

Nondeterministic Finite Automata

CSCI 3130 Formal Languages and Automata Theory

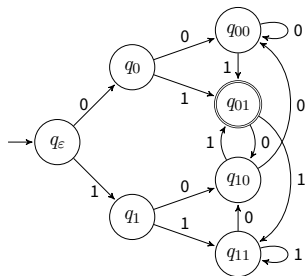
Siu On CHAN

Chinese University of Hong Kong

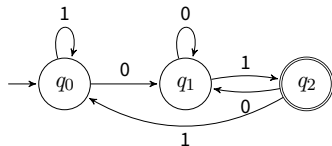
Fall 2017

Example from last lecture with a simpler solution

Construct a DFA over alphabet $\{0, 1\}$ that accepts all strings ending in 01



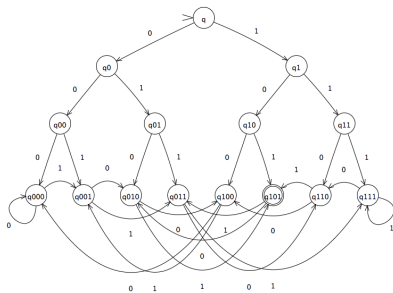
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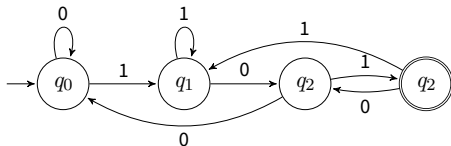
Three weeks later: DFA minimization

Another example from last lecture

Construct a DFA over alphabet $\{0, 1\}$ that accepts all strings ending in 101

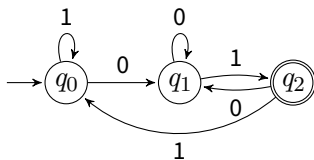


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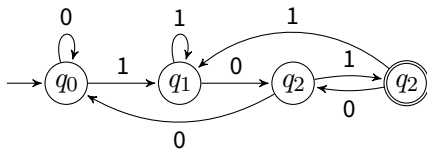


String matching DFAs

Ending in **01**



Ending in **101**



Fast string matching algorithms to turn a **pattern** into a string matching DFA

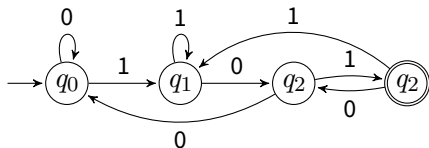
and execute the DFA:

Boyer–Moore (BM) and Knuth–Morris–Pratt (KMP)

(won't cover in class)

Nondeterminism

In a few lectures



What problems can finite state machines solve?

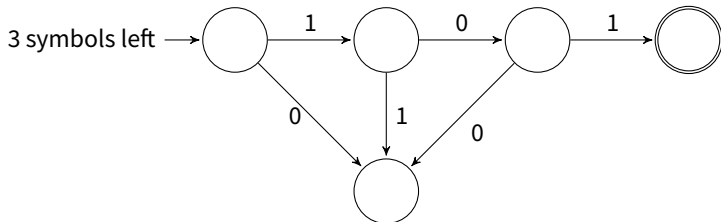
We'll answer this question in the next few lectures

Useful to consider hypothetical machines that are **nondeterministic**

Even easier with guesses

Suppose we could **guess** when the input string has only 3 symbols left

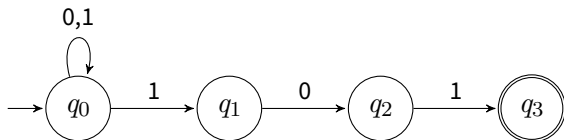
Accept strings ending in 101:



This is **not** a DFA!

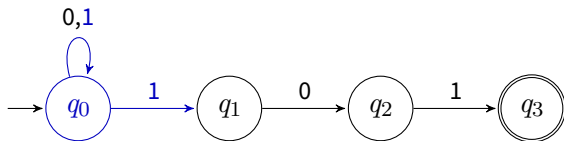
Nondeterministic finite automata

A machine that allows us to make **guesses**



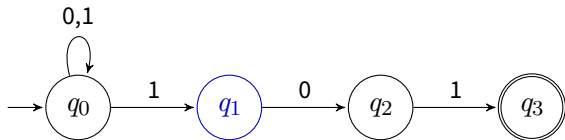
Each state can have **zero, one, or more** outgoing transitions labeled by the same symbol

Choosing where to go



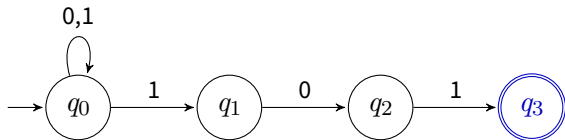
State q_0 has two transitions labeled 1
Upon reading 1, we have the **choice** of staying at q_0 or moving to q_1

Ability to choose



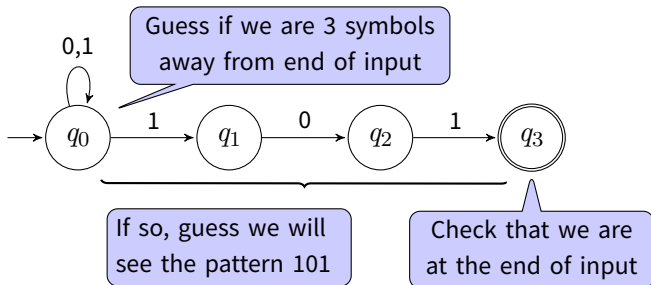
State q_1 has **no transition labeled 1**
Upon reading 1 at q_1 , die; upon reading 0, continue to q_2

