# **Church-Turing Thesis**

CSCI 3130 Formal Languages and Automata Theory

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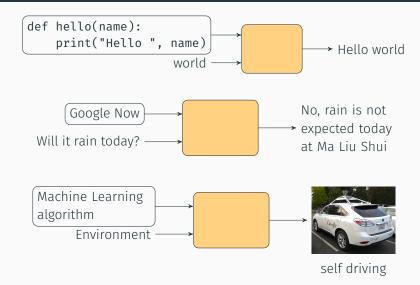
Chinese University of Hong Kong

#### What is a computer?

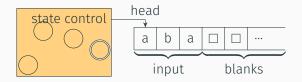


A computer is a machine that manipulates data according to a list of instructions

## What is a computer?



### Turing machines



Can both read from and write to the tape

Head can move both left and right

Unlimited tape space

Has two special states accept and reject

#### Example

$$L_1 = \{ w \# w \mid w \in \{ a, b \}^* \}$$

#### Strategy:

Read and remember the first symbol <u>abbaa#abbaa</u>

Cross it off <u>x</u>bbaa#abbaa

Read the first symbol past # xbbaa#<u>a</u>bbaa

If they don't match, reject

If they do, cross it off  $xbbaa#\underline{x}bbaa$ 

### Example

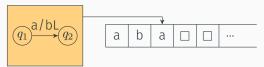
$$L_1 = \{ w \# w \mid w \in \{a, b\}^* \}$$

#### Strategy:

Find & remember first uncrossed symbol	x <u>b</u> baa#xbbaa
Cross it off	x <u>x</u> baa#xbbaa
Read first symbol past #	xxbaa#x <u>b</u> baa
If they do, cross it off, else reject	xxbaa#x <u>x</u> baa
At the end, there should be only $\mathbf{x}$ 's	xxxxx#xxxx <u>x</u>
if so, accept; otherwise reject	

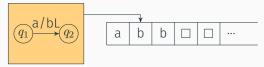
## How Turing machines operate

current state:  $q_1$ 



Replace a with b, and move head left

new state:  $q_2$ 



Computing devices: from practice

to theory

# Brief history of computing devices



Antikythera Mechanism (~100BC)



Abacus (Sumer 2700-2300BC, China 1200)



Its reproduction



Babbage Difference engine (1840s)

Photo source: Wikipedia

## Brief history of computing devices: programmable devices



Z3 (Germany, 1941)



Personal computers (since 1970s)



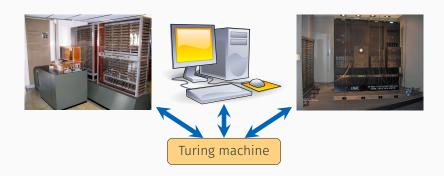
ENIAC (Pennsylvania, US, 1945)



Mobile phones

Photo source: Wikipedia

## Computation is universal

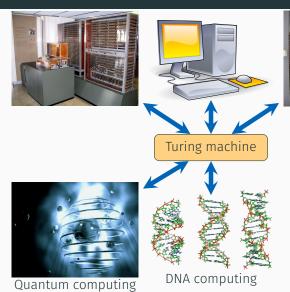


In principle, all computers have the same problem solving ability

If an algorithm can be implemented on any realistic computer, then

it can be implemented on a Turing machine

# Church-Turing Thesis



# **Alan Turing**



Alan Turing aged 16 (1912–1954)

Invented the Turing Test to tell apart humans from computers

Broke German encryption machines during World War II

Turing Award is the "Nobel prize of Computer Science"

Turing's motivation: Understand the limitations of human computation by studying his "automatic machines"

#### Hilbert's Entscheidungsproblem, 1928 reformulation



David Hilbert

#### Entscheidungsproblem (Decision Problem)

"Write a program" to solve the following task:

**Input:** mathematical statement

(in first-order logic)

Output: whether the statement is true

In fact, he didn't ask to "write a program", but to "design a procedure"

Examples of statements expressible in first-order logic:

#### Fermat's last theorem:

$$x^n + y^n = z^n$$
  
has no integer solution  
for integer  $n \geqslant 3$ 

#### Twin prime conjecture:

There are infinitely many pairs of primes of the form p and p+2

#### Undecidability

#### Entscheidungsproblem (Decision Problem)

Design a procedure to solve the following task:

Input: mathematical statement

(in first-order logic)

Output: whether the statement is true

Church (1935-1936) and Turing (1936-1937) independently showed the procedure that Entscheidungsproblem asks for cannot exist!

Definitions of procedure/algorithm:

 $\lambda$ -calculus (Church) and automatic machine (Turing)

### Church-Turing Thesis

#### Church-Turing Thesis

Intuitive notion of algorithms coincides with those implementable on Turing machines

#### Supporting arguments:

- 1. Turing machine is intuitive
- Many independent definitions of "algorithms" turn out to be equivalent

#### References:

Alan Turing, "On Computable Numbers, with an Application to the Entscheidungsproblem", 1937

Alonzo Church, "A Note on the Entscheidungsproblem", 1936

#### Formal definition of Turing machine

A Turing Machine is  $(Q, \Sigma, \Gamma, \delta, q_0, q_{\text{acc}}, q_{\text{rej}})$ , where

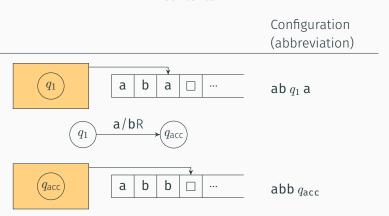
- · Q is a finite set of states
- $\Sigma$  is the finite input alphabet, not containing the blank symbol  $\square$
- $\Gamma$  is the finite tape alphabet  $(\Sigma \subset \Gamma)$  including  $\square$
- $q_0 \in Q$  is the initial state
- $q_{
  m acc},\,q_{
  m rej}\in \mathit{Q}$  are the accepting and rejecting states ( $q_{
  m acc} 
  eq q_{
  m rej}$ )
- $\cdot$   $\delta$  is the transition function

$$\delta: (Q \setminus \{q_{\mathrm{acc}}, q_{\mathrm{rej}}\}) \times \Gamma \to Q \times \Gamma \times \{\mathsf{L}, \mathsf{R}\}$$

Turing machines are deterministic

#### Configurations

A configuration consists of current state, head position, and tape contents



#### Configurations

The start configuration of the TM on input w is  $q_0$  w

We say a configuration C yields C' if the TM can go from C to C' in one step

Example:  $ab q_1 a$  yields  $abb q_{acc}$ 

An accepting configuration is one that contains  $q_{
m acc}$ A rejecting configuration is one that contains  $q_{
m rej}$ 

#### The language of a Turing machine

A Turing machine M accepts x if there is a sequence of configurations  $C_0, C_1, \ldots, C_k$  where

 $C_0$  is starting  $C_i$  yields  $C_{i+1}$   $C_k$  is accepting

The language recognized by M is the set of all strings that M accepts