

Church–Turing Thesis

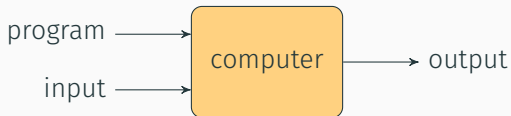
CSCI 3130 Formal Languages and Automata Theory

Siu On CHAN

Fall 2020

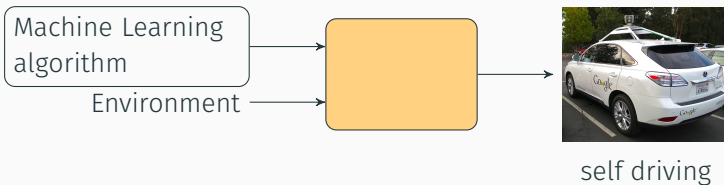
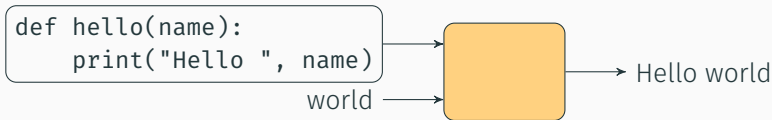
Chinese University of Hong Kong

What is a computer?

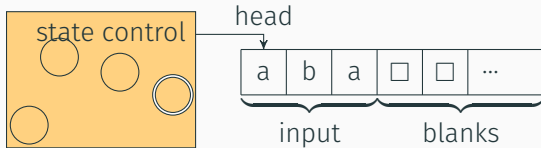


A **computer** is a machine that manipulates data according to a list of instructions

What is a computer?



Turing machines



Can both **read** from and **write** to the tape

Head can move both **left** and **right**

Unlimited tape space

Has two special states **accept** and **reject**

Example

$$L_1 = \{w#w \mid w \in \{a, b\}^*\}$$

Strategy:

| | |
|------------------------------------|--------------------------|
| Read and remember the first symbol | <u>a</u> bbaa#abbaa |
| Cross it off | x bbaa#abbaa |
| Read the first symbol past # | xbbaa# <u>a</u> bbaa |
| If they don't match, reject | |
| If they do, cross it off | xbbaa# x bbaa |

Example

$$L_1 = \{w#w \mid w \in \{a, b\}^*\}$$

Strategy:

Find & remember first uncrossed symbol xbbaa#xbbaa

Cross it off xxbbaa#xbbaa

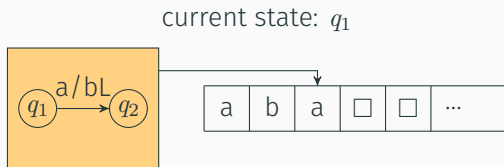
Read first symbol past # xxbaa#xbbaa

If they do, cross it off, else reject xxbaa#xxbbaa

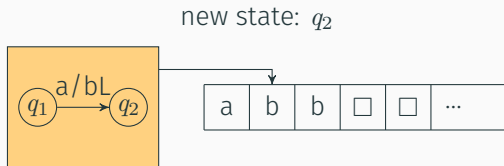
At the end, there should be only x's xxxxx#xxxxx

if so, accept; otherwise reject

How Turing machines operate



Replace a with b, and move head left



Computing devices: from practice to theory

Brief history of computing devices



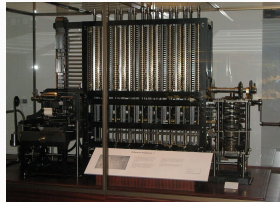
Antikythera Mechanism (~100BC)



Its reproduction



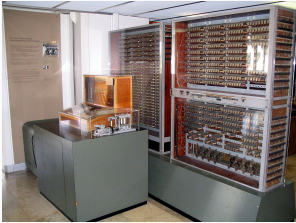
Abacus (Sumer 2700-2300BC, China 1200)



Babbage Difference engine (1840s)

Photo source: Wikipedia

Brief history of computing devices: programmable devices



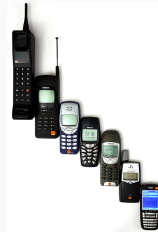
Z3 (Germany, 1941)



ENIAC (Pennsylvania, US, 1945)



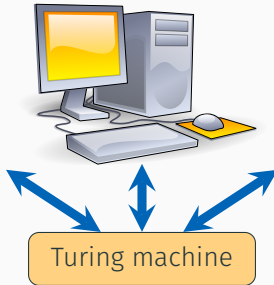
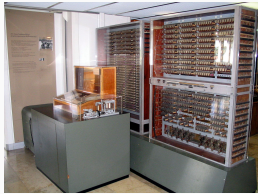
Personal computers (since 1970s)



