

Chapter 6

Strong Typicality

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6.1 Strong AEP

Setup

- $\{X_k, k \geq 1\}$, X_k i.i.d. $\sim p(x)$.
- X denotes generic r.v. with entropy $H(X) < \infty$.
- $|\mathcal{X}| < \infty$

6.1 Strong AEP

Definition 6.1 The strongly typical set $T_{[X]\delta}^n$ with respect to $p(x)$ is the set of sequences $\mathbf{x} = (x_1, x_2, \dots, x_n) \in \mathcal{X}^n$ such that $N(x; \mathbf{x}) = 0$ for $x \notin \mathcal{S}_X$, and

$$\sum_x \left| \frac{1}{n} N(x; \mathbf{x}) - p(x) \right| \leq \delta,$$

where $N(x; \mathbf{x})$ is the number of occurrences of x in the sequence \mathbf{x} , and δ is an arbitrarily small positive real number. The sequences in $T_{[X]\delta}^n$ are called strongly δ -typical sequences.

Theorem 6.2 (Strong AEP) There exists $\eta > 0$ such that $\eta \rightarrow 0$ as $\delta \rightarrow 0$, and the following hold:

1) If $\mathbf{x} \in T_{[X]\delta}^n$, then

$$2^{-n(H(X)+\eta)} \leq p(\mathbf{x}) \leq 2^{-n(H(X)-\eta)}.$$

2) For n sufficiently large,

$$\Pr\{\mathbf{X} \in T_{[X]\delta}^n\} > 1 - \delta.$$

3) For n sufficiently large,

$$(1 - \delta)2^{n(H(X)-\eta)} \leq |T_{[X]\delta}^n| \leq 2^{n(H(X)+\eta)}.$$

Proof

1. If the relative frequency is about right, then everything else, including the empirical entropy, would be about right.
2. A consequence of WLLN.
3. Exactly the same as the proof of Part 3) of Theorem 5.3.

Theorem 6.3 For sufficiently large n , there exists $\varphi(\delta) > 0$ such that

$$\Pr\{\mathbf{X} \notin T_{[X]\delta}^n\} < 2^{-n\varphi(\delta)}.$$

Proof Chernoff bound.

6.2 Strong Typicality vs Weak Typicality

- Weak typicality: empirical entropy $\approx H(X)$
- Strong typicality: relative frequency $\sim p(x)$
- Strong typicality \Rightarrow weak typicality (Proposition 6.5)
- Weak typicality $\not\Rightarrow$ strong typicality (see example in text)
- Both have AEP, but strong typicality has stronger conditional asymptotic properties (Theorem 6.10).
- Strong typicality works only for finite alphabet, i.e., $|\mathcal{X}| < \infty$.

Strong Typicality Implies Weak Typicality

Proposition 6.5 For any $\mathbf{x} \in \mathcal{X}^n$, if $\mathbf{x} \in T_{[X]\delta}^n$, then $\mathbf{x} \in W_{[X]\eta}^n$, where $\eta \rightarrow 0$ as $\delta \rightarrow 0$.

Proof By strong AEP and the definitions.

6.3 Joint Typicality

Setup

- $\{(X_k, Y_k), k \geq 1\}$, (X_k, Y_k) i.i.d. $\sim p(x, y)$.
- (X, Y) denotes pair of generic r.v. with entropy $H(X, Y) < \infty$.
- $|\mathcal{X}|, |\mathcal{Y}| < \infty$

Definition 6.6 The strongly jointly typical set $T_{[XY]\delta}^n$ with respect to $p(x, y)$ is the set of $(\mathbf{x}, \mathbf{y}) \in \mathcal{X}^n \times \mathcal{Y}^n$ such that $N(x, y; \mathbf{x}, \mathbf{y}) = 0$ for $(x, y) \notin \mathcal{S}_{XY}$, and

$$\sum_x \sum_y \left| \frac{1}{n} N(x, y; \mathbf{x}, \mathbf{y}) - p(x, y) \right| \leq \delta,$$

where $N(x, y; \mathbf{x}, \mathbf{y})$ is the number of occurrences of (x, y) in the pair of sequences (\mathbf{x}, \mathbf{y}) , and δ is an arbitrarily small positive real number. A pair of sequences (\mathbf{x}, \mathbf{y}) is called strongly jointly δ -typical if it is in $T_{[XY]\delta}^n$.

