



香港中文大學
The Chinese University of Hong Kong

7.5 A Discussion

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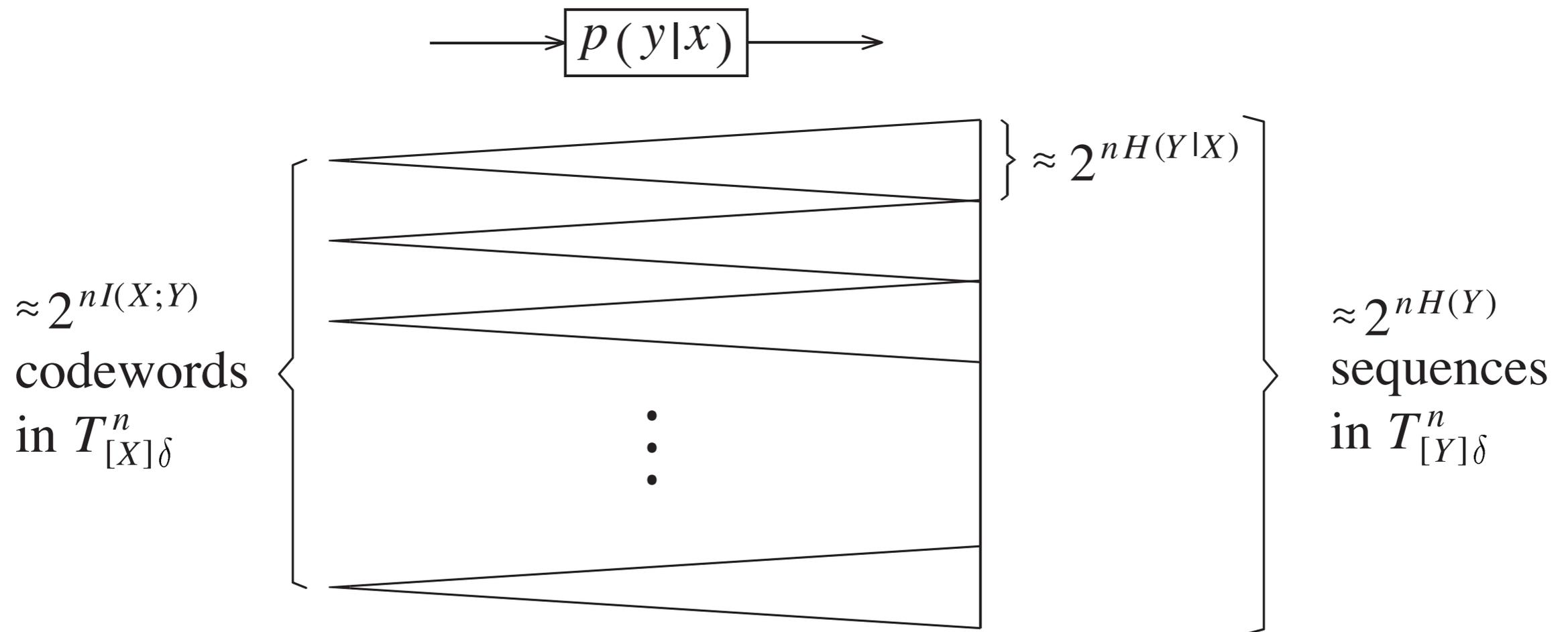
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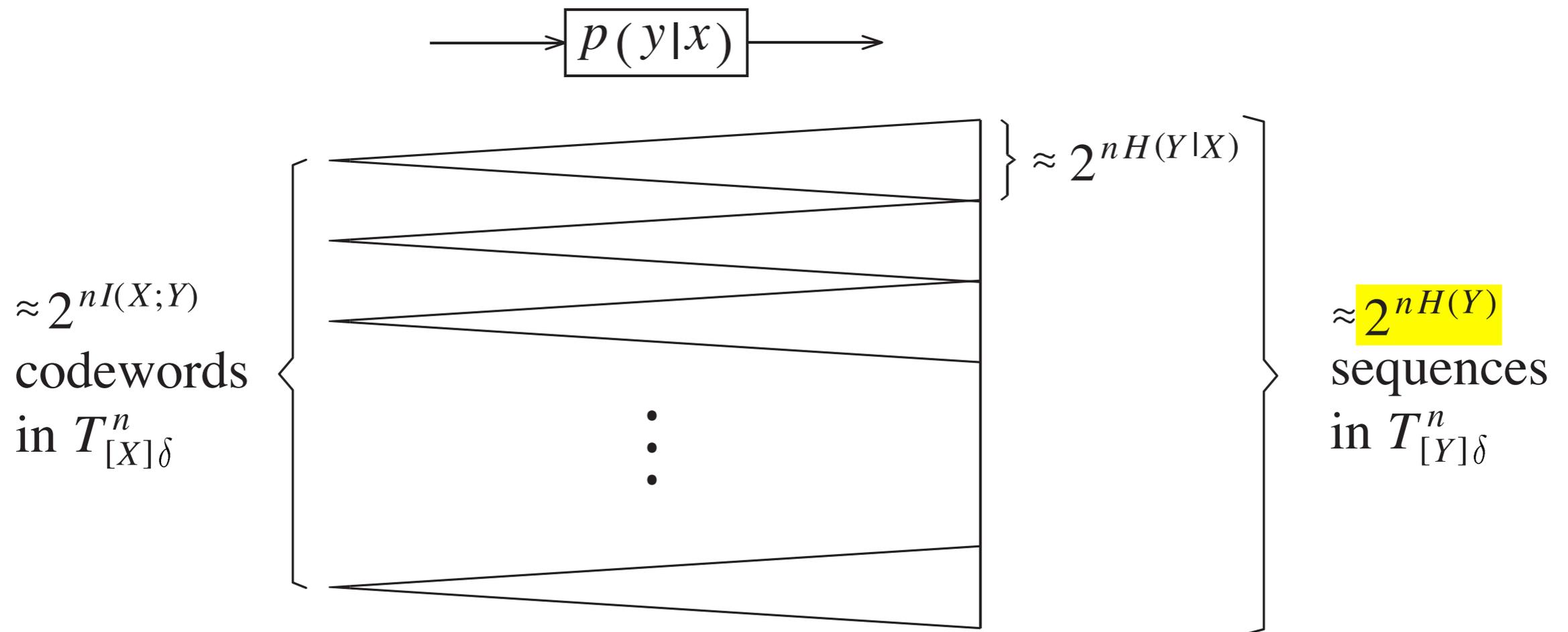
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- In particular, the repetition code is not a good