

# More on Binary Search Trees

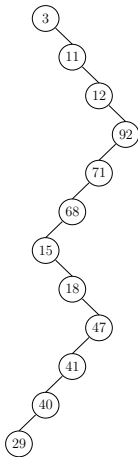
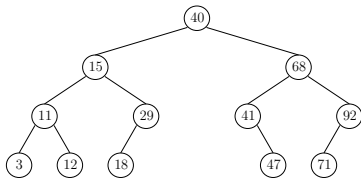
CSCI2100 Tutorial 10

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## Binary Search Tree Example

Two possible BSTs on  $S = \{3, 11, 12, 15, 18, 29, 40, 41, 47, 68, 71, 92\}$ :

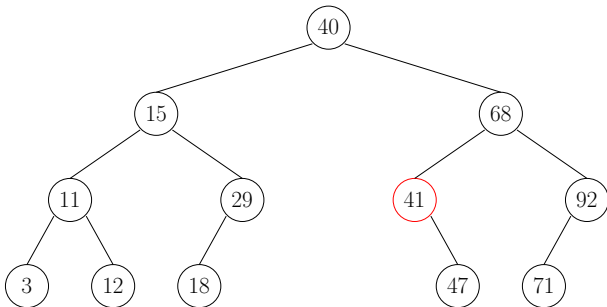


## Predecessor Query

Let  $S$  be a set of integers. A predecessor query for a given integer  $q$  is to find its **predecessor** in  $S$ , which is the largest integer in  $S$  that does not exceed  $q$ .

### Example

Suppose that  $S = \{3, 11, 12, 15, 18, 29, 40, 41, 47, 68, 71, 92\}$  and we have a balanced BST  $T$  on  $S$ :



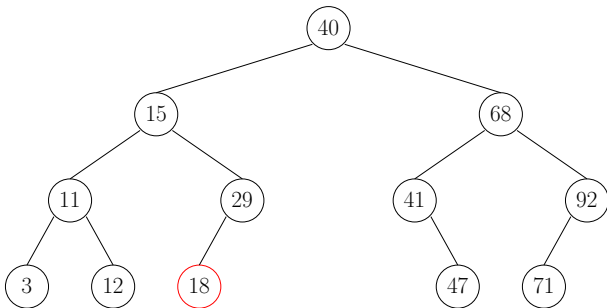
We want to find the predecessor of  $q = 42$  in  $S$ .

## Successor Query

Let  $S$  be a set of integers. A successor query for a given integer  $q$  is to find its **successor** in  $S$ , which is the smallest integer in  $S$  that is no smaller than  $q$ .

## Example

We want to find the successor of  $q = 17$  in  $S$ .



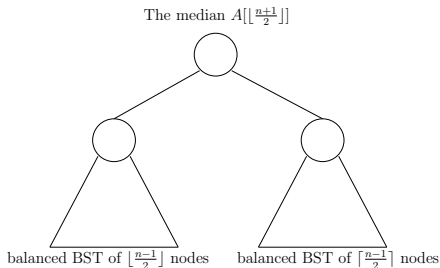
## Construction of a Balanced BST

In the following, we will discuss how to construct a balanced BST  $T$  on a given **sorted** set  $S$  of  $n$  integers in  $O(n)$  time.

## Construction of a Balanced BST

Assume that  $S$  is stored an array  $A$  and  $A$  is sorted.

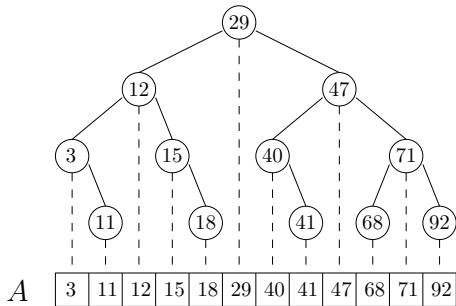
- **Observation:** The subtree of **any** node in a balanced BST is also a balanced BST.
- **Main idea:** A BST of  $n$  nodes constructed by the following form:





### Example

Let us construct a balanced BST  $T$  on the following sorted array  $A$ .



## Construction of a Balanced BST

Let  $f(n)$  be the maximum running time for constructing a balanced BST from an array of length  $n$ . We have:

$$f(1) = O(1)$$

$$f(n) = O(1) + 2 \cdot f(\lceil n/2 \rceil)$$

Solving the recurrence gives  $f(n) = O(n)$ .