Ability to choose

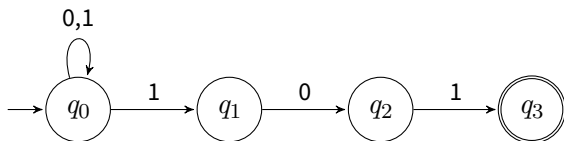


State q_1 has **no transition** going out
Upon reading 0 or 1 at q_3 , die

Meaning of NFA



How to run an NFA



input: 01101

The NFA can have **several active states** at the same time
NFA accepts if at the end, **one of its active states** is accepting

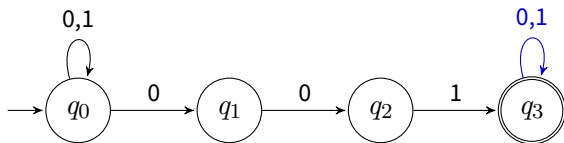
Example

Construct an NFA over alphabet $\{0, 1\}$ that accepts all strings containing the pattern 001 somewhere

11001010, 001001, 111001 should be accepted
 ϵ , 000, 010101 should not

Example

Construct an NFA over alphabet $\{0, 1\}$ that accepts all strings containing the pattern 001 somewhere



Definition

A **nondeterministic finite automaton** (NFA) is a 5-tuple $(Q, \Sigma, \delta, q_0, F)$
where

- ▶ Q is a finite set of states
- ▶ Σ is an alphabet
- ▶ $\delta : Q \times (\Sigma \cup \{\epsilon\}) \rightarrow$ **subsets of Q** is a transition function
- ▶ $q_0 \in Q$ is the initial state
- ▶ $F \subseteq Q$ is a set of accepting states

Differences from DFA:

- ▶ transition function δ can go into several states
- ▶ allows **ϵ -transitions**

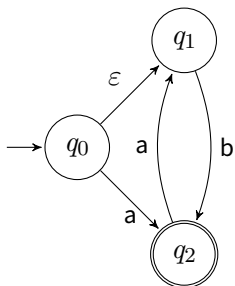
Language of an NFA

The NFA **accepts** string x if there is some path that, starting from q_0 , ends at an accepting state as x is read from left to right

The **language of an NFA** is the set of all strings accepted by the NFA

ϵ -transitions

ϵ -transitions can be taken **for free**:



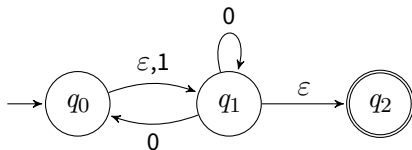
accepts

a, b, aab, bab, aabab, ...

rejects

ϵ , aa, ba, bb, ...

Example



alphabet $\Sigma = \{0, 1\}$

states $Q = \{q_0, q_1, q_2\}$

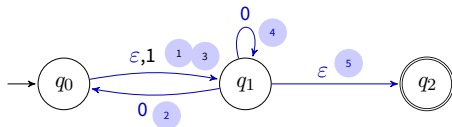
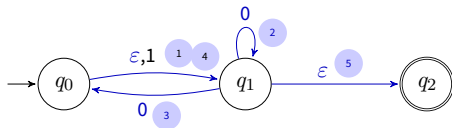
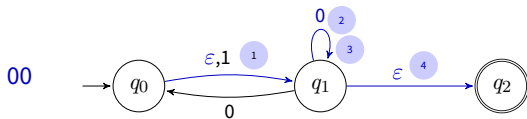
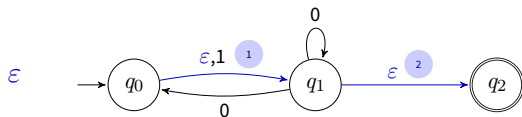
initial state q_0

accepting states $F = \{q_2\}$

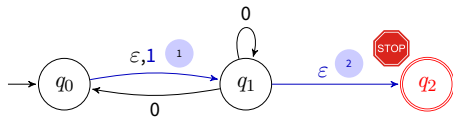
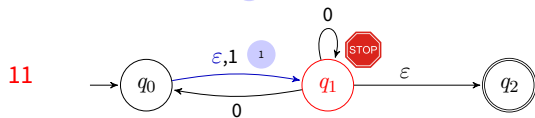
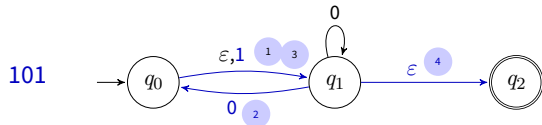
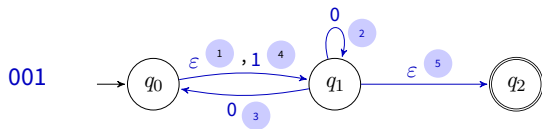
table of transition function δ

		inputs		
		0	1	ϵ
states	q_0	\emptyset	$\{q_1\}$	$\{q_1\}$
	q_1	$\{q_0, q_1\}$	\emptyset	$\{q_2\}$
	q_2	\emptyset	\emptyset	\emptyset

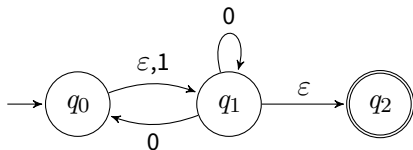
Running NFA



Running NFA



Language of this NFA



What is the language of this NFA?

Example of ϵ -transitions

Construct an NFA that accepts all strings with an even number of 0s **or** an odd number of 1s

Example of ϵ -transitions

Construct an NFA that accepts all strings with an even number of 0s **or** an odd number of 1s

