

**1. Statements with many quantifiers.**

When we carefully analyse the logical structure of a mathematical statement, say,  $S$ , we will most likely find that  $S$  is of the form

$$(\mathbf{q}_x x)((\mathbf{q}_y y)(\cdots((\mathbf{q}_z z)((\mathbf{q}_w w)P(x, y, \cdots, z, w))) \cdots)),$$

in which:

- $P(x, y, \cdots, z, w)$  is a predicate with variables  $x, y, \cdots, z, w$ , and
- each of  $\mathbf{q}_x, \mathbf{q}_y, \dots, \mathbf{q}_z, \mathbf{q}_w$  stands for the universal quantifier  $\forall$  or the existential quantifier  $\exists$ .

Starting with the predicate  $P(x, y, \cdots, z, w)$ , we obtain successive predicates with fewer and fewer variables, arriving at the statement  $S$  in the end, by ‘closing the variables  $w, z, \cdots, y, x$  with quantifiers’ one by one:

- $P(x, y, \cdots, z, w)$ ,
- $(\mathbf{q}_w w)P(x, y, \cdots, z, w)$ ,
- $(\mathbf{q}_z z)((\mathbf{q}_w w)P(x, y, \cdots, z, w))$ ,
- ...
- $(\mathbf{q}_y y)(\cdots((\mathbf{q}_z z)((\mathbf{q}_w w)P(x, y, \cdots, z, w))) \cdots)$ ,
- $(\mathbf{q}_x x)((\mathbf{q}_y y)(\cdots((\mathbf{q}_z z)((\mathbf{q}_w w)P(x, y, \cdots, z, w))) \cdots))$ .

**2. Statements starting with two quantifiers.**

From a predicate  $Q(x, y)$  with two variables  $x, y$ , eight statements can be formed:

- |   |   |
|---|---|
| (1) $(\forall x)[(\forall y)Q(x, y)]$ . | (5) $(\forall x)[(\exists y)Q(x, y)]$ . |
| (2) $(\forall y)[(\forall x)Q(x, y)]$ . | (6) $(\exists y)[(\forall x)Q(x, y)]$ . |
| (3) $(\exists x)[(\exists y)Q(x, y)]$ . | (7) $(\exists x)[(\forall y)Q(x, y)]$ . |
| (4) $(\exists y)[(\exists x)Q(x, y)]$ . | (8) $(\forall y)[(\exists x)Q(x, y)]$ . |

We accept (1), (2) to be logically equivalent.

Examples of (1), (2).

- (a) For any  $x > 0$ , for any  $y > 0$ ,  $x + y > 0$ .
- (b) Let  $x, y \in \mathbb{Z}$ . Suppose  $x$  is divisible by  $y$  and  $y$  is divisible by  $x$ . Then  $|x| = |y|$ .

We accept (3), (4) to be logically equivalent (in most situations).

Examples of (3), (4).

- (a) There exist some irrational numbers  $x, y$  such that  $x + y$  is a rational number.
- (b) There exist some integers  $q, r$  such that  $10000 = 333q + r$  and  $0 \leq r \leq 332$ .

Care must be taken with (5), (6), (7), (8).

**3. Statements starting with one universal quantifier and one existential quantifier.**

Non-mathematical examples. Compare and contrast the statements in each pair (b), (#) below:

- (a) (b) Every student gets A in some MATH course.  
(No big deal; everyone has his/her own ‘lucky’ course.)
- (#) In some MATH course, every student gets A.  
(Then you will rush to enrol in such a course.)
- (b) (b) In every MATH course, some student gets A.  
(No big deal; you don’t expect us to be excessively harsh.)
- (#) Some student gets A in every MATH course.  
(Then you will look for ‘source’ from him/her.)

Now replace ‘A’ by ‘F’, and compare and contrast the resultant statements.



- In (‡),  $y$  does not ‘depend’ on  $x$ .

- (c) If you are in doubt, recall some examples which help you distinguish the meanings of (b) and (‡). For example, refer to ‘non-mathematical examples’.
- (d) Ask yourself whether what you write is the same as what you will be understood. For instance, if what you mean is

‘for any  $x \in \mathbb{R}$ , there exists some  $y \in \mathbb{R}$  such that  $x < y$ ’,

do not write the statement as

‘for any  $x \in \mathbb{R}$ ,  $x < y$  for some  $y \in \mathbb{R}$ ’,

or worse,

‘there exists some  $y \in \mathbb{R}$  such that  $x < y$  for any  $x \in \mathbb{R}$ ’.