## Range Count Problem

Let  $S$  be a set of  $n$  integers. Given two integers  $a$  and  $b$  such that  $a \leq b$ . Find the **number** of integers in  $S$  which are in the range of  $[a, b]$ .

In the following, we will discuss how to **augment** a balanced BST on  $S$  to achieve:

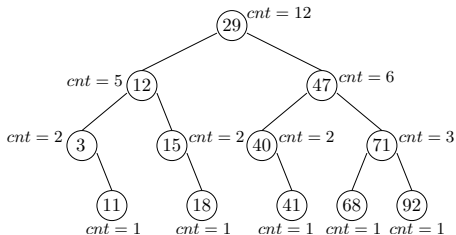
- $O(n)$  space consumption,
- $O(\log n)$  time for each query.

## Range Count Problem

Augment a balanced BST  $T$  on  $S$  by storing one additional information in each node  $u$  that is:

- the number of nodes in the subtree of  $u$ .

For example,

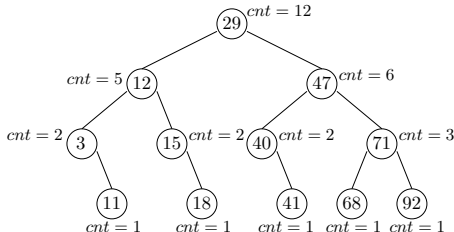


## Range Count Problem

Define a concept first.

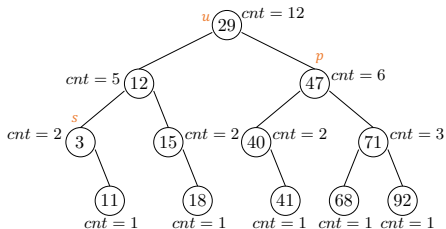
- **Lowest Common Ancestor:** Let  $t$  be the root. The lowest common ancestor of nodes  $v_1$  and  $v_2$  is the **lowest node** that is on both of the paths  $P(t, v_1)$  and  $P(t, v_2)$ .

For example, the lowest common ancestor of node with key 3 and node with key 15 is the node with key 12.



## Range Count Problem

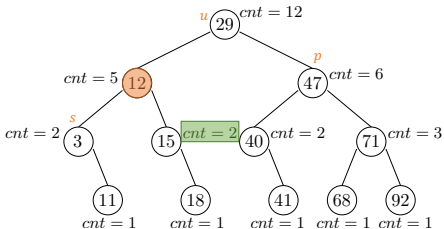
For a range  $[2, 48]$ , let  $s$  be the successor of 2,  $p$  the predecessor of 48 and  $u$  the **lowest common ancestor** of  $s$  and  $p$ .  
Initialize a count  $c = 1$  (since  $u$  is within the range)



## Range Count Problem

Traverse the path from  $u$ 's left child to  $s$ .  
For every node  $v$  being visited, if  $v.\text{key} \geq 2$ :

- $c += 1$
- $c +=$  the counter of  $v$ 's right child

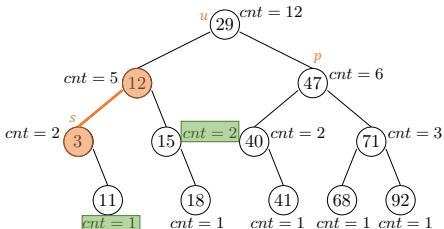


$C$  is incremented by  $1 + 2$ .

## Range Count Problem

Traverse the path from  $u$ 's left child to  $s$ .  
For every node  $v$  being visited, if  $v.\text{key} \geq 2$ :

- $c += 1$
- $c +=$  the counter of  $v$ 's right child



$C$  is incremented by  $1 + 1$ .

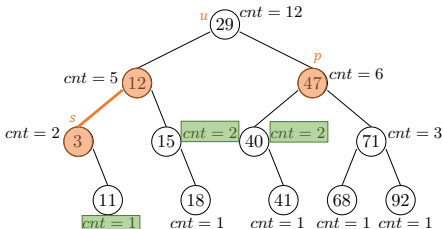


## Range Count Problem

Traverse the path from  $u$ 's right child to  $p$ .

For every node  $v$  being visited, if  $v.\text{key} \leq 48$ :

- $c += 1$
- $c +=$  the counter of  $v$ 's left child

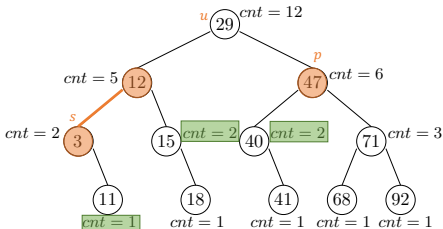


$C$  is incremented by  $1 + 2$ . Finally,  $c$  becomes 9.

## Range Count Problem

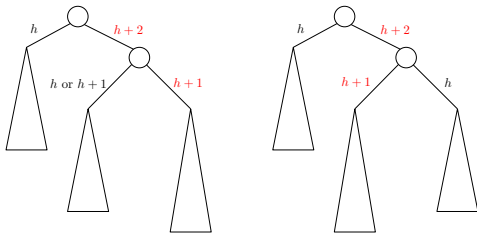
We walked through two paths, at most  $\log_2 n$  nodes in each path.  
For each node visited, we perform constant-time operations, which takes  $O(1)$ .

Time complexity:  $O(\log n)$



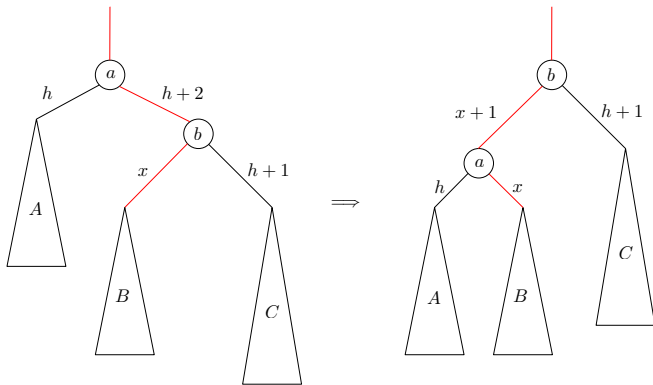
## Rebalancing

In lectures we explored the Left-Left and Left-Right cases in detail, so here we will look at Right-Right and Right-Left:



## Right-Right

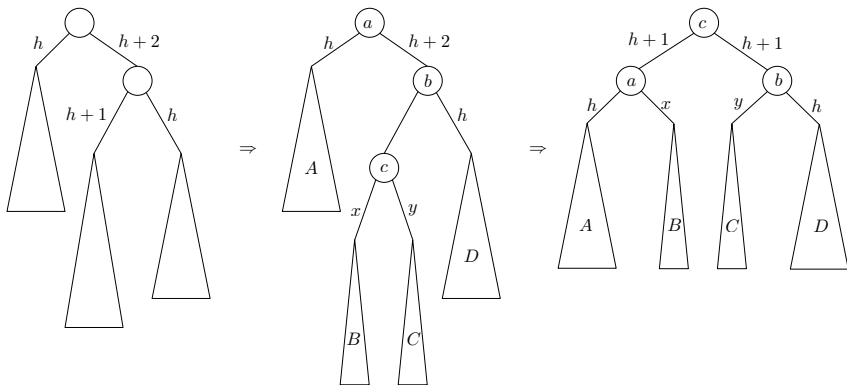
Fix by a **rotation** (symmetric to left-left):



Note that  $x = h$  or  $h + 1$ , and the ordering from left to right of  $A, a, B, b, C$  is preserved after rotation.

## Right-Left

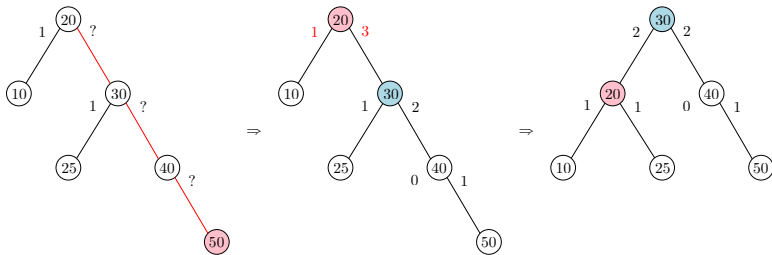
Fix by a **double rotation** (symmetric to left-right):



Note that  $x$  and  $y$  must be  $h$  or  $h - 1$ . Furthermore at least one of them must be  $h$ .

## Right-Right Example

Inserting 50:



## Right-Left Example

Inserting 38:

