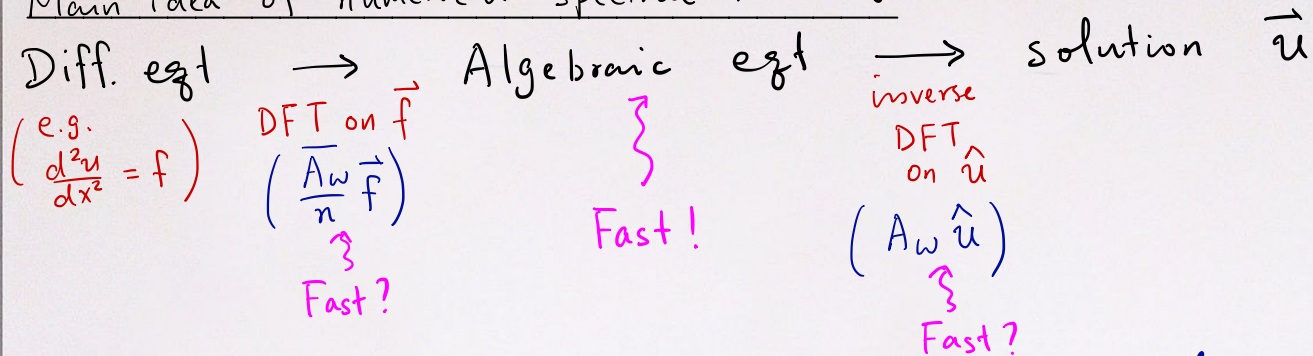


## Lecture 12:

Main idea of numerical spectral method



Remark: To develop an efficient numerical spectral method, we need to compute  $A_w \hat{u}$  and  $\frac{\overline{A_w} \vec{f}}{n}$  fast.

- Computational cost for  $A_w \hat{u}$  is  $\mathcal{O}(n^2)$ .  
( $n \times n$ )

Goal: Reduce the computational cost to  $\mathcal{O}(n \log n)$

e.g.  $n = 2^{10}$ ,  $n^2 = 2^{20}$ ,  $n \log n = 10 \cdot 2^{10} < 2^{14}$ .  $\therefore 2^6 = 64$  times faster!

# Fast Fourier Transform (FFT) (Colley and Tukey, 1965)

Let  $F_n = \begin{pmatrix} 1 & 1 & \dots & 1 \\ 1 & \omega_n & \dots & \omega_n^{n-1} \\ \vdots & \vdots & \dots & \vdots \\ 1 & \omega_n^{n-1} & \dots & \omega_n^{(n-1)^2} \end{pmatrix}$  where  $\omega_n = e^{i\left(\frac{2\pi}{n}\right)}$

Let  $\vec{y} = F_n \vec{x}$ , where  $\vec{y} = \begin{pmatrix} y_0 \\ y_1 \\ \vdots \\ y_{n-1} \end{pmatrix}$  and  $\vec{x} = \begin{pmatrix} x_0 \\ x_1 \\ \vdots \\ x_{n-1} \end{pmatrix}$ . Suppose  $n=2m$ .

Then, for each  $0 \leq j \leq n-1$ ,

$$y_j = \sum_{k=0}^{n-1} \omega_n^{jk} x_k = \sum_{k=0}^{2m-1} \omega_{2m}^{kj} x_k.$$

Divide  $k=0, 1, 2, \dots, 2m-1$  into two parts:

Part 1:  $0, 2, 4, 6, \dots, 2(m-1)$  (Even)

Part 2:  $1, 3, 5, 7, \dots, 2m-1$  (Odd)

$$\text{Then: } y_j = \underbrace{\sum_{k=0}^{m-1} \omega_n^{2kj} X_{2k}}_{\text{Part 1}} + \underbrace{\sum_{k=0}^{m-1} \omega_n^{(2k+1)j} X_{2k+1}}_{\text{Part 2}}$$

$$= \sum_{k=0}^{m-1} \omega_m^{kj} \vec{X}'_k + \sum_{k=0}^{m-1} \omega_n^{\dot{j}} \omega_m^{kj} \vec{X}''_k$$