#### 4. Negations of statements starting with two quantifiers.

We apply the rules for negating statements with one quantifier repeatedly for statements with two quantifiers:

- (a) The negation of ‘ $(\forall x)[(\exists y)Q(x, y)]$ ’ is ‘ $(\exists x)[(\forall y)(\sim Q(x, y))]$ ’.
- (b) The negation of ‘ $(\exists y)[(\forall x)Q(x, y)]$ ’ is ‘ $(\forall y)[(\exists x)(\sim Q(x, y))]$ ’.
- (c) The negation of ‘ $(\forall x)[(\forall y)Q(x, y)]$ ’ is ‘ $(\exists x)[(\exists y)(\sim Q(x, y))]$ ’.
- (d) The negation of ‘ $(\exists x)[(\exists y)Q(x, y)]$ ’ is ‘ $(\forall x)[(\forall y)(\sim Q(x, y))]$ ’.

Examples. How to write down the negations of the statements below?

- (a) *There exists some  $y \in S$  such that for any  $x \in T$ ,  $x < y$ .*

Convert the statement to be negated into a ‘chain of symbols’:

- $(\exists y \in S)[(\forall x \in T)(x < y)]$ .

Now repeatedly apply the rules for negating statements with one quantifier:

- $\sim\{(\exists y \in S)[(\forall x \in T)(x < y)]\}$  is equivalent to  $(\forall y \in S)\{\sim[(\forall x \in T)(x < y)]\}$ .

The latter is logically equivalent to  $(\forall y \in S)\{(\exists x \in T)[\sim(x < y)]\}$ .

The latter is logically equivalent to  $(\forall y \in S)[(\exists x \in T)(x \geq y)]$ .

Now convert this last ‘chain of symbols’ into words:

- ‘For any  $y \in S$ , there exists some  $x \in T$  such that  $x \geq y$ .’

This is the required negation of the given statement.

- (b) *For any  $a, b \in \mathbb{Z}$ ,  $a + b$  is divisible by 2.*

Negation: *There exist some  $a, b \in \mathbb{Z}$  such that  $a + b$  is not divisible by 2.*

- (c) *For any  $z \in \mathbb{C}$ , there exists some  $w \in \mathbb{R}$  such that  $\operatorname{Re}(z + w) = \operatorname{Im}(z + w)$ .*

Negation: *There exists some  $z \in \mathbb{C}$  such that for any  $w \in \mathbb{R}$ ,  $\operatorname{Re}(z + w) \neq \operatorname{Im}(z + w)$ .*

- (d) *There exist some  $s, t \in \mathbb{Q}$  such that  $(s + t \in \mathbb{Z} \text{ and } st \notin \mathbb{Z})$ .*

Negation: *For any  $s, t \in \mathbb{Q}$ ,  $(s + t \notin \mathbb{Z} \text{ or } st \in \mathbb{Z})$ .*

#### 5. Statements with many quantifiers.

The principles in the discussion above can be extended to statements with three or more quantifiers.

Questions. How to read and/or write them? How to negate them?

#### 6. Examples from linear algebra.

- (a) Let  $\mathbf{u}, \mathbf{v}, \mathbf{w} \in \mathbb{R}^n$ , and  $S$  be a subset of  $\mathbb{R}^n$ .

How to formulate ‘every vector in  $S$  is a linear combination of  $\mathbf{u}, \mathbf{v}, \mathbf{w}$  over  $\mathbb{R}$ ’?

- Formulation in words:

For any  $\mathbf{x} \in S$ , there exist some  $a, b, c \in \mathbb{R}$  such that  $\mathbf{x} = a\mathbf{u} + b\mathbf{v} + c\mathbf{w}$ .

- Formulation in symbols:

$(\forall \mathbf{x} \in S)[(\exists a \in \mathbb{R})(\exists b \in \mathbb{R})(\exists c \in \mathbb{R})(\mathbf{x} = a\mathbf{u} + b\mathbf{v} + c\mathbf{w})]$ .

This statement is of the form  $(\forall \mathbf{x})((\exists a)(\exists b)(\exists c)Q(\mathbf{x}, a, b, c))$ .

How to formulate ‘not every vector in  $S$  is a linear combination of  $\mathbf{u}, \mathbf{v}, \mathbf{w}$  over  $\mathbb{R}$ ’?

- Formulation in symbols:  
 $(\exists \mathbf{x} \in S)[(\forall a \in \mathbb{R})(\forall b \in \mathbb{R})(\forall c \in \mathbb{R})(\mathbf{x} \neq a\mathbf{u} + b\mathbf{v} + c\mathbf{w})]$ .