Theorem 6.7 (Consistency) If $(\mathbf{x}, \mathbf{y}) \in T_{[XY]\delta}^n$, then $\mathbf{x} \in T_{[X]\delta}^n$ and $\mathbf{y} \in T_{[Y]\delta}^n$.

Theorem 6.8 (Preservation) Let $Y = f(X)$. If

$$\mathbf{x} = (x_1, x_2, \dots, x_n) \in T_{[X]\delta}^n,$$

then

$$f(\mathbf{x}) = (y_1, y_2, \dots, y_n) \in T_{[Y]\delta}^n,$$

where $y_i = f(x_i)$ for $1 \leq i \leq n$.

Theorem 6.9 (Strong JAEP) Let

$$(\mathbf{X}, \mathbf{Y}) = ((X_1, Y_1), (X_2, Y_2), \dots, (X_n, Y_n)),$$

where (X_i, Y_i) are i.i.d. with generic pair of random variables (X, Y) . Then there exists $\lambda > 0$ such that $\lambda \rightarrow 0$ as $\delta \rightarrow 0$, and the following hold:

1) If $(\mathbf{x}, \mathbf{y}) \in T_{[XY]\delta}^n$, then

$$2^{-n(H(X,Y)+\lambda)} \leq p(\mathbf{x}, \mathbf{y}) \leq 2^{-n(H(X,Y)-\lambda)}.$$

2) For n sufficiently large,

$$\Pr\{(\mathbf{X}, \mathbf{Y}) \in T_{[XY]\delta}^n\} > 1 - \delta.$$

3) For n sufficiently large,

$$(1 - \delta)2^{n(H(X,Y)-\lambda)} \leq |T_{[XY]\delta}^n| \leq 2^{n(H(X,Y)+\lambda)}.$$

Stirling's Approximation

Lemma 6.11 (simplified) $\ln n! \sim n \ln n$.