Mobile phones

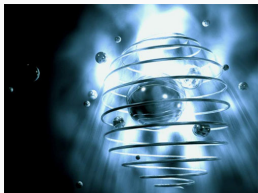
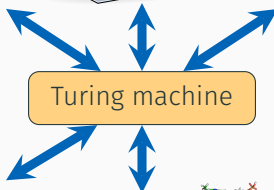
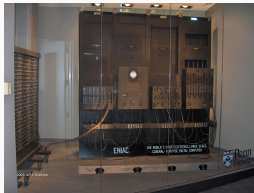
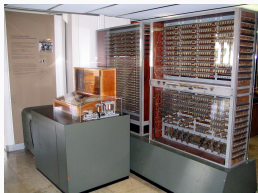
Photo source: Wikipedia

Computation is universal

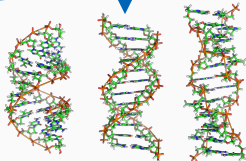


In principle, all computers have **the same** problem solving ability
If an algorithm can be implemented on any **realistic** computer, then
it can be implemented on a Turing machine

Church-Turing Thesis

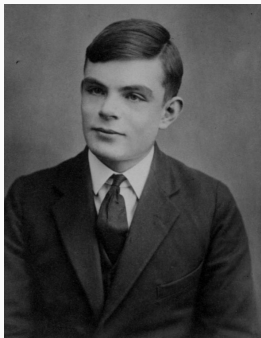


Quantum computing



DNA computing

Alan Turing



Alan Turing aged 16
(1912–1954)

Invented the **Turing Test** to tell apart humans from computers

Broke German encryption machines during World War II

Turing Award is the “Nobel prize of Computer Science”

Turing’s motivation: Understand the limitations of human computation by studying his “automatic machines”

Hilbert's Entscheidungsproblem, 1928 reformulation



David Hilbert

Entscheidungsproblem (Decision Problem)

“Write a program” to solve the following task:

Input: mathematical statement
(in first-order logic)

Output: whether the statement is true

In fact, he didn't ask to “write a program”, but to “design a procedure”

Examples of statements expressible in first-order logic:

Fermat's last theorem:

$$x^n + y^n = z^n$$

has no integer solution

for integer $n \geq 3$

Twin prime conjecture:

There are infinitely many pairs
of primes of the form p and

$$p + 2$$

Entscheidungsproblem (Decision Problem)

Design a procedure to solve the following task:

Input: mathematical statement
(in first-order logic)

Output: whether the statement is true

Church (1935-1936) and Turing (1936-1937) independently showed the procedure that Entscheidungsproblem asks for cannot exist!

Definitions of [procedure/algorithm](#):

[\$\lambda\$ -calculus](#) (Church) and [automatic machine](#) (Turing)

Church–Turing Thesis

Church–Turing Thesis

Intuitive notion of algorithms coincides with those implementable on Turing machines

Supporting arguments:

1. Turing machine is intuitive
2. Many independent definitions of “algorithms” turn out to be equivalent

References:

Alan Turing, “On Computable Numbers, with an Application to the Entscheidungsproblem”, 1937

Alonzo Church, “A Note on the Entscheidungsproblem”, 1936

Formal definition of Turing machine

A Turing Machine is $(Q, \Sigma, \Gamma, \delta, q_0, q_{\text{acc}}, q_{\text{rej}})$, where

- Q is a finite set of **states**
- Σ is the finite **input alphabet**, not containing the blank symbol \square
- Γ is the finite **tape alphabet** ($\Sigma \subset \Gamma$) including \square
- $q_0 \in Q$ is the **initial state**
- $q_{\text{acc}}, q_{\text{rej}} \in Q$ are the **accepting** and **rejecting states** ($q_{\text{acc}} \neq q_{\text{rej}}$)
- δ is the transition function

$$\delta : (Q \setminus \{q_{\text{acc}}, q_{\text{rej}}\}) \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}$$

Turing machines are deterministic

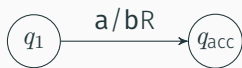
Configurations

A **configuration** consists of current state, head position, and tape contents

Configuration
(abbreviation)



$ab\ q_1\ a$



$abb\ q_{acc}$

Configurations

The **start configuration** of the TM on input w is $q_0 w$

We say a configuration C **yields** C' if the TM can go from C to C' in one step

Example: **ab** q_1 **a** yields **abb** q_{acc}

An **accepting configuration** is one that contains q_{acc}

A **rejecting configuration** is one that contains q_{rej}

The language of a Turing machine

A Turing machine M accepts x if there is a sequence of configurations C_0, C_1, \dots, C_k where

C_0 is starting C_i yields C_{i+1} C_k is accepting

The language recognized by M is the set of all strings that M accepts