$\left( \because \omega_{2m}^{2k} = e^{i\left(\frac{2\pi}{2m}\right)2k} \right)$   
 $= e^{i\left(\frac{2\pi}{m}\right)k}$   
 $= \omega_m^k$

Denote  $\vec{X}' = \begin{pmatrix} X_0 \\ X_2 \\ \vdots \\ X_{2m-2} \end{pmatrix}$ ,  $\vec{X}'' = \begin{pmatrix} X_1 \\ X_3 \\ \vdots \\ X_{2m-1} \end{pmatrix}$ . Let  $\vec{y}' = F_m \vec{X}'$  and  $\vec{y}'' = F_m \vec{X}''$ .

$$\therefore y_j = \overbrace{(F_m \vec{X}')_j}^{m \times m} + \omega_n^{\dot{j}} \overbrace{(F_m \vec{X}'')_j}^{m \times m} \text{ for } j=0, 1, 2, \dots, m-1$$

$$= (\vec{y}')_j + \omega_n^{\dot{j}} (\vec{y}'')_j$$

$\uparrow$   
j-th entry of  $\vec{y}'$

$\uparrow$   
j-th entry of  $\vec{y}''$

$$y_{j+m} = \sum_{k=0}^{m-1} \omega_n^{2k(j+m)} x_{2k} + \sum_{k=0}^{m-1} \omega_n^{(2k+1)(j+m)} x_{2k+1} \quad \text{for } j=0, 1, 2, \dots, m-1$$

↑  
Capture

$$y_m, y_{m+1}, \dots, y_{2m-1} = \sum_{k=0}^{m-1} \omega_m^{kj} \omega_m^{km} (\bar{x}')_k + \sum_{k=0}^{m-1} \omega_m^{k(j+m)} \omega_n^{j+m} (\bar{x}'')_k$$

$e^{i(\frac{2\pi}{m})km}$   
 $\omega_m^{kj} \omega_m^{km}$   
 $\omega_m^{kj} \omega_n^{j+m}$   
 $\omega_n^{j+m} e^{i(\frac{2\pi}{2m}) \cdot m}$

$$\therefore y_{j+m} = \sum_{k=0}^{m-1} \omega_m^{kj} (\bar{x}')_k - \omega_n^j \sum_{k=0}^{m-1} \omega_m^{kj} (\bar{x}'')_k - \omega_n^j$$

$$y_{j+m} = \underbrace{(F_m \bar{x}')}_m - \omega_n^j \underbrace{(F_m \bar{x}'')}_m \quad \text{for } j=0, 1, 2, \dots, m-1$$

Note:  $n \times n$  matrix multiplication becomes  $\frac{n}{2} \times \frac{n}{2} = m \times m$  matrix multiplication.

For simplicity, we denote:  $\begin{pmatrix} v_1 \\ v_2 \\ \vdots \\ v_n \end{pmatrix} \otimes \begin{pmatrix} w_1 \\ w_2 \\ \vdots \\ w_n \end{pmatrix} = \begin{pmatrix} v_1 w_1 \\ v_2 w_2 \\ \vdots \\ v_n w_n \end{pmatrix}$

Then:  $\begin{pmatrix} y_0 \\ y_1 \\ \vdots \\ y_{m-1} \end{pmatrix} = \vec{y}' + \begin{pmatrix} w_n^0 \\ w_n^1 \\ \vdots \\ w_n^{m-1} \end{pmatrix} \otimes \vec{y}''$  and  $\begin{pmatrix} y_m \\ y_{m+1} \\ \vdots \\ y_{2m-1} \end{pmatrix} = \vec{y}' - \begin{pmatrix} w_n^0 \\ w_n^1 \\ \vdots \\ w_n^{m-1} \end{pmatrix} \otimes \vec{y}''$