Proof Write

$$\ln n! = \ln 1 + \ln 2 + \cdots + \ln n.$$

Since $\ln x$ is a monotonically increasing, we have

$$\int_{k-1}^k \ln x \, dx < \ln k < \int_k^{k+1} \ln x \, dx.$$

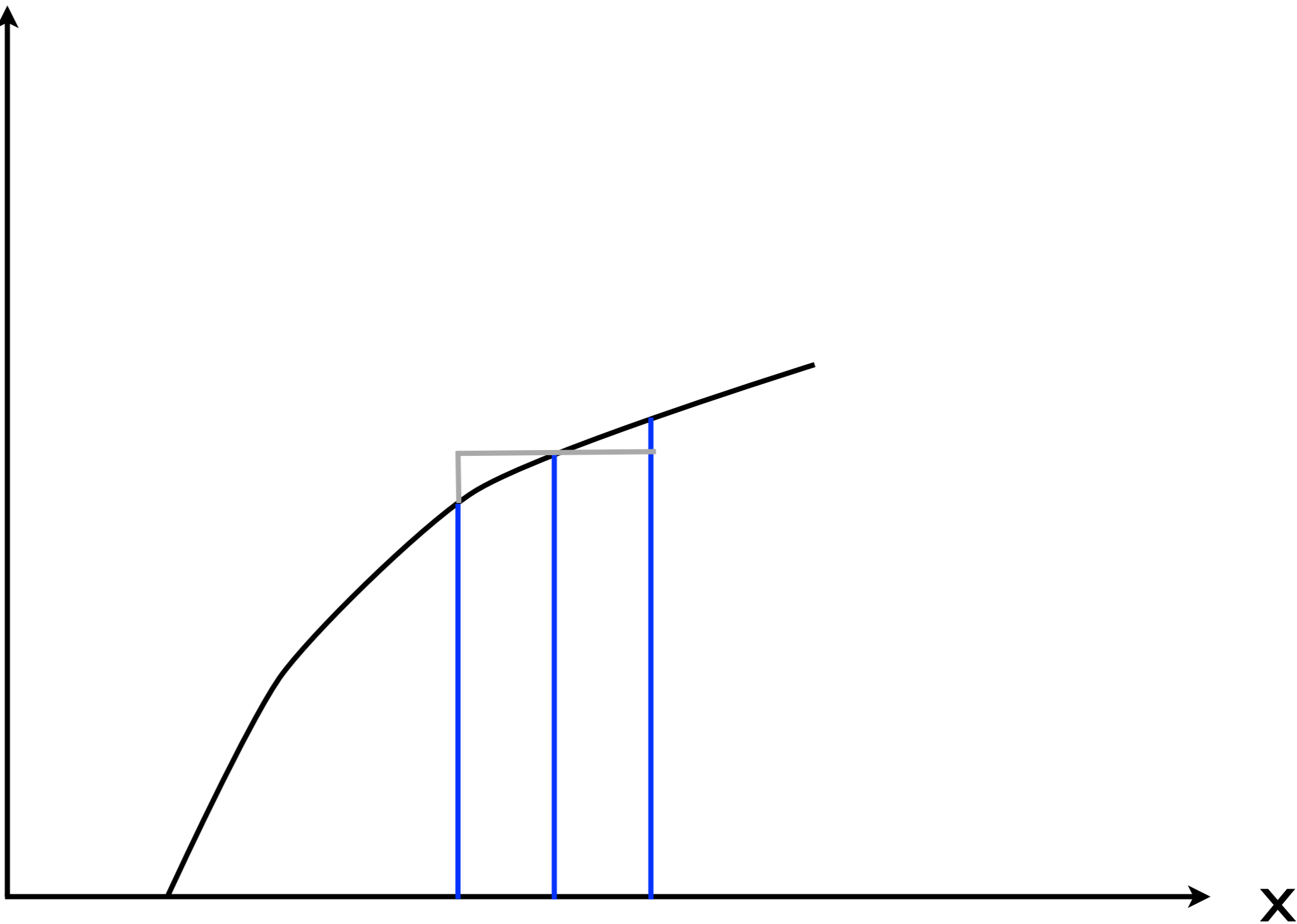
Summing over $1 \leq k \leq n$, we have

$$\int_0^n \ln x \, dx < \ln n! < \int_1^{n+1} \ln x \, dx,$$

or

$$n \ln n - n < \ln n! < (n + 1) \ln(n + 1) - n.$$

$\ln x$



$k-1$

k

$k+1$

x

An Application

Lemma For large n ,

$$\binom{n}{np, n(1-p)} \approx 2^{nH_2(\{p, 1-p\})}$$

Proof

$$\begin{aligned} \ln \binom{n}{np, n(1-p)} &\approx n \ln n - (np) \ln(np) - (n(1-p)) \ln(n(1-p)) \\ &= n \ln n - np[\ln n + \ln p] - n(1-p)[\ln n + \ln(1-p)] \\ &= -n[p \ln p + (1-p) \ln(1-p)] \end{aligned}$$

Changing to the base 2, we have

$$\log_2 \binom{n}{np, n(1-p)} \approx nH_2(\{p, 1-p\})$$

In general, for large n ,

$$\binom{n}{np_1, np_2, \dots, np_m} = \frac{n!}{\prod_i (np_i)!} \approx 2^{nH(\{p_i\})}$$

Theorem 6.10 (Conditional Strong AEP) For any $\mathbf{x} \in T_{[X]\delta}^n$, define

$$T_{[Y|X]\delta}^n(\mathbf{x}) = \{\mathbf{y} \in T_{[Y]\delta}^n : (\mathbf{x}, \mathbf{y}) \in T_{[XY]\delta}^n\}.$$

If $|T_{[Y|X]\delta}^n(\mathbf{x})| \geq 1$, then

$$2^{n(H(Y|X)-\nu)} \leq |T_{[Y|X]\delta}^n(\mathbf{x})| \leq 2^{n(H(Y|X)+\nu)},$$

where $\nu \rightarrow 0$ as $n \rightarrow \infty$ and $\delta \rightarrow 0$.