code because the numbers of ‘0’ and ‘1’s in the codewords are not roughly the same.



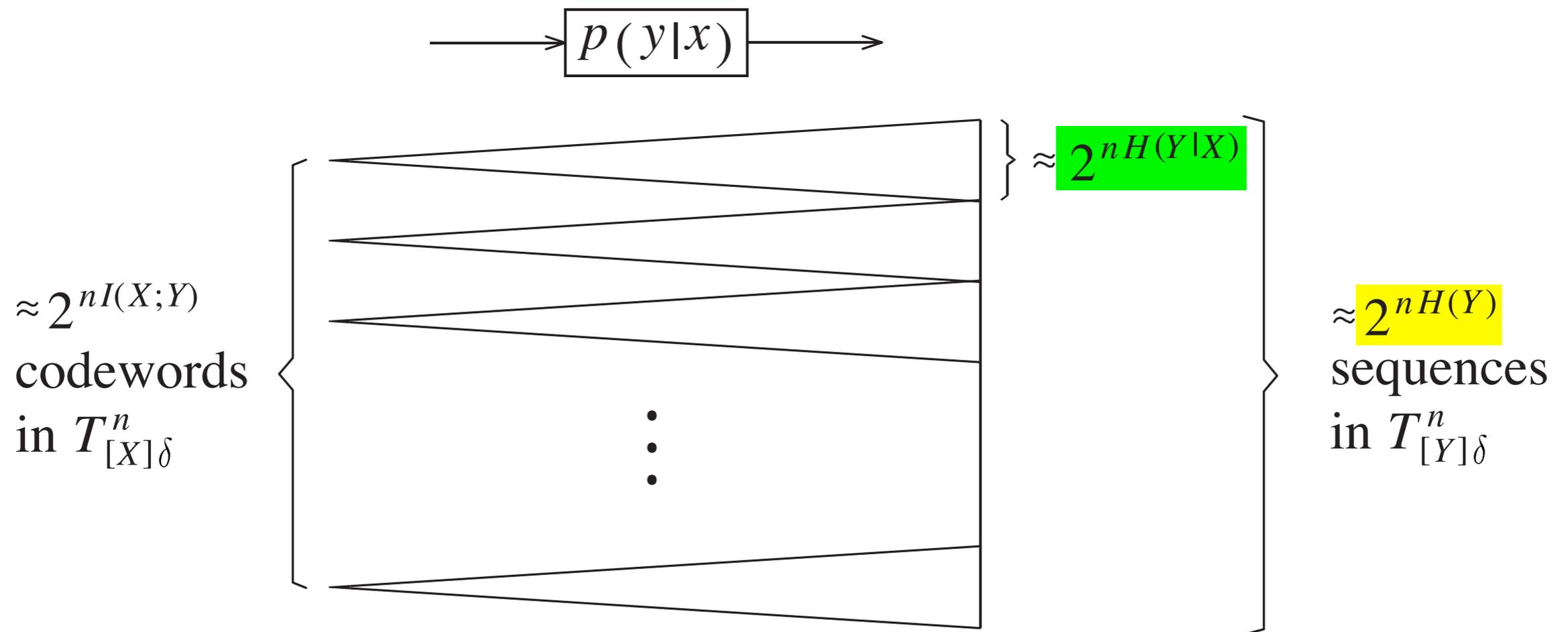
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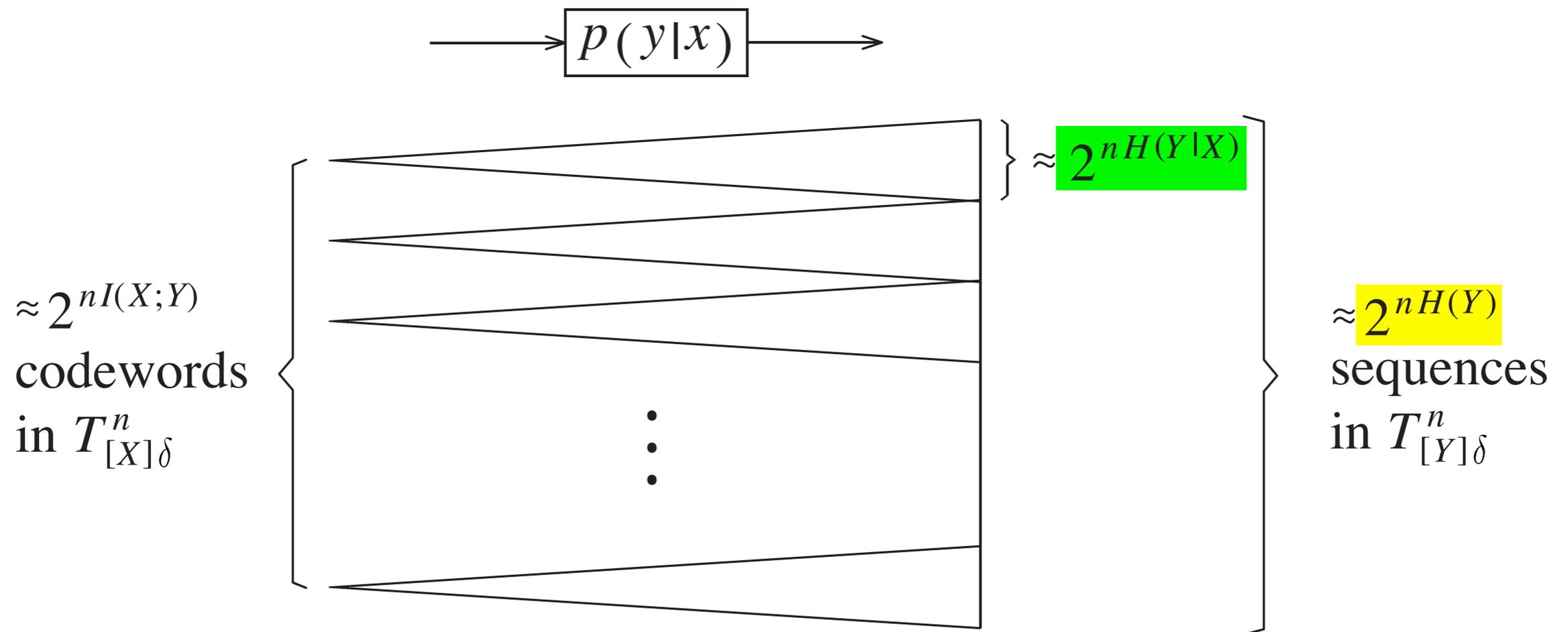
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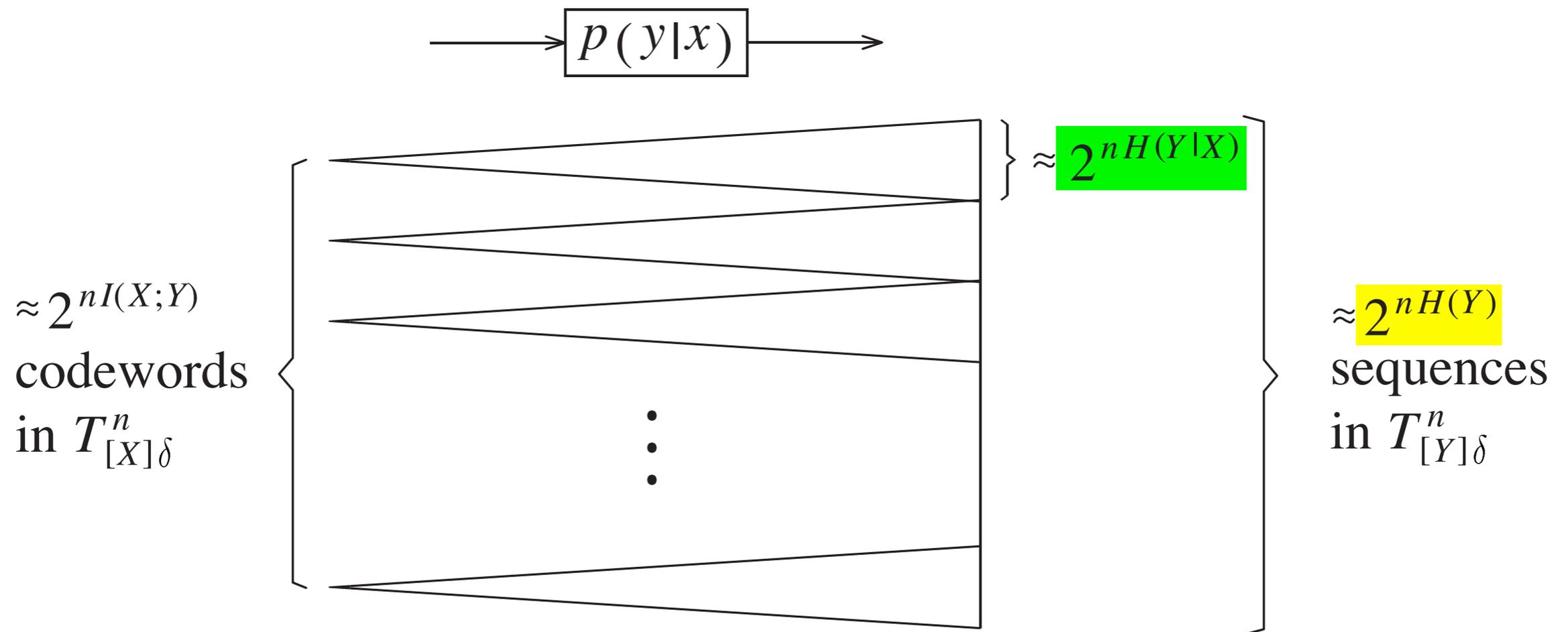
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