### Summary of FFT

Step 1: Split  $\vec{x}$  into  $\vec{x}' = \begin{pmatrix} x_0 \\ x_2 \\ \vdots \\ x_{2(m-1)} \end{pmatrix}$  and  $\vec{x}'' = \begin{pmatrix} x_1 \\ x_3 \\ \vdots \\ x_{2m-1} \end{pmatrix}$  " $m \times m$ "

Step 2: Compute  $\vec{y}' = F_m \vec{x}'$  and  $\vec{y}'' = F_m \vec{x}''$ , where  $F_m = \frac{n}{2} \times \frac{n}{2}$  " $\mathcal{O}(m^2)$ " matrix

Step 3: Compute:  $\vec{w}_m$  " $\mathcal{O}(m)$ "

$\begin{pmatrix} y_0 \\ y_1 \\ \vdots \\ y_{m-1} \end{pmatrix} = \vec{y}' + \begin{pmatrix} w_n^0 \\ w_n^1 \\ \vdots \\ w_n^{m-1} \end{pmatrix} \otimes \vec{y}''$  " $\mathcal{O}(m)$ " and  $\begin{pmatrix} y_m \\ y_{m+1} \\ \vdots \\ y_{2m-1} \end{pmatrix} = \vec{y}' - \begin{pmatrix} w_n^0 \\ w_n^1 \\ \vdots \\ w_n^{m-1} \end{pmatrix} \otimes \vec{y}''$  " $\mathcal{O}(m)$ "

Remark: Computational cost =  $O(m^2)$  +  $O(m)$  (multiplication + addition)  
ss  
 $O(m^2)$

Computational cost for FFT: Assume  $n = 2^l$ .

Let  $C_m$  = computational cost of  $F_m$ . Then  $C_1 = 1$ .

Claim:  $C_{2m} = 2C_m + 3m$

Proof: Step 2:  $\vec{y}' = F_m \vec{x}'$ ,  $\vec{y}'' = F_m \vec{x}''$  (=  $2C_m$ )

Step 3:  $y_j = \vec{y}'_j + \omega_n^j \vec{y}''_j$   
 $y_{j+m} = \vec{y}'_j - \omega_n^j \vec{y}''_j$   
for  $j = 0, 1, 2, \dots, m-1$

( = 1 multiplication  
+  
1 addition  
+  
1 subtraction )

↓  
Total:  $3m$

$\therefore C_{2m} = 2C_m + 3m$ .

Now,  $n = 2^l$ .

$$\therefore C_{2^l} = 2C_{2^{l-1}} + 3 \cdot 2^{l-1}$$

$$\begin{aligned}\therefore 2^{-l}C_{2^l} &= 2^{-(l-1)}C_{2^{l-1}} + \frac{3}{2} = 2^{-(l-2)}C_{2^{l-2}} + 2\left(\frac{3}{2}\right) \\ &= \vdots\end{aligned}$$

$$\begin{aligned}\therefore C_{2^l} &= \underbrace{2^l}_n + \frac{3}{2} l \underbrace{2^l}_n = 2^0 C_{2^0} + l\left(\frac{3}{2}\right) = 1 + \frac{3}{2}l \\ &= n + \frac{3}{2}n \log_2 n = \mathcal{O}(n \log_2 n)\end{aligned}$$

# Butterfly diagram (Algorithmic visualization)

Consider  $F_4$  (4x4 matrix).

[ Recall:  $\vec{y}_e = \vec{y}' = F_m \vec{x}'$       Denote  $\vec{x}' := \vec{x}_e$ ,  $\vec{x}'' = \vec{x}_0$  ]  
 $\vec{y}_o = \vec{y}'' = F_m \vec{x}''$

$$\vec{y}_e := F_2 \vec{x}_e = \begin{pmatrix} x_0 \\ x_2 \end{pmatrix} = \begin{pmatrix} \vdots \\ \vdots \end{pmatrix} \xrightarrow{\quad} \begin{pmatrix} y_0 \\ y_1 \end{pmatrix}$$

$$\vec{y}_o := F_2 \vec{x}_0 = \begin{pmatrix} x_1 \\ x_3 \end{pmatrix} = \begin{pmatrix} \vdots \\ \vdots \end{pmatrix} \xrightarrow{\quad} \begin{pmatrix} y_2 \\ y_3 \end{pmatrix}$$

