## Solutions of MATH5360 Assignment 2



1. (a)Set up the tableau and apply pivoting operations, we have

Thus an optimal vector for the primal problem is  $(y_1, y_2, y_3) = (1, 6, 0)$  and the maximum value  $f$  is 45.

The dual problem is

$$
\begin{array}{ll}\n\text{min} & g = 15x_1 + 24x_2 + 12 \\
\text{subject to} & 3x_1 \ge 3 \\
& 2x_1 + 4x_2 \ge 5 \\
& 2x_1 + 5x_2 \ge 4\n\end{array}
$$

Thus an optimal vector for the primal problem is  $(x_1, x_2) = (1, 3/4)$  and the minimum value  $g$  is 45. (b)Set up the tableau and apply pivoting operations, we have





Thus an optimal vector for the primal problem is  $(y_1, y_2, y_3, y_4) = (0, 52/5, 0, 2/5)$ and the maximum value  $f$  is 42.

The dual problem is

min 
$$
g = 12x_1 + 7x_2 + 10x_3
$$
  
subject to  $3x_1 + x_2 + 2x_3 \ge 2$   
 $x_1 - 3x_2 + x_3 \ge 4$   
 $x_1 + 2x_2 + 3x_3 \ge 3$   
 $4x_1 + 3x_2 - x_3 \ge 1$ 

Thus an optimal vector for the primal problem is  $(x_1, x_2, x_3) = (1, 0, 3)$  and the minimum value  $q$  is 42.

2.(a) Add  $k = 3$  to every entry to get

$$
\begin{pmatrix} 5 & 0 & 6 \ 1 & 6 & 4 \ 4 & 4 & 8 \end{pmatrix}.
$$

Set up the tableau and apply pivoting operations, we have

y<sup>1</sup> y<sup>2</sup> y<sup>3</sup> x<sup>1</sup> y<sup>2</sup> y<sup>3</sup> <sup>∗</sup> 0 6 1 x<sup>1</sup> 5 y<sup>1</sup> 1/5 0 6/5 1/5 x<sup>2</sup> 1 6 4 1 → x<sup>2</sup> −1/5 6 14/5 4/5 → x<sup>3</sup> −4/5 4<sup>∗</sup> 16/5 1/5 x<sup>3</sup> 4 4 8 1 −1/5 1 −1/5 −1/5 1 1 1 0 x<sup>1</sup> x<sup>3</sup> y<sup>3</sup> y<sup>1</sup> 1/5 0 6/5 1/5 → x<sup>2</sup> 1 −3/2 −2 1/2 . y<sup>2</sup> −1/5 1/4 4/5 1/20 0 −1/4 −1 −1/4

Therefore,  $d = 1/4$  and a maximin strategy for the row player is

$$
\mathbf{p} = \frac{1}{d}(x_1, x_2, x_3) = (0, 0, 1),
$$

a minimax strategy for the column player is

$$
\mathbf{q} = \frac{1}{d}(y_1, y_2, y_3) = (4/5, 1/5, 0),
$$

the value of the game is  $v = \frac{1}{d} - k = 1$ .

(b) Add  $k = 5$  to every entry to get

$$
\begin{pmatrix} 8 & 6 & 0 \\ 4 & 3 & 11 \\ 3 & 4 & 7 \end{pmatrix}.
$$

Set up the tableau and apply pivoting operations, we have



Therefore,  $d = 7/33$  and a maximin strategy for the row player is

$$
\mathbf{p} = \frac{1}{d}(x_1, x_2, x_3) = (4/7, 3/7, 0),
$$

a minimax strategy for the column player is

$$
\mathbf{q} = \frac{1}{d}(y_1, y_2, y_3) = (0, 11/14, 3/14),
$$

the value of the game is  $v = \frac{1}{d} - k = -2/7$ .

