

A space efficient streaming algorithm for triangle counting using the birthday paradox

Madhav Jha

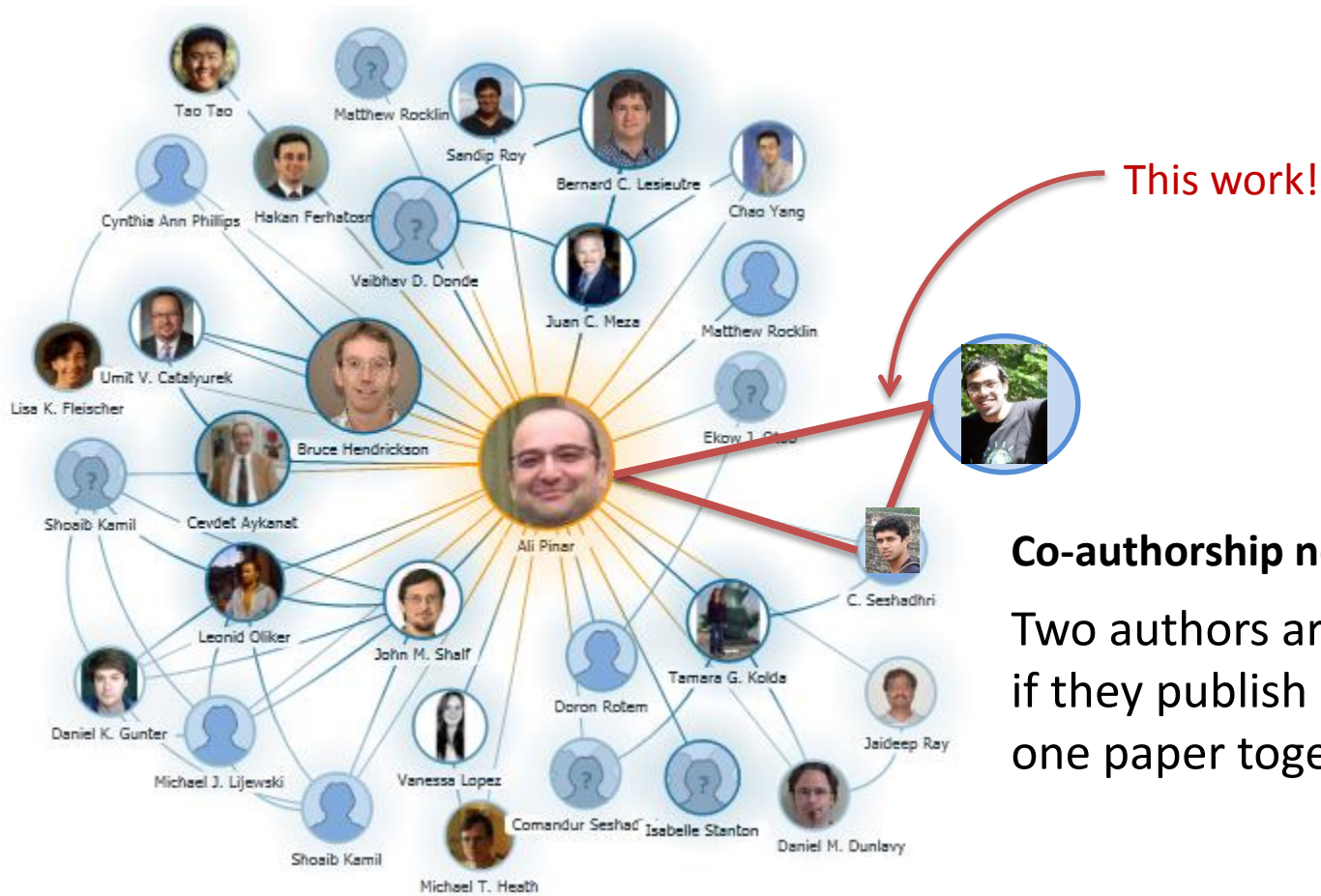
(Penn State → Sandia National Labs)

Joint work with
and

C. Seshadhri (Sandia National Labs)

Ali Pinar (Sandia National Labs)

