Church–Turing Thesis CSCI 3130 Formal Languages and Automata Theory

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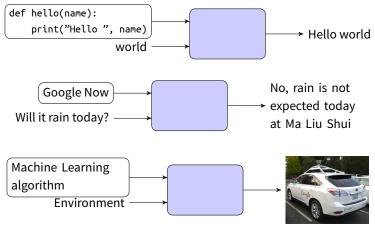
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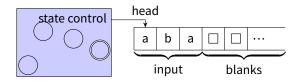
A computer is a machine that manipulates data according to a list of instructions

What is a computer?





Turing machines



Can both read from and write to the tape

Head can move both left and right

Unlimited tape space

Has two special states accept and reject

Example

 $L_1 = \{ w \# w \mid w \in \{ \mathsf{a}, \mathsf{b} \}^* \}$

Strategy:

Read and remember the first symbol	<u>a</u> bbaa#abbaa
Cross it off	<u>x</u> bbaa#abbaa
Read the first symbol past #	xbbaa# <u>a</u> bbaa
If they don't match, reject	
If they do, cross it off	xbbaa# <u>x</u> bbaa

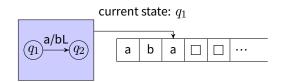
Example

$L_1 = \{ w \# w \mid w \in \{ \mathsf{a}, \mathsf{b} \}^* \}$

Strategy:

Look for and remember the first uncrossed symbol	x <u>b</u> baa#xbbaa
Cross it off	x <u>x</u> baa#xbbaa
Read the first symbol past #	xxbaa#x <u>b</u> baa
If they do, cross it off, else reject	xxbaa#x <u>x</u> baa
At the end, there should be only x's	xxxxx#xxxx <u>x</u>
if so, accept; otherwise reject	

How Turing machines operate



Replace a with b, and move head left

new state: q_2





Computing devices: from practice to theory

Brief history of computing devices



Antikythera Mechanism (~100BC)



Abacus (Sumer 2700-2300BC, China 1200)



Its reproduction



Babbage Difference engine (1840s)

Photo source: Wikipedia

Brief history of computing devices: programmable devices



Z3 (Germany, 1941)



Personal computers (since 1970s)



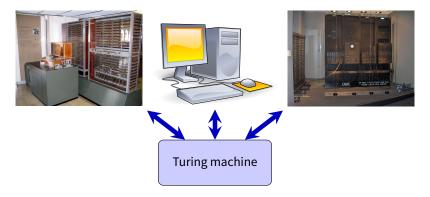
ENIAC (Pennsylvania, US, 1945)



Mobile phones

Photo source: Wikipedia

Computation is universal



In principle, all computers have the same problem solving ability

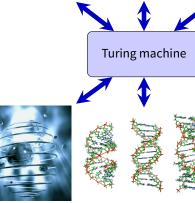
If an algorithm can be implemented on any realistic computer, then it can be implemented on a Turing machine

Church–Turing Thesis









Quantum computing

DNA computing

Alan Turing



Alan Turing aged 16 (1912-1954)

Invented the Turing Test to tell apart humans from computers

Broke German encryption machines during World War II

Turing Award is the "Nobel prize of Computer Science"

Turing's motivation: Understand the limitations of human computation by studying his "automatic machines"

Hilbert's Entscheidungsproblem, 1928 reformulation



David Hilbert

Entscheidungsproblem (Decision Problem)

"Write a program" to solve the following task: Input: mathematical statement (in first-order logic) Output: whether the statement is true

In fact, he didn't ask to "write a program", but to "design a procedure"

Examples of statements expressible in first-order logic:

Fermat's last theorem:

 $\begin{aligned} x^n + y^n &= z^n \\ \text{has no integer solution} \\ \text{for integer } n \geqslant 3 \end{aligned}$

Twin prime conjecture:

There are infinitely many pairs of primes of the form p and p + 2

Undecidability

Entscheidungsproblem (Decision Problem)

Design a procedure to solve the following task: Input: mathematical statement (in first-order logic) Output: whether the statement is true

Church (1935-1936) and Turing (1936-1937) independently showed the procedure that Entscheidungsproblem asks for cannot exist!

 $\label{eq:loss} \begin{array}{l} \mbox{Definitions of procedure/algorithm:} \\ \lambda\mbox{-calculus (Church) and automatic machine (Turing)} \end{array}$

Church-Turing Thesis

Church-Turing Thesis

Intuitive notion of algorithms coincides with those implementable on Turing machines

Supporting arguments:

- 1. Turing machine is intuitive
- 2. Many independent definitions of "algorithms" turn out to be equivalent

References:

Alan Turing, "On Computable Numbers, with an Application to the Entscheidungsproblem", 1937 Alonzo Church, "A Note on the Entscheidungsproblem", 1936

Formal definition of Turing machine

A Turing Machine is $(Q, \Sigma, \Gamma, \delta, q_0, q_{\mathsf{acc}}, q_{\mathsf{rej}})$, where

- Q is a finite set of states
- Σ is the input alphabet, not containing the blank symbol \Box
- Γ is the tape alphabet ($\Sigma \subseteq \Gamma$) including \Box
- $q_0 \in Q$ is the initial state
- ▶ $q_{\sf acc}, q_{\sf rej} \in Q$ are the accepting and rejecting states
- $\blacktriangleright \delta$ is the transition function

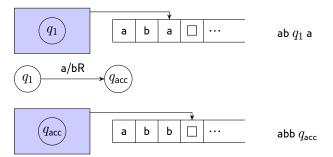
$$\delta: (Q \setminus \{q_{\mathsf{acc}}, q_{\mathsf{rej}}\}) \times \Gamma \to Q \times \Gamma \times \{\mathsf{L}, \mathsf{R}\}$$

Turing machines are deterministic

Configurations

A configuration consists of current state, head position, and tape contents

Configuration (abbreviation)



The start configuration of the TM on input w is $q_0 w$

We say a configuration C yields C' if the TM can go from C to C' in one step $\mbox{Example:} \quad \mbox{ab } q_1 \mbox{a } \mbox{yields} \quad \mbox{abb } q_{\rm acc}$

An accepting configuration is one that contains q_{acc} A rejecting configuration is one that contains q_{rej}

The language of a Turing machine

A Turing machine M accepts x if there is a sequence of configurations $C_0,\,C_1,\,\ldots,\,C_k \text{ where }$

 C_0 is starting C_i yields C_{i+1} C_k is accepting

The language recognized by M is the set of all strings that M accepts