(c) Add  $k = 2$  to every entry to get

$$
\begin{pmatrix} 5 & 2 & 3 \ 1 & 4 & 0 \ 2 & 3 & 1 \end{pmatrix}.
$$

Set up the tableau and apply pivoting operations, we have



Therefore,  $d=5/12$  and a maximin strategy for the row player is

$$
\mathbf{p} = \frac{1}{d}(x_1, x_2, x_3) = (4/5, 1/5, 0),
$$

a minimax strategy for the column player is

$$
\mathbf{q} = \frac{1}{d}(y_1, y_2, y_3) = (0, 3/5, 2/5),
$$

the value of the game is  $v = \frac{1}{d} - k = 2/5$ .

(d) Add  $k = 3$  to every entry to get

$$
\begin{pmatrix} 5 & 3 & 1 \ 2 & 0 & 6 \ 1 & 5 & 3 \end{pmatrix}.
$$

Set up the tableau and apply pivoting operations, we have



Therefore,  $d = 12/35$  and a maximin strategy for the row player is

$$
\mathbf{p} = \frac{1}{d}(x_1, x_2, x_3) = (5/12, 1/4, 1/3),
$$

a minimax strategy for the column player is

$$
\mathbf{q} = \frac{1}{d}(y_1, y_2, y_3) = (1/3, 7/24, 3/8),
$$

the value of the game is  $v = \frac{1}{d} - k = -1/12$ .

(e) Add  $k = 2$  to every entry to get

$$
\begin{pmatrix} 3 & 1 & 3 \ 0 & 2 & 1 \ 3 & 0 & 4 \ 1 & 3 & 0 \end{pmatrix}.
$$

Set up the tableau and apply pivoting operations, we have



Therefore,  $d=5/9$  and a maximin strategy for the row player is

$$
\mathbf{p} = \frac{1}{d}(x_1, x_2, x_3, x_4) = (3/5, 0, 0, 2/5),
$$

a minimax strategy for the column player is

$$
\mathbf{q} = \frac{1}{d}(y_1, y_2, y_3) = (0, 3/5, 2/5),
$$

the value of the game is  $v = \frac{1}{d} - k = -1/5$ .

(f) Add  $k = 3$  to every entry to get

$$
\begin{pmatrix} 0 & 5 & 3 \\ 4 & 1 & 2 \\ 2 & 3 & 5 \\ 4 & 4 & 0 \end{pmatrix}.
$$

Set up the tableau and apply pivoting operations, we have



 $y_2$  |  $-5/17$  4/17 2/17 | 1/17

 $-1/17$   $-7/68$   $-3/17$   $-23/68$ 

.

Therefore,  $d = 23/68$  and a maximin strategy for the row player is

$$
\mathbf{p} = \frac{1}{d}(x_1, x_2, x_3, x_4) = (0, 4/23, 12/23, 7/23),
$$

a minimax strategy for the column player is

0  $-1/4$  1  $|-1/4$ 

 $y_2$   $\begin{array}{|l}$   $-1/3$   $1/3$   $-2/3$  0

$$
\mathbf{q} = \frac{1}{d}(y_1, y_2, y_3) = (13/23, 4/23, 6/23),
$$

the value of the game is  $v = \frac{1}{d} - k = -1/23$ .

3.(a) For any  $x_1, x_2 \in C_1 \cap C_2$ ,  $\lambda \in [0,1]$ , by the convexity of  $C_1, C_2$ ,  $\lambda x_1 + (1-\lambda)x_2 \in C_1$  and  $\lambda x_1 + (1-\lambda)x_2 \in C_2$ . Hence  $\lambda x_1 + (1-\lambda)x_2 \in C_1 \cap C_2$ , that is  $C_1 \cap C_2$  is convex.

(b) For any  $x = x_1 + x_2 \in C_1 + C_2$ ,  $y = y_1 + y_2 \in C_1 + C_2$  and  $\lambda \in [0, 1]$ , by the convexity of  $C_1$ ,  $C_2$ ,  $\lambda x_1 + (1 - \lambda)y_1 \in C_1$  and  $\lambda x_2 + (1 - \lambda)y_2 \in C_2$ . Hence  $\lambda x + (1 - \lambda)y \in C_1 + C_2$ , that is  $C_1 + C_2$  is convex.

