)
 - ESC/Java (PLDI 02)
 - ARCHER (FSE 03)
- Drawbacks

Runtime Check

■ Previous Work

Runtime Check

■ Previous Work

- Purify (Winter USENIX Conf 1992)

Runtime Check

■ Previous Work

- Purify (Winter USENIX Conf 1992)
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Runtime Check

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Runtime Check

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■ Drawbacks

Symbolic Execution

■ Previous Work

Symbolic Execution

- Previous Work
 - DART(Random Testing) (PLDI 05)

Symbolic Execution

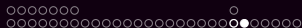
■ Previous Work

- DART(Random Testing) (PLDI 05)
- CUTE (FSE 05)

Symbolic Execution

■ Previous Work

- DART(Random Testing) (PLDI 05)
- CUTE (FSE 05)
- KLEE (OSDI 08)



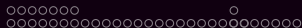
Symbolic Execution

■ Previous Work

- DART(Random Testing) (PLDI 05)
- CUTE (FSE 05)
- KLEE (OSDI 08)
- Length Abstraction (ISSTA 08)

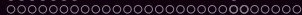
Symbolic Execution

- Previous Work
 - DART(Random Testing) (PLDI 05)
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 - KLEE (OSDI 08)
 - Length Abstraction (ISSTA 08)
- Drawbacks



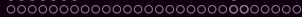
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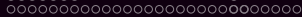
Improvement

- Compared with dynamic analysis



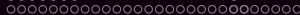
Improvement

- Compared with dynamic analysis
 - symbolic execution



Improvement

- Compared with dynamic analysis
 - symbolic execution
 - runtime check
- Compared with static analysis



Improvement

- Compared with dynamic analysis
 - symbolic execution
 - runtime check
- Compared with static analysis
- Compared with Length Abstraction (Unsure -, -)

Key Idea (by example)

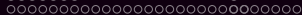
■ Sendmail 8.12.7 (daemon.c)

```

/*      char buf[50];
        char *p=&buf[0];
        char *s;
        char ibuf[50];
        int i, nleft;      */

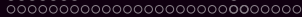
nleft=sizeof(ibuf)-1;
while((i=read(s,p,nleft))>0)
{
    p+=i;
    nleft-=i;
    *p='\0';
    if(strchr(ibuf, '\n') != NULL || nleft <= 0)
        break;
}
close(s);
if(i<0||p==&ibuf[0])
    goto noident;
if(*--p=='\n' && *--p == '\r')
    p--;
*++p = '\n';
noident:
return 0;

```

Q & A

■ Q & A



Q & A

- Q & A
- Thanks