$\vec{w}_4 = \begin{pmatrix} w_4^0 \\ w_4^1 \end{pmatrix}$

Diagram means:

$$\begin{pmatrix} y_0 \\ y_1 \end{pmatrix} = \vec{y}_e + \vec{w}_4 \otimes \vec{y}_o$$

$$\begin{pmatrix} y_2 \\ y_3 \end{pmatrix} = \vec{y}_o - \vec{w}_4 \otimes \vec{y}_o$$

} Depend on  $\vec{y}_o$  and  $\vec{y}_e$ .



For  $F_2 \vec{x}_e$ :  $\overset{1}{=} F_1 \vec{x}_{ee} = (x_0)$   $\overset{2}{=} F_1 \vec{x}_{e0} = (x_2)$

$$\begin{pmatrix} \cdot \\ \cdot \end{pmatrix} = F_2 \vec{x}_e = \vec{y}_e$$

For  $F_2 \vec{x}_0$ :  $\overset{1}{=} F_1 \vec{x}_{0e} = (x_1)$   $\overset{2}{=} F_1 \vec{x}_{00} = (x_3)$

$$\begin{pmatrix} \cdot \\ \cdot \end{pmatrix} = F_2 \vec{x}_0 = \vec{y}_0$$

Remark:

$$\vec{x}_e = \begin{pmatrix} x_0 \\ x_2 \end{pmatrix}, \quad \vec{x}_0 = \begin{pmatrix} x_1 \\ x_3 \end{pmatrix}$$

$$\downarrow$$

$$\vec{x}_{ee} = (x_0)$$

$$\vec{x}_{0e} = (x_1)$$

$$\vec{x}_{e0} = (x_2)$$

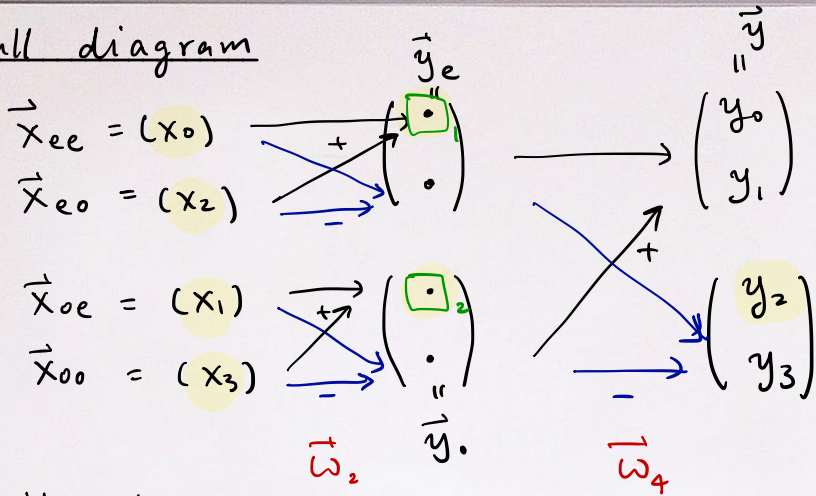
$$\vec{x}_{00} = (x_3)$$

$$\vec{w}_2^0 = (w_2^0) = (1)$$

$$\vec{y}_e = \begin{pmatrix} x_0 + \vec{w}_2^0 \otimes (x_2) \\ x_0 - \vec{w}_2^0 \otimes (x_2) \end{pmatrix} = \begin{pmatrix} x_0 + x_2 \\ x_0 - x_2 \end{pmatrix}$$

$$\vec{y}_0 = \begin{pmatrix} x_1 + \vec{w}_2^0 \otimes (x_3) \\ x_1 - \vec{w}_2^0 \otimes (x_3) \end{pmatrix} = \begin{pmatrix} x_1 + x_3 \\ x_1 - x_3 \end{pmatrix}$$

# Overall diagram



Using the diagram, find  $y_2$ .

$$y_2 = \square_1 - (\vec{W}_4)_0 \cdot \square_2 \quad ; \quad \square_1 = X_0 + X_2$$

$$= \square_1 - \square_2 \quad \square_2 = X_1 + X_3$$

$$\therefore y_2 = X_0 + X_2 - (X_1 + X_3)$$

# Iterative method to solve huge linear system

Recall: Numerical spectral method handles periodic functions.

Consider: (\*)  $\frac{d^2 u}{dx^2} = f$ ,  $u(0) = A$ ,  $u(1) = B$   
 $x \in [0, 1]$ .

Partition  $[0, 1]$  into  $x_j = jh$  where  $h = \frac{1}{n+1}$

Then: (\*) is discretized as:

$$\frac{u_{i+1} - 2u_i + u_{i-1}}{h^2} = f(x_i) \quad \text{or}$$

for all  $i=1, 2, \dots, n$

$$\tilde{A} \begin{pmatrix} u_1 \\ u_2 \\ \vdots \\ u_n \end{pmatrix} = \begin{pmatrix} f(x_1) - \frac{A}{h^2} \\ f(x_2) \\ \vdots \\ f(x_n) - \frac{B}{h^2} \end{pmatrix}$$

$u_i \stackrel{\text{def}}{=} u(x_i)$

$D = \begin{pmatrix} -2 & 1 & 0 & \dots & 1 \\ 1 & -2 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & \dots & 1 & -2 & 1 \end{pmatrix}$  has eigenvector  $e^{itx}$

where

$$\tilde{A} = \frac{1}{h^2} \begin{pmatrix} -2 & 1 & & & 0 \\ & 1 & -2 & 1 & \\ & & \ddots & \ddots & \\ 0 & & & 1 & -2 \end{pmatrix}$$

Question: How to solve BIG linear system?

Method 1: Gaussian elimination

Comp. cost:  $\mathcal{O}(n^3)$

Sol: exact.

Method 2: LU factorization.

Decompose  $A = LU$   
           $\uparrow$   $\nwarrow$   
          lower upper

(If  $A$  is SPD, then  $A = LL^T$   
by Cholesky decomposition)

$LU\vec{x} = \vec{b}$  by solving  $\begin{cases} L\vec{y} = \vec{b} \\ U\vec{x} = \vec{y} \end{cases}$  — easy

Comp. cost:  $\mathcal{O}(n^3)$

Sol: exact.

Goal: Develop iterative method: find a sequence  $\vec{x}_0, \vec{x}_1, \vec{x}_2, \dots$   
such that  $\vec{x}_k \rightarrow \vec{x}^* = \text{sol. of } A\vec{x} = \vec{f}$  as  $k \rightarrow \infty$ .

Remark: We can stop when error is small enough.

Method: Splitting method

Consider a linear system  $A\vec{x} = \vec{f}$  where  $A \in M_{n \times n}$  ( $n$  is BIG)

Split  $A$  as follows:  $A = N + (A - N) = N - \underbrace{(N - A)}_P$

Then:  $A\vec{x} = \vec{f} \Leftrightarrow (N - P)\vec{x} = \vec{f} \Leftrightarrow N\vec{x} = P\vec{x} + \vec{f}$

Develop an iterative scheme as follows:

$$(\star) N\vec{x}^{n+1} = P\vec{x}^n + \vec{f}$$

If  $\{\vec{x}^n\}_{n=1}^{\infty}$  converges, then it converges to the sol  $\vec{x}^*$  of  $A\vec{x} = \vec{f}$

- Remark:
- $N$  should be simple = easy to find inverse.
  - $N$  should have an inverse
  - $N$  should be "related to"  $A$ .