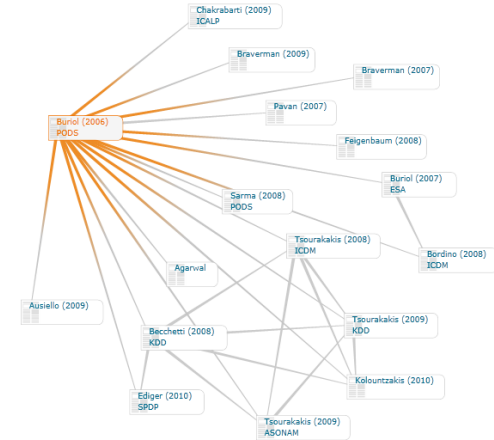
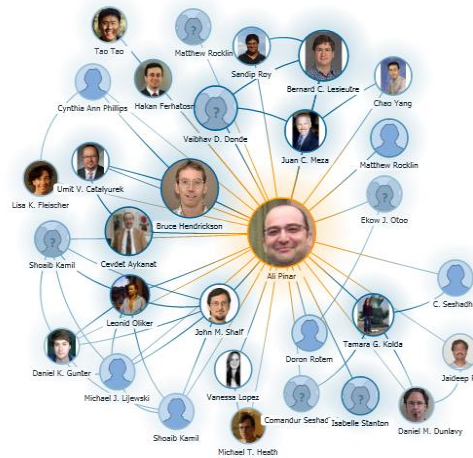
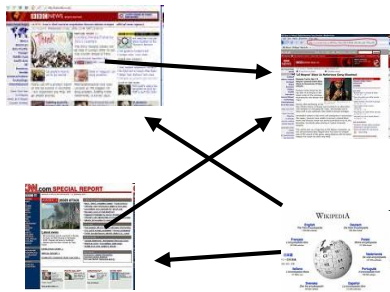
Real-world graphs: An Example



Co-authorship network
 Two authors are connected if they publish at least one paper together

| Graph [SNAP] | # nodes (n) | # edges (m) | # triangles (T) |
|--------------|-------------|-------------|-----------------|
| Ca-HepPh | 12K | 118K | 3.35M |

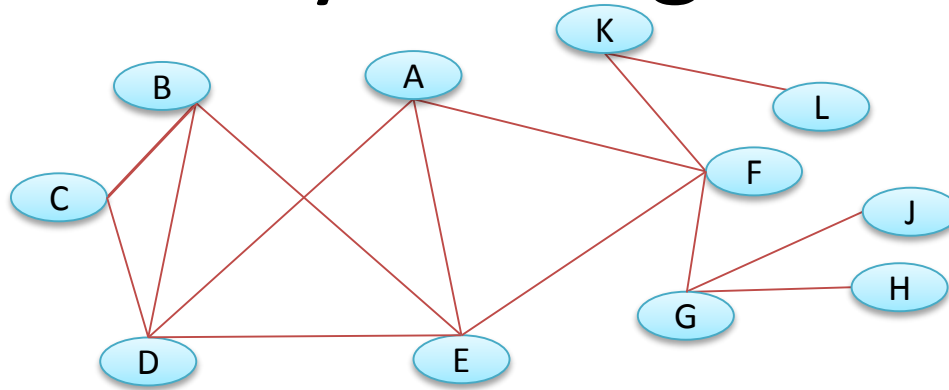
Real-world graphs



1. Graphs are everywhere.
2. Real-world graphs are huge. (Lots of vertices and edges.)
3. Real-world graphs have lots of triangles.

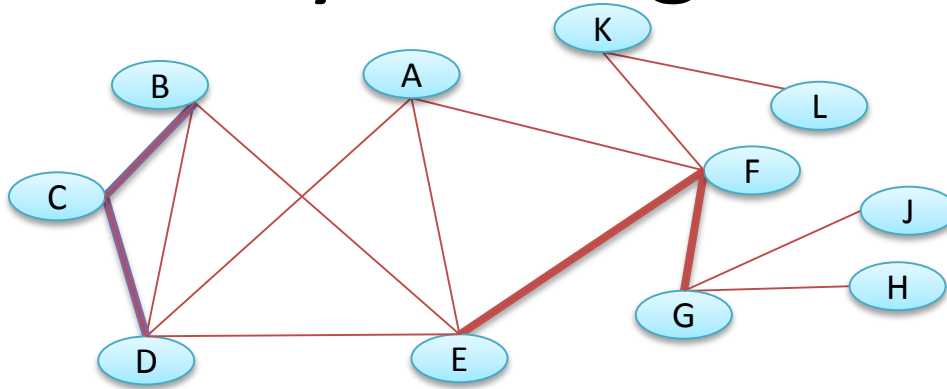
| Graph [SNAP] | # nodes (n) | # edges (m) | # triangles (T) |
|--------------|-------------|-------------|-----------------|
| web-BerkStan | 0.6M | 6M | 64M |
| orkut | 3M | 22M | 627M |
| Ca-HepPH | 12K | 118K | 3.35M |
| cit-Patents | 3M | 16M | 7M |

Transitivity: Triangle “density”



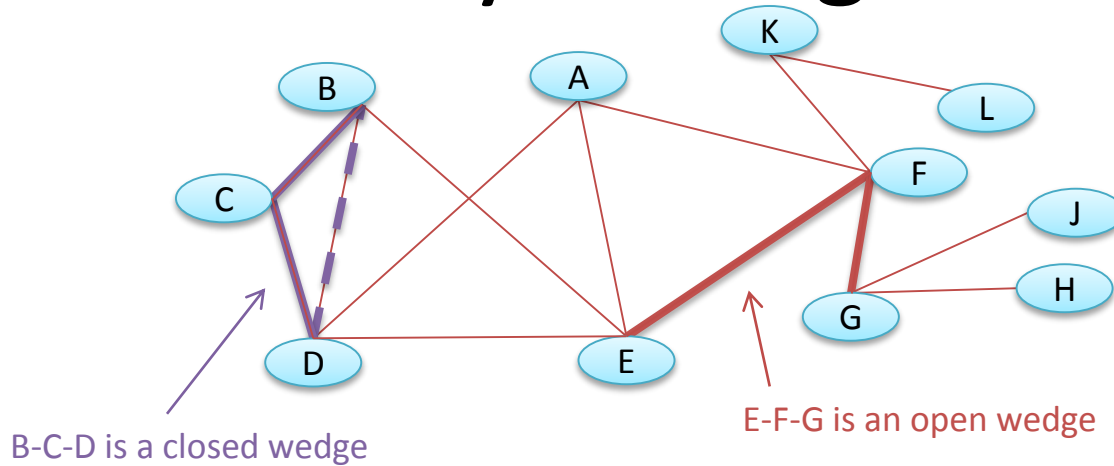
- A **wedge** is a length 2 path. Namely, a “potential” triangle.
- **Transitivity** = $\tau = 3 \text{ #Triangles} / \text{#Wedges}$ = fraction of closed wedges

Transitivity: Triangle “density”



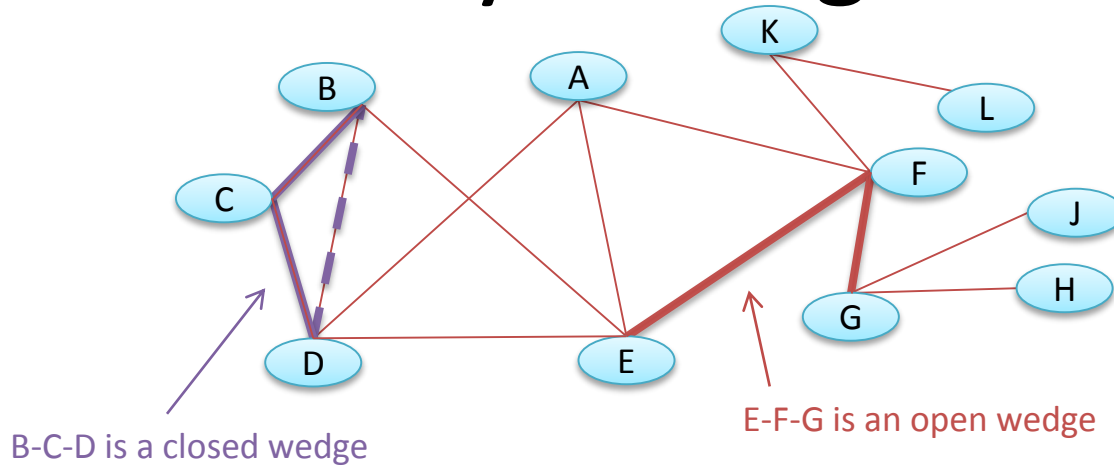
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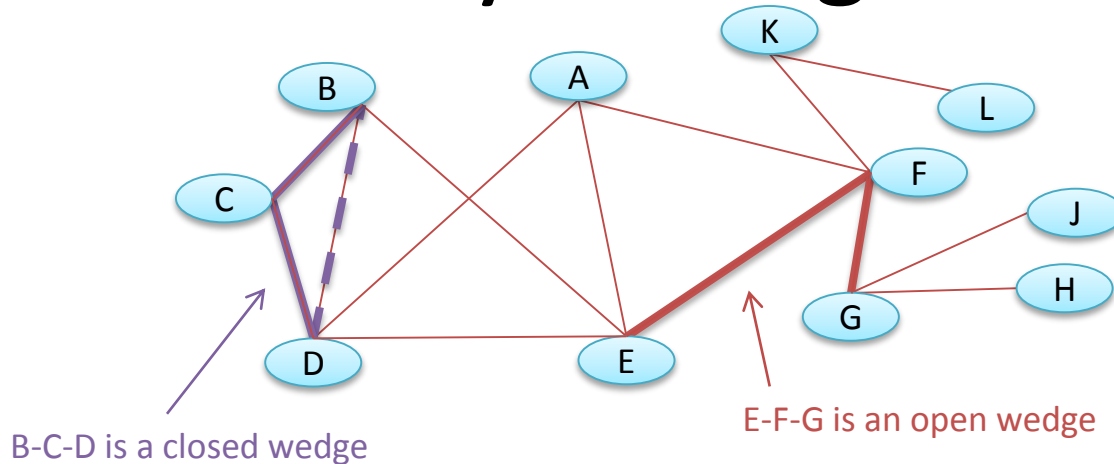
Transitivity: Triangle “density”



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| Graph [SNAP] | # nodes (n) | # edges (m) | # triangles (T) | Transitivity |
|--------------|-------------|-------------|-----------------|--------------|
| web-BerkStan | 0.6M | 6M | 64M | 0.007 |
| orkut | 3M | 223M | 627M | 0.041 |
| Ca-HepPH | 12K | 118K | 3.35M | 0.39 |
| cit-Patents | 3M | 16M | 7M | 0.067 |

Transitivity: Triangle “density”



- A **wedge** is a length 2 path. Namely, a “potential” triangle.
 - **Transitivity** = $\tau = 3 \text{ #Triangles} / \text{#Wedges}$ = fraction of closed wedges
- [Seshadhri Pinar Kolda 2013] gave algorithm for computing transitivity given access to the entire graph. This algorithm is the starting point of work.

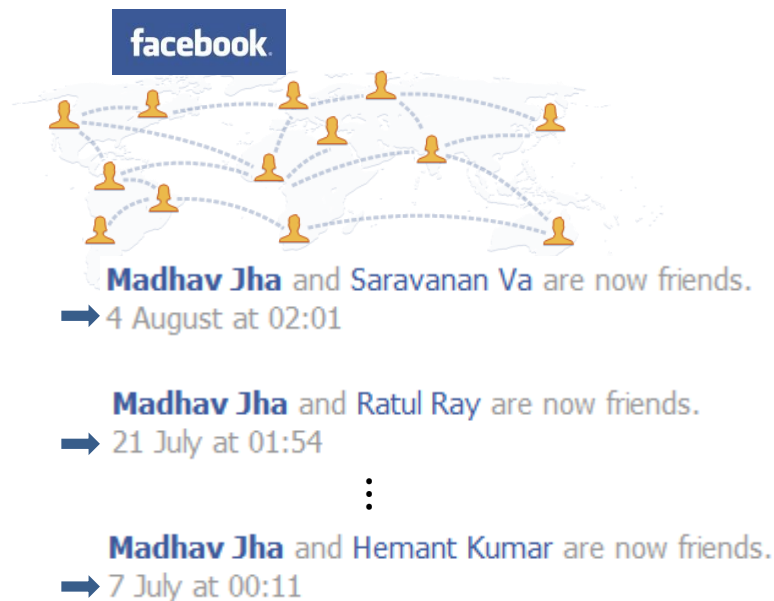
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Why Count Triangles in Graphs?

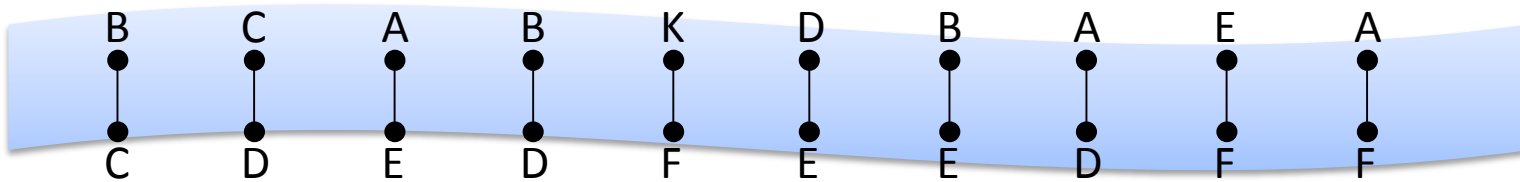
- Useful in Social Science for positing various theses on behavior
[Burt 09], [Coleman 88], [Welles, Devender, Contractor 10], [Portes 88]
- Applied to spam detection [Becchetti Boldi Castillo Gionis 08]
- Relevant for finding topics on WWW [Eckmann Moses 02]
- Proposed as a guide for community structure
Stated as a core feature for graph models [Vivar Banks 11]
Cornerstone for Block Two-level Erdos-Renyi (BTER) [Seshadhri Pinar Kolda 12]
- Good descriptor of the underlying graph [Durak Seshadhri Pinar Kolda 12]
- Rich set of algorithmic results spanning various models
(exact/approximate/deterministic/randomized/...) X (streaming, map-reduce, parallel etc.)
- Very well-studied: [Ahn Guha McGregorGraph 2012], [Durak Pinar Kolda Seshadhri 2012], [Pagh Tsourakakis 2012], [Suri Vassilvitskii 2011], [Tsourakakis Kolountzakis Miller 2011], [Chu Cheng 2011], [Yoon Kim 2011][Kolountzakis Miller Peng Tsourakakis 2010], [Avron 2010],[Tsourakakis Drineas Michelakis Koutis Faloutsos 2009], [Tsourakakis Kang Miller Faloutsos 2009], [Latapy 2008], [Becchetti Boldi Castillo Gionis 2008], [Tsourakakis 08], [Buriol Frahling Leonardi Marchetti-Spaccamela Sohler 2006], [Jowhari Ghodsi 2005], [Schank Wagner 2005], [Bar-Yossef Kumar Sivakumar 2002], ...

Graph as stream of edges

- Real-world graphs have a natural time-stamp



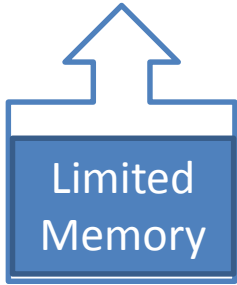
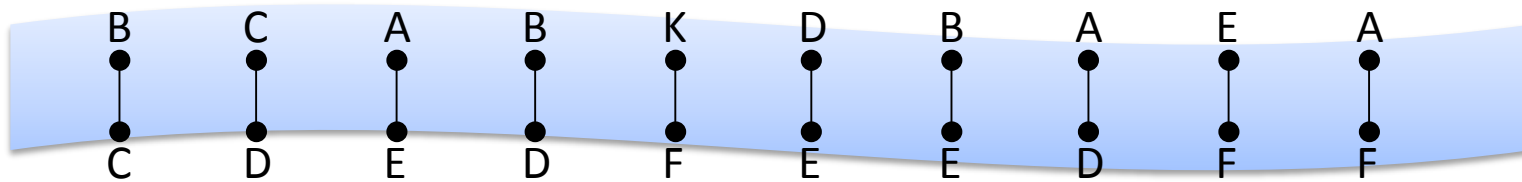
Graph as stream of edges



Triangles so far:

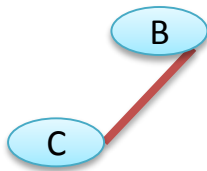
Graph seen so far:

Graph as stream of edges

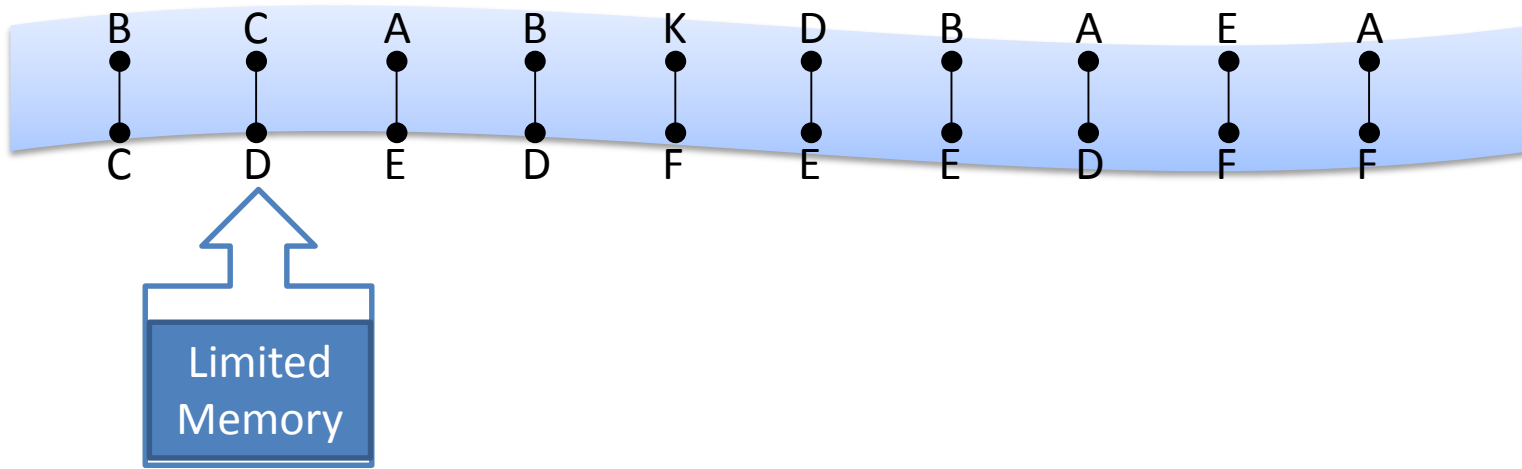


Triangles so far:

Graph seen so far:

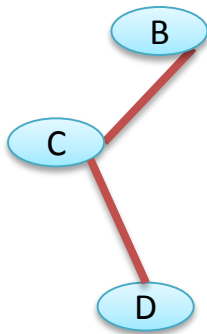


Graph as stream of edges

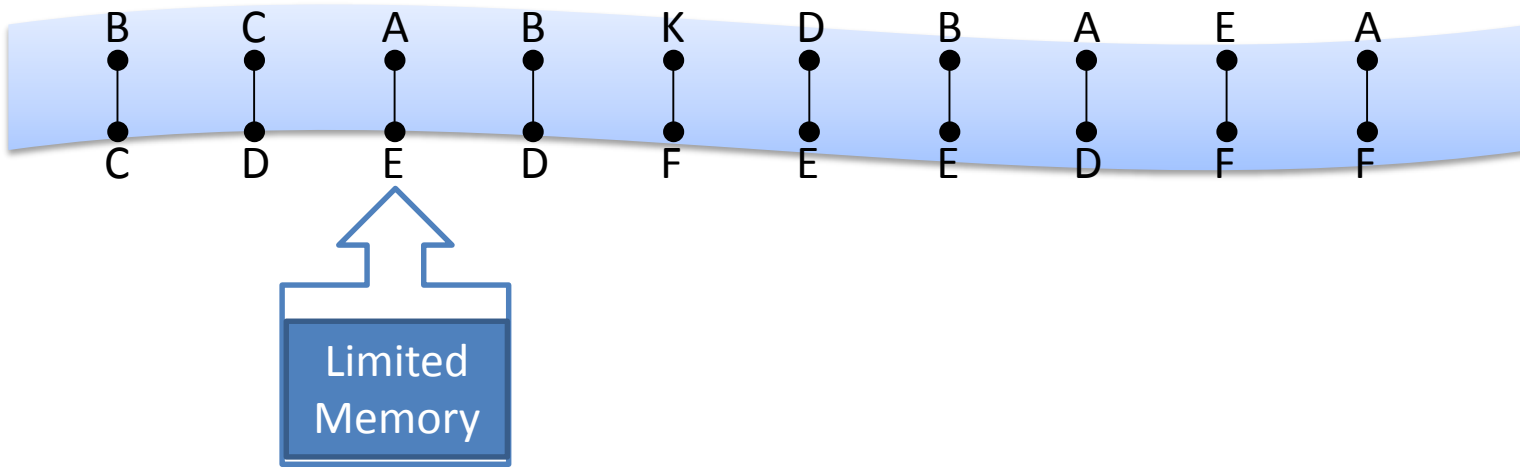


Triangles so far:

Graph seen so far:

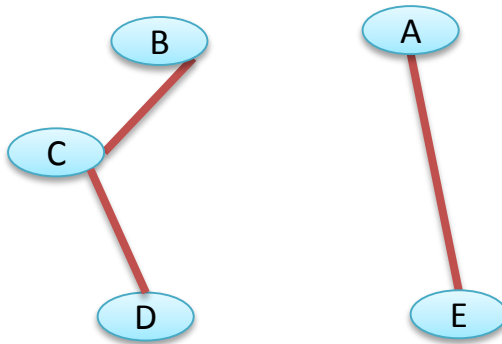


Graph as stream of edges

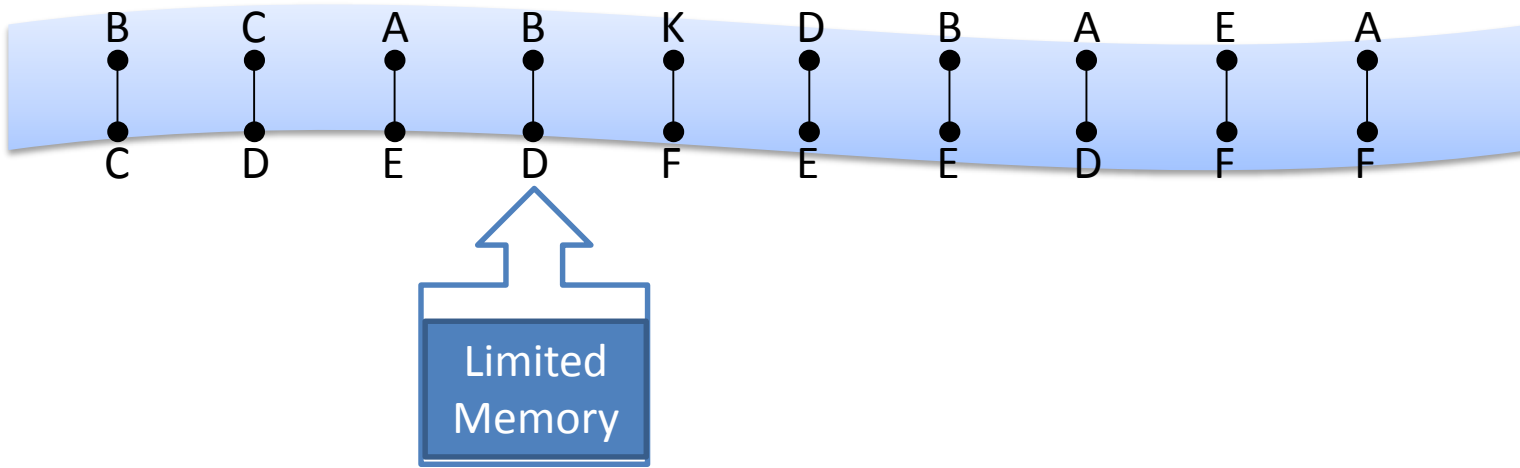


Triangles so far:

Graph seen so far:

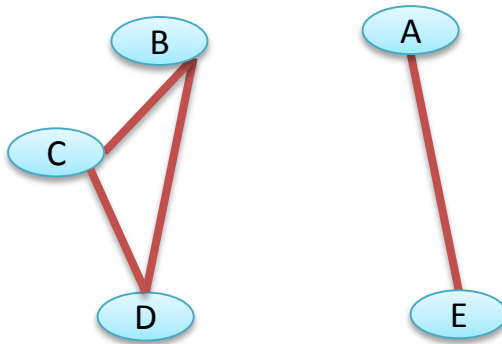


Graph as stream of edges

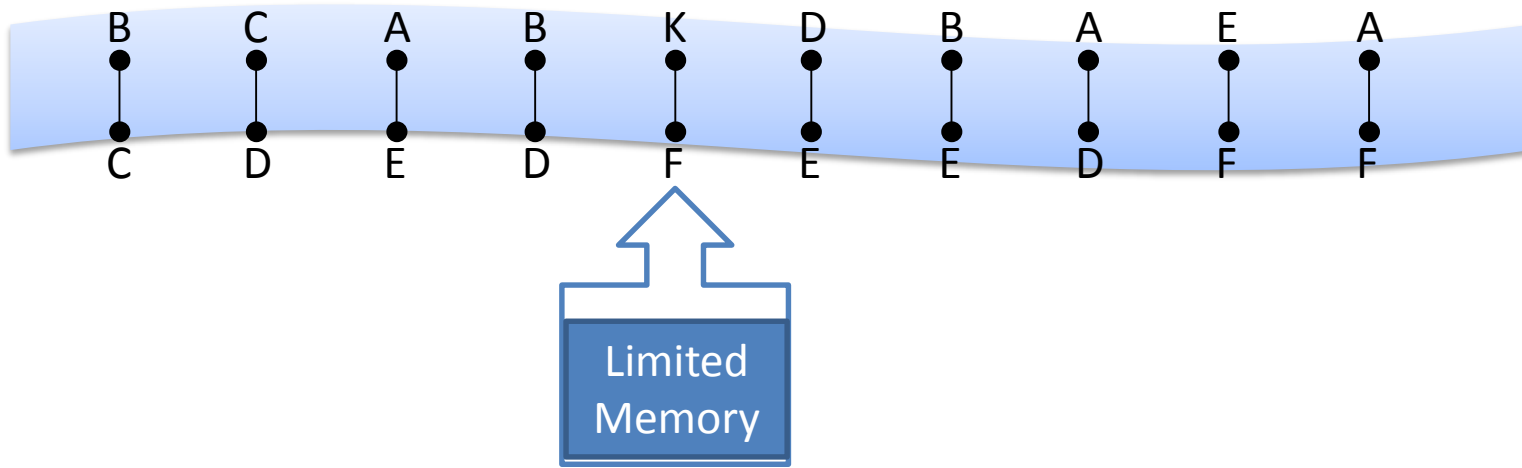


Triangles so far: 1

Graph seen so far:

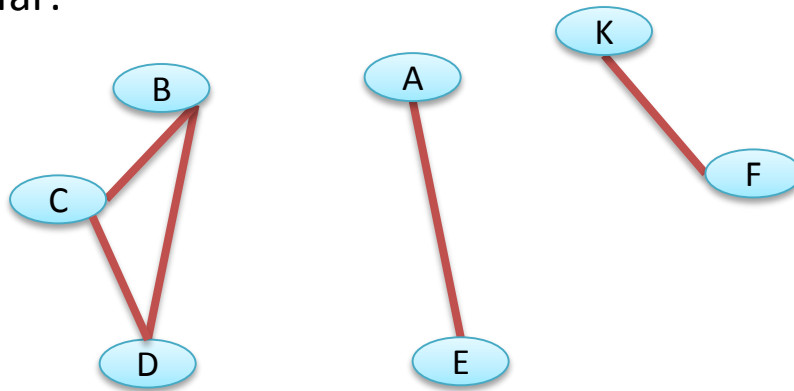


Graph as stream of edges

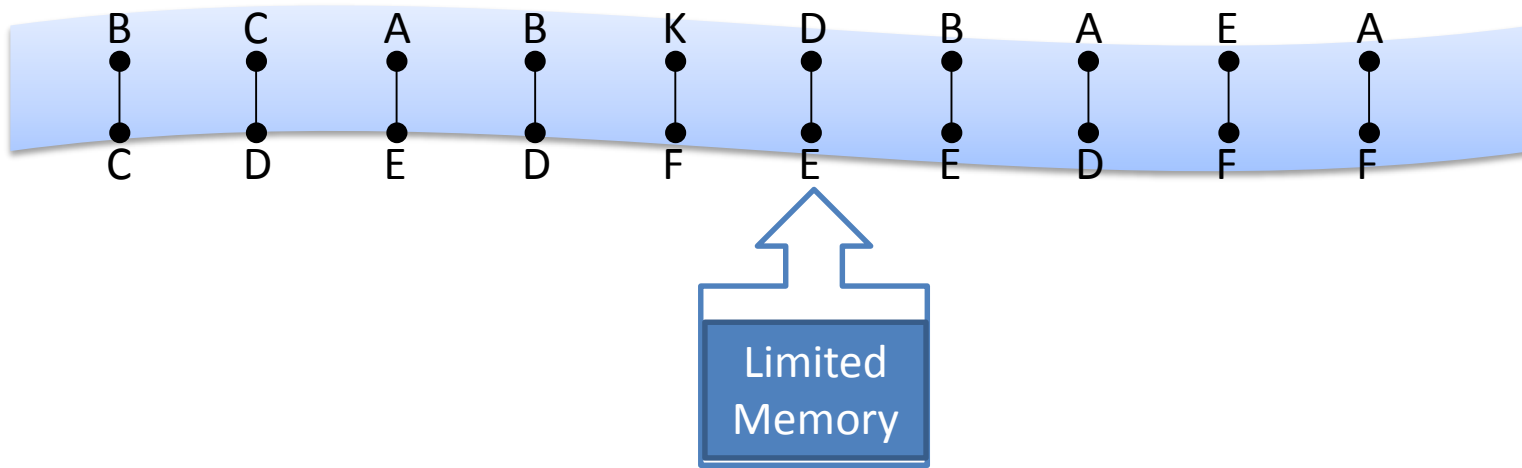


Triangles so far: 1

Graph seen so far:

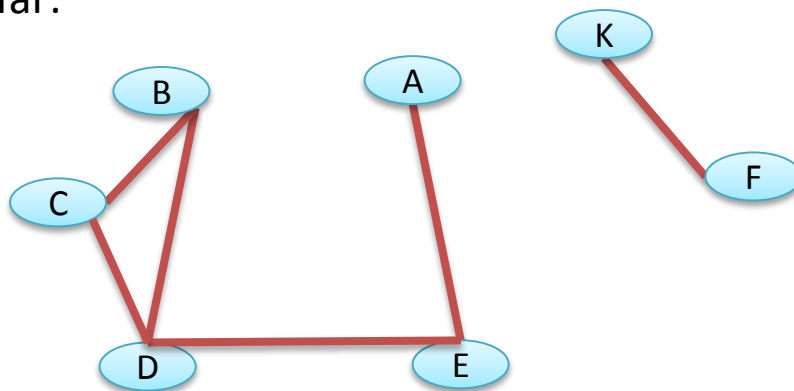


Graph as stream of edges

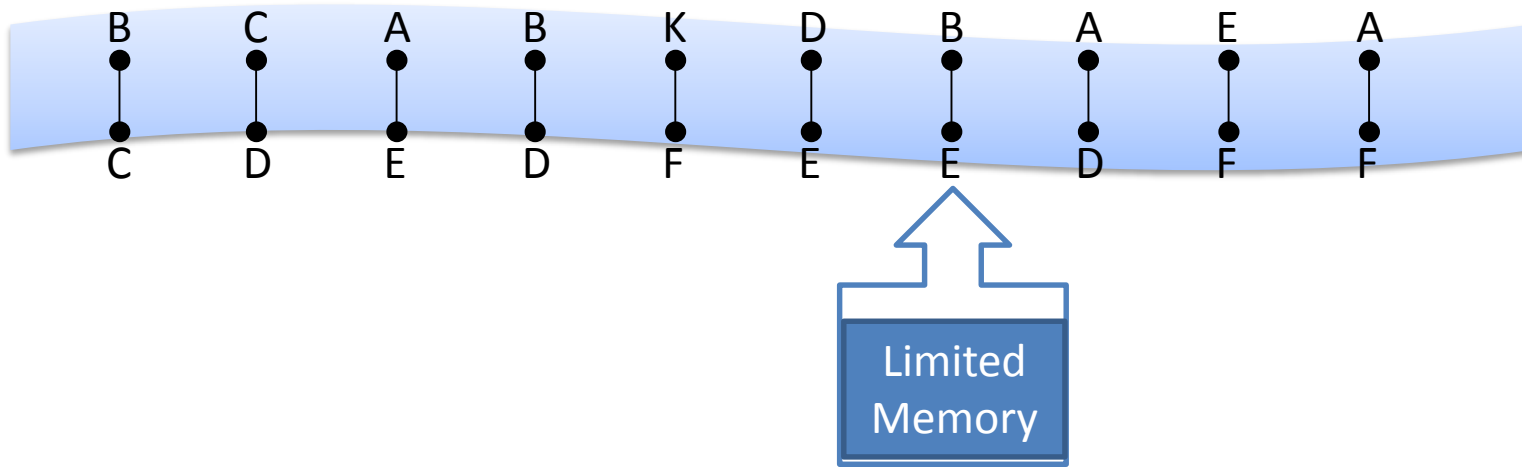


Triangles so far: 1

Graph seen so far:

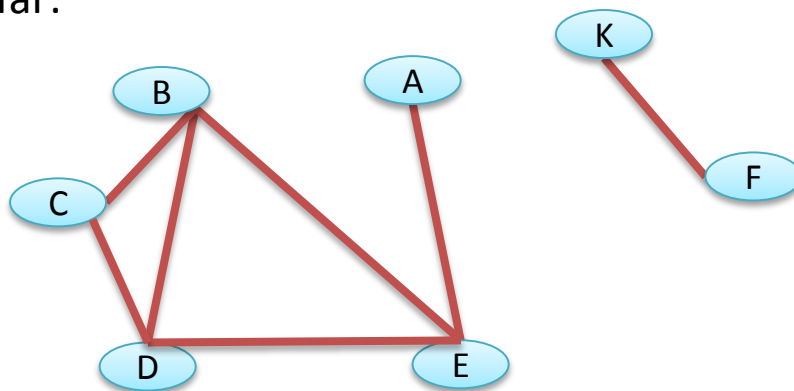


Graph as stream of edges

