

CSCI3160: Special Exercise Set 12

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Problem 1 (Textbook Exercise 35.3-1). Consider $\mathcal{S} = \{\text{arid, dash, drain, heard, lost, nose, shun, slate, snare, thread}\}$. Treat each word in \mathcal{S} as a set of letters. Run the set-cover algorithm discussed in the lecture and describe its output.

Problem 2. Recall that our set-cover algorithm in each iteration picks a set with the largest *benefit*. Prove: if we lay out the sets in the order they are picked, their benefits are non-ascending.

Problem 3*. Give an example to show that the approximation ratio of our set-cover algorithm cannot be bounded by 2.

Problem 4. Consider the set \mathcal{S} given in Problem 1. This time, we want to treat it as an input to the hitting set problem. Re-formulate the problem as a set-cover problem.

Problem 5. Let M be an $n \times m$ matrix where each cell is either 0 or 1. It is guaranteed that every row of M has at least one 1. A set S of columns is a *column cover* if every row of M has a 1 in at least one column of S . If OPT is the minimize size of all column covers, describe a $\text{poly}(n, m)$ -time algorithm (i.e., polynomial in n and m) that finds a column cover of size $O(\text{OPT} \cdot \log n)$.