Remark Weak Typicality guarantees that the number of \mathbf{y} that are jointly typical with a typical \mathbf{x} is approximately equal to $2^{n(H(Y|X))}$ on the average. Strong typicality guarantees that this is so for each typical \mathbf{x} as long as there exists at least one \mathbf{y} that is jointly typical with \mathbf{x} .

Upper Bound in Theorem 6.10

- For any $\nu > 0$, consider

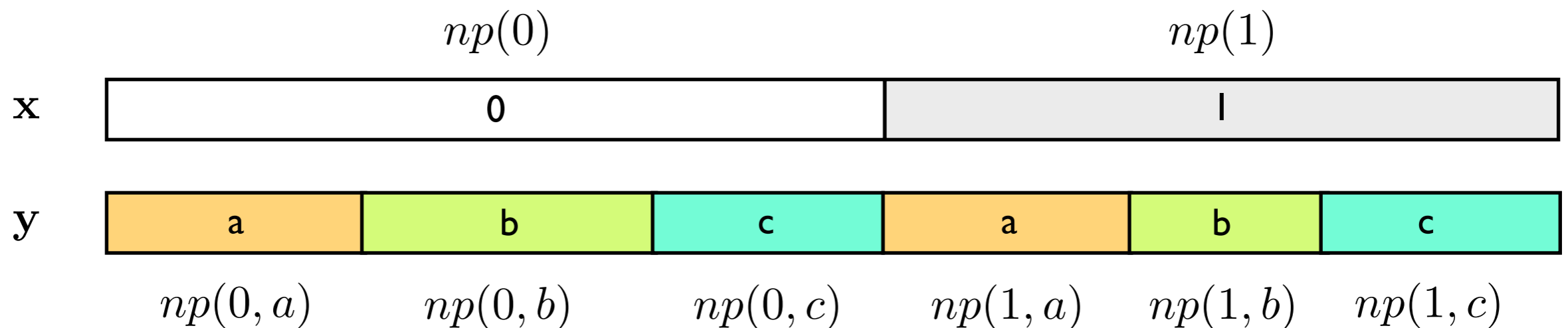
$$\begin{aligned} 2^{-n(H(X)-\nu/2)} &\stackrel{a)}{\geq} p(\mathbf{x}) \\ &= \sum_{\mathbf{y} \in \mathcal{Y}^n} p(\mathbf{x}, \mathbf{y}) \\ &\geq \sum_{\mathbf{y} \in T_{[Y|X]\delta}^n(\mathbf{x})} p(\mathbf{x}, \mathbf{y}) \\ &\stackrel{b)}{\geq} \sum_{\mathbf{y} \in T_{[Y|X]\delta}^n(\mathbf{x})} 2^{-n(H(X,Y)+\nu/2)} \\ &= |T_{[Y|X]\delta}^n(\mathbf{x})| 2^{-n(H(X,Y)+\nu/2)}, \end{aligned}$$

a) and b) follows from strong joint AEP.

- Similar to the proof of the upper bound on $|T_{[X]\delta}^n|$ in Theorem 6.2 (SAEP).

Lower Bound in Theorem 6.10

$$\mathcal{X} = \{0, 1\}, \mathcal{Y} = \{a, b, c\}$$



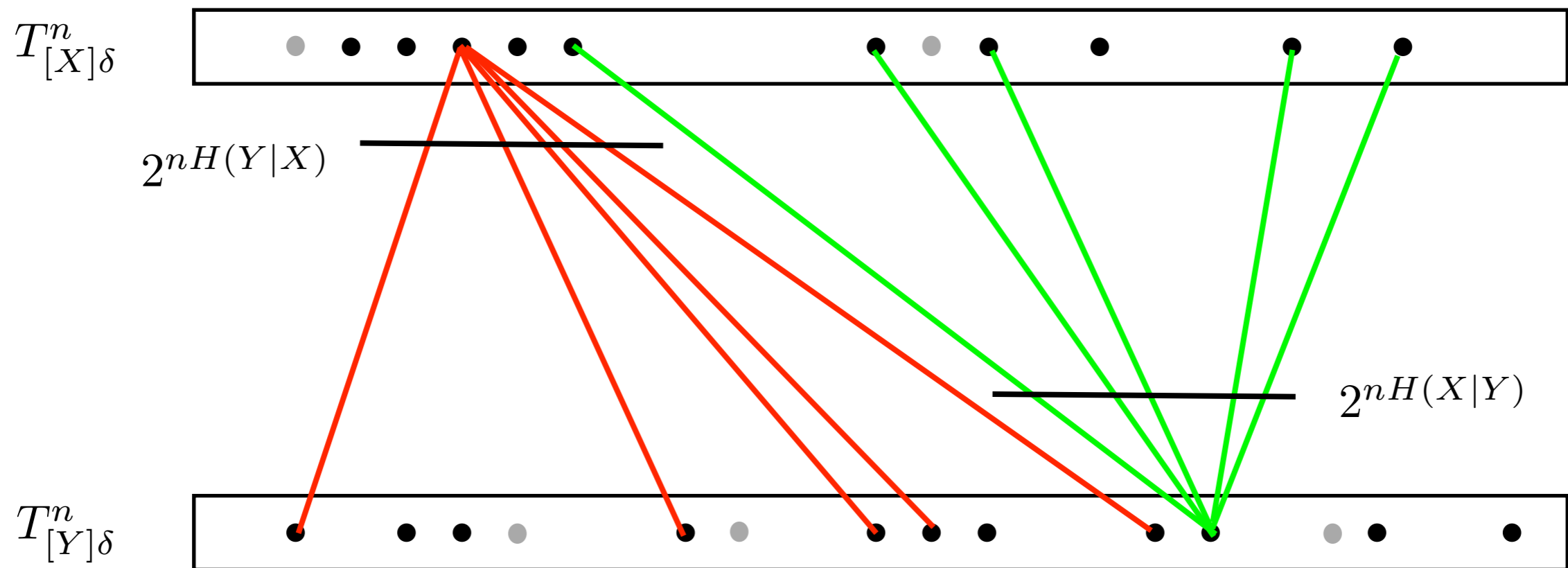
Rearrange the components of \mathbf{y} corresponding to $x_k = 0$ and rearrange the components of \mathbf{y} corresponding to $x_k = 1$. This preserves joint typicality.

$$\begin{aligned} \#arrangements &\approx \binom{np(0)}{np(0, a), np(0, b), np(0, c)} \binom{np(1)}{np(1, a), np(1, b), np(1, c)} \\ &\approx 2^{np(0)H(\{p(\cdot|0)\})} 2^{np(1)H(\{p(\cdot|1)\})} \\ &= 2^{n(p(0)H(Y|X=0) + p(1)H(Y|X=1))} \\ &= 2^{nH(Y|X)} \end{aligned}$$

Hence,

$$|T_{[Y|X]\delta}^n(\mathbf{x})| \geq 2^{n(H(Y|X) - \nu)}.$$

An Illustration of Conditional SAEP



Corollary 6.12 For a joint distribution $p(x, y)$ on $\mathcal{X} \times \mathcal{Y}$, let $S_{[X]_\delta}^n$ be the set of all sequences $\mathbf{x} \in T_{[X]_\delta}^n$ such that $T_{[Y|X]_\delta}^n(\mathbf{x})$ is nonempty. Then

$$|S_{[X]_\delta}^n| \geq (1 - \delta)2^{n(H(X) - \psi)},$$

where $\psi \rightarrow 0$ as $n \rightarrow \infty$ and $\delta \rightarrow 0$.

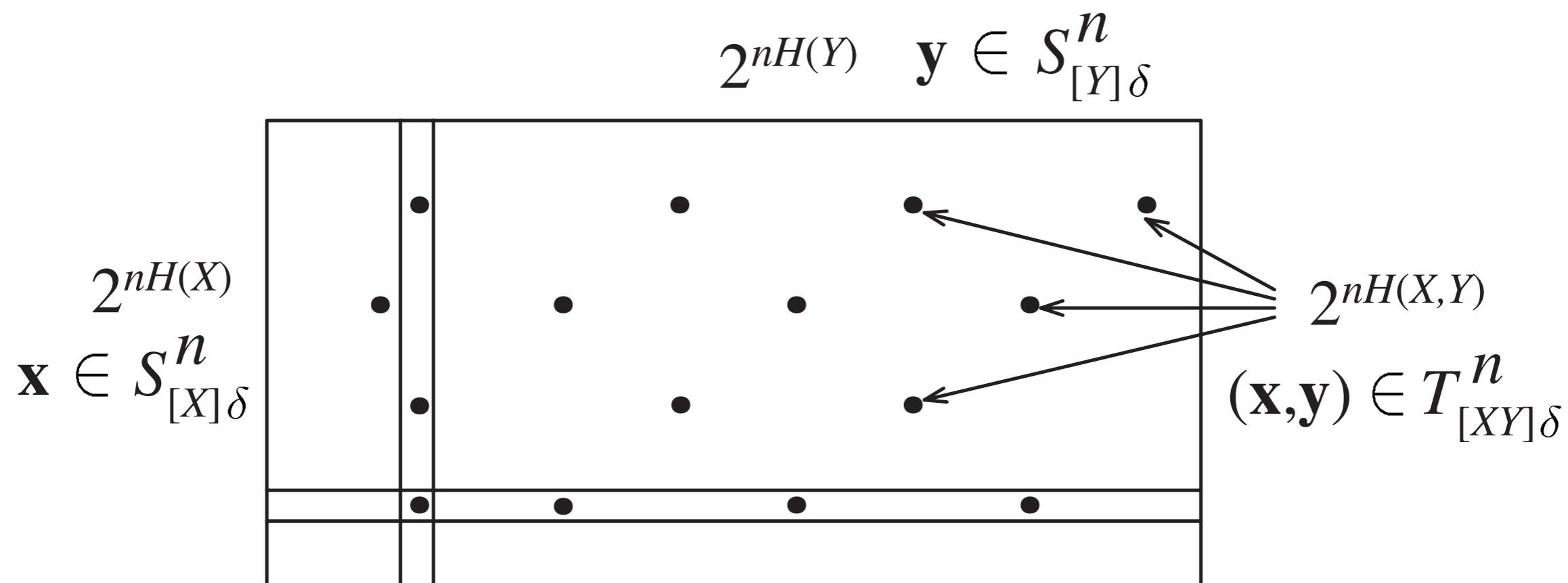
Proposition 6.13 With respect to a joint distribution $p(x, y)$ on $\mathcal{X} \times \mathcal{Y}$, for any $\delta > 0$,

$$\Pr\{\mathbf{X} \in S_{[X]_\delta}^n\} > 1 - \delta$$

for n sufficiently large.

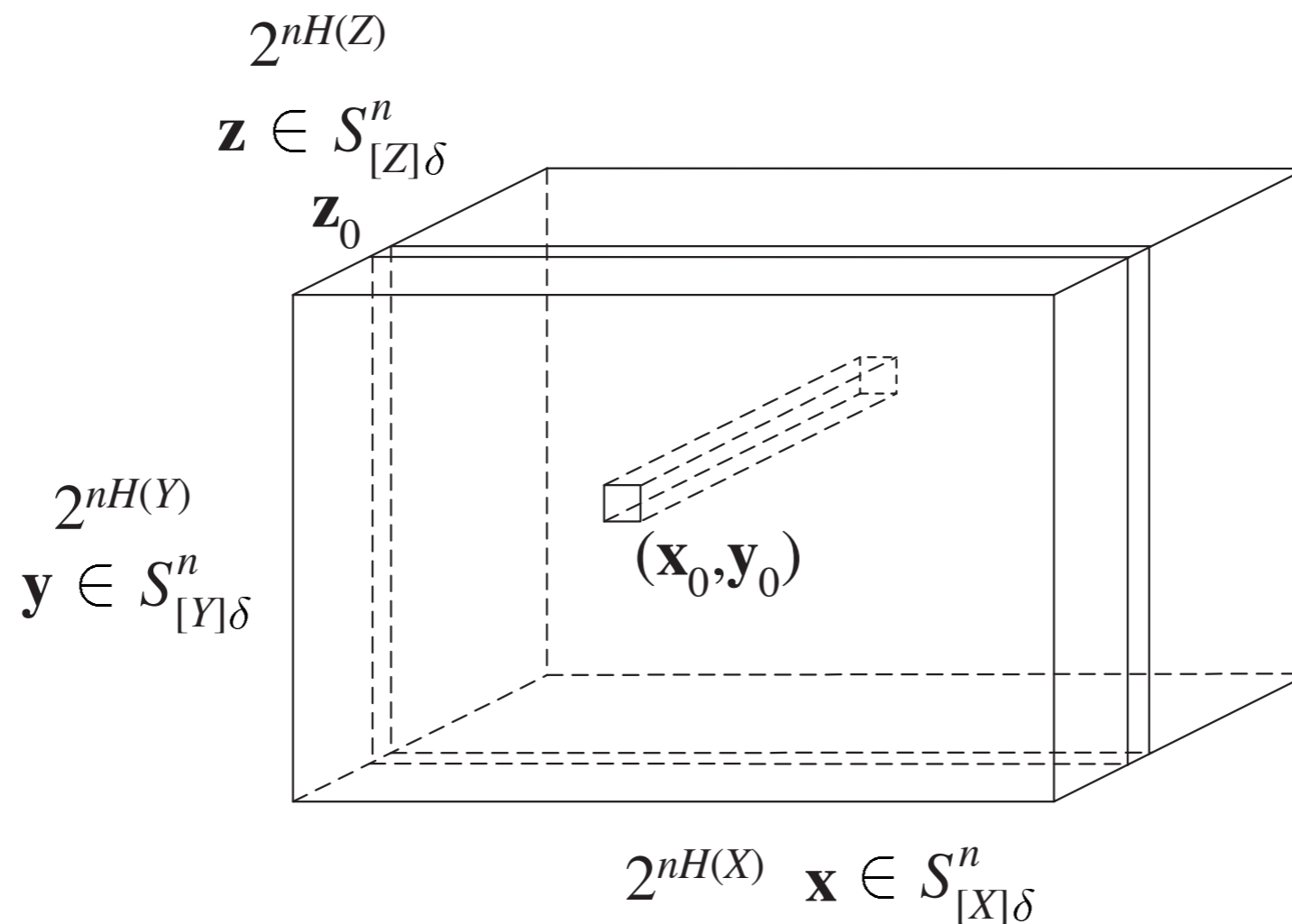
Strongly Joint Typicality Array

- Exhibits an “asymptotic quasi-uniform” structure.
- Two-Dimensional:



Strongly Joint Typicality Array

- Three-Dimensional:



Quasi-Uniform Array

- Provides a combinatorial interpretation of information inequalities.
- Related to many branches of information sciences: combinatorics, group theory (Ch. 16), Kolmogorov complexity, network coding, probability theory, matrix theory, quantum mechanics, ...
- Resources:
 - http://temple.birs.ca/~09w5103/yeung_09w5103_talk.pdf
(slides)
 - <http://www.birs.ca/events/2009/5-day-workshops/09w5103/videos>
(video)