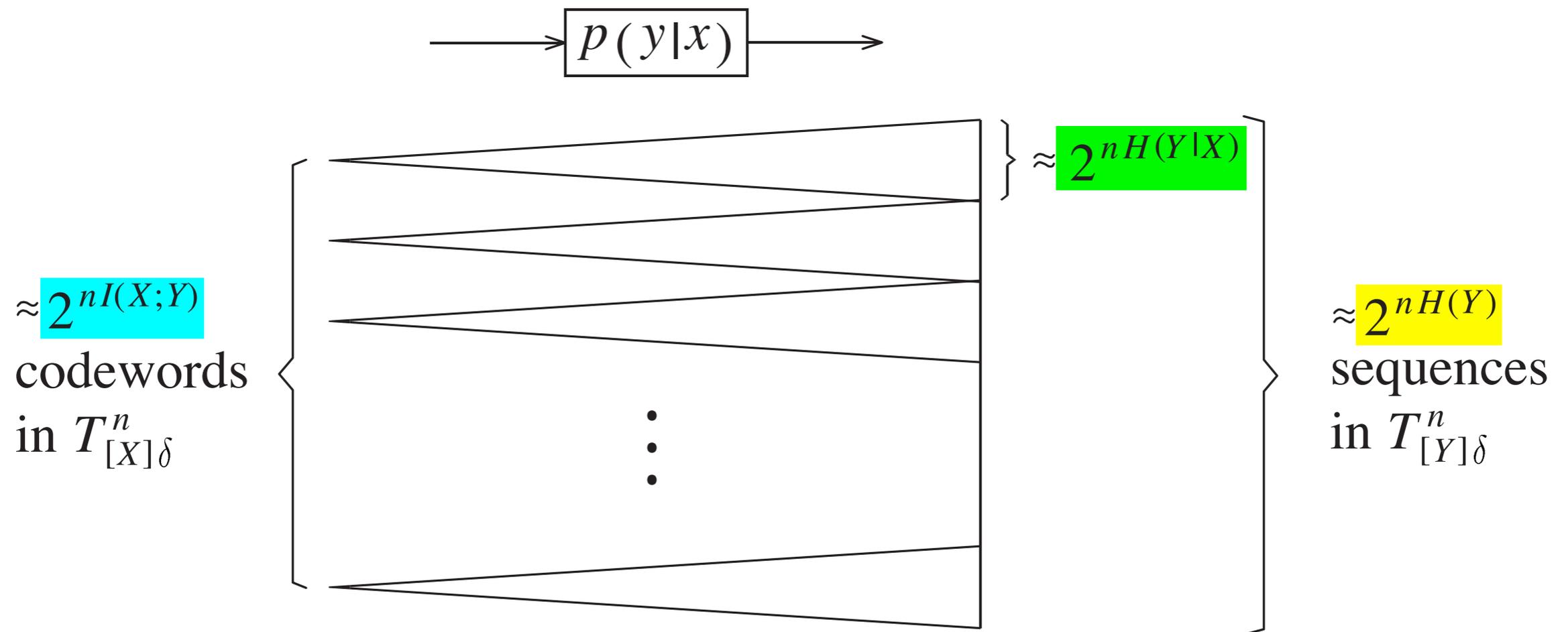
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- In particular, turbo code and LDPC code are almost “capacity achieving.”