- Formulation in words:  
*There exists some  $\mathbf{x} \in S$  such that for any  $a, b, c \in \mathbb{R}$ ,  $\mathbf{x} \neq a\mathbf{u} + b\mathbf{v} + c\mathbf{w}$ .*

(b) Let  $\mathbf{u}, \mathbf{v}, \mathbf{w} \in \mathbb{R}^n$ .

How to formulate ' $\mathbf{u}, \mathbf{v}, \mathbf{w}$  are linearly independent over  $\mathbb{R}$ '?

- Formulation in words:  
*For any  $a, b, c \in \mathbb{R}$ , if  $a\mathbf{u} + b\mathbf{v} + c\mathbf{w} = \mathbf{0}$  then ( $a = 0$  and  $b = 0$  and  $c = 0$ ).*
- Formulation in symbols:  
 $(\forall a \in \mathbb{R})(\forall b \in \mathbb{R})(\forall c \in \mathbb{R})\{(a\mathbf{u} + b\mathbf{v} + c\mathbf{w} = \mathbf{0}) \longrightarrow [(a = 0) \wedge (b = 0) \wedge (c = 0)]\}$ .

This statement is of the form  $(\forall a)(\forall b)(\forall c)Q(a, b, c)$ .

How to formulate ' $\mathbf{u}, \mathbf{v}, \mathbf{w}$  are linearly dependent over  $\mathbb{R}$ '?

- Formulation in symbols:  
 $(\exists a \in \mathbb{R})(\exists b \in \mathbb{R})(\exists c \in \mathbb{R})\{(a\mathbf{u} + b\mathbf{v} + c\mathbf{w} = \mathbf{0}) \wedge [(a \neq 0) \vee (b \neq 0) \vee (c \neq 0)]\}$ .
- Formulation in words:  
*There exist some  $a, b, c \in \mathbb{R}$  such that ( $a\mathbf{u} + b\mathbf{v} + c\mathbf{w} = \mathbf{0}$  and  $a, b, c$  are not all 0).*

## 7. Examples from calculus of one variable.

(a) Let  $f$  be a real-valued function on  $\mathbb{R}$ , and  $c \in \mathbb{R}$ .

How to formulate ' $f$  attains a relative minimum at  $c$ '?

- Formulation in words:  
*There exists some  $\delta > 0$  such that for any  $x \in \mathbb{R}$ , (if  $|x - c| < \delta$  then  $f(x) \geq f(c)$ ).*
- Formulation in symbols:  
 $(\exists \delta > 0)\{(\forall x \in \mathbb{R})[ (|x - c| < \delta) \longrightarrow (f(x) \geq f(c)) ]\}$ .

This statement is of the form  $(\exists \delta)((\forall x)Q(\delta, x))$ .

How to formulate ' $f$  does not attain a relative minimum at  $c$ '?

- Formulation in symbols:  
 $(\forall \delta > 0)\{(\exists x \in \mathbb{R})[ (|x - c| < \delta) \wedge (f(x) < f(c)) ]\}$ .
- Formulation in words:  
*For any  $\delta > 0$ , there exists some  $x \in \mathbb{R}$  such that ( $|x - c| < \delta$  and  $f(x) < f(c)$ ).*

(b) Let  $f$  be a real-valued function on  $\mathbb{R}$ , and  $c \in \mathbb{R}$ .

How to formulate ' $f$  is continuous at  $c$ '?

- Formulation in words:  
*For any  $\varepsilon > 0$ , there exists some  $\delta > 0$  such that for any  $x \in \mathbb{R}$ , (if  $|x - c| < \delta$  then  $|f(x) - f(c)| < \varepsilon$ ).*
- Formulation in symbols:  
 $(\forall \varepsilon > 0)\{(\exists \delta > 0)[(\forall x \in \mathbb{R})((|x - c| < \delta) \longrightarrow (|f(x) - f(c)| < \varepsilon))]\}$ .

This statement is of the form  $(\forall \varepsilon)[(\exists \delta)((\forall x)Q(\varepsilon, \delta, x))]$ .

How to formulate ' $f$  is not continuous at  $c$ '?

- Formulation in symbols:  
 $(\exists \varepsilon > 0)\{(\forall \delta > 0)[(\exists x \in \mathbb{R})((|x - c| < \delta) \wedge (|f(x) - f(c)| \geq \varepsilon))]\}$ .
- Formulation in words:  
*There exists some  $\varepsilon > 0$  such that for any  $\delta > 0$ , there exists some  $x \in \mathbb{R}$  such that ( $|x - c| < \delta$  and  $|f(x) - f(c)| \geq \varepsilon$ ).*

Now you see why *mathematical analysis* is hard: even very basic notions involve a heavy presence of quantifiers.