4. Let C be the set of maximin strategy for the row player of A and  $\nu$ be the value of the game with game matrix A. For any  $\mathbf{u} = (u_1, u_2, ..., u_n)$ ,  $\mathbf{v} = (v_1, v_2, ..., v_n) \in C$  and  $\lambda \in [0, 1]$ , then, by the definition of C, we have  $\mathbf{u} A \mathbf{y}^T \ge \nu$  and  $\mathbf{v} A \mathbf{y}^T \ge \nu$  for any  $\mathbf{y} \in P^m$ , and  $\sum_{i=1}^n u_i = \sum_{i=1}^n v_i = 1$ . Let  $\mathbf{w} = (w_1, w_2, ..., w_n) = \lambda \mathbf{u} + (1 - \lambda)\mathbf{v}, \text{ then } w_i = \lambda \overline{u_i + (1 - \lambda)v_i} \text{ with } \sum_{i=1}^n w_i =$  $\lambda + (1 - \lambda) = 1$ , and  $\mathbf{w} A \mathbf{y}^T = (\lambda \mathbf{u} + (1 - \lambda) \mathbf{v}) A \mathbf{y}^T \ge \lambda \nu + (1 - \lambda) \nu = \nu$  for any  $y \in P^m$ . Hence,  $w \in C$ , that is, C is convex.

5.(a)  $z = \lambda x + (1 - \lambda)y$  with  $\lambda \in \mathbb{R}$ . Since z is orthogonal to  $x - y$ , we have  $\langle \mathbf{x} - \mathbf{y}, \mathbf{z} \rangle = \langle \mathbf{x} - \mathbf{y}, \lambda \mathbf{x} + (1 - \lambda) \mathbf{y} \rangle = 0$ . Thus,

$$
\lambda = \frac{\|{\bf y}\|^2 - {\bf x},{\bf y} >}{\|{\bf x}\|^2 - 2 < {\bf x},{\bf y} > + \|{\bf y}\|^2}.
$$

 $(b) < x, y > 0$  implies  $\lambda \in [0, 1]$ . Thus,  $z \in C$  since C is convex.

6.  $\nu_c(A) \leq 0$  implies there exists a minimax strategy  $\mathbf{q} = (q_1, q_2, ..., q_n)$  for the column player such that  $(-\lambda_1, -\lambda_2, ..., -\lambda_m)^T := A\mathbf{q}^T \leq \mathbf{0}^T$ , that is,  $\lambda_i \geq$ 0. Thus,  $\mathbf{0}^T = A\mathbf{q}^T + (\lambda_1, \lambda_2, ..., \lambda_m)^T = q_1\mathbf{a_1}^T + q_2\mathbf{a_2}^T + ... + q_n\mathbf{a_n}^T + \lambda_1\mathbf{e_1}^T +$  $\ldots + \lambda_m \mathbf{e_m}^T$ . Therefore,  $\mathbf{0}^T = l_1 \mathbf{a_1}^T + l_2 \mathbf{a_2}^T + \ldots + l_n \mathbf{a_n}^T + l_{n+1} \mathbf{e_1}^T + \ldots + l_{n+m} \mathbf{e_m}^T$ with  $l_i = \frac{q_i}{\sum q_i +}$  $\frac{q_i}{q_i + \sum \lambda_j} \in [0, 1], i = 1, 2, ..., n \text{ and } l_{n+j} = \frac{\lambda_j}{\sum q_i + \sum \lambda_j}$  $\frac{\lambda_j}{q_i + \sum \lambda_j} \in [0,1],$  $j = 1, 2, ..., m$  and  $\sum_{k=1}^{n+m} l_k = 1$ , that is,  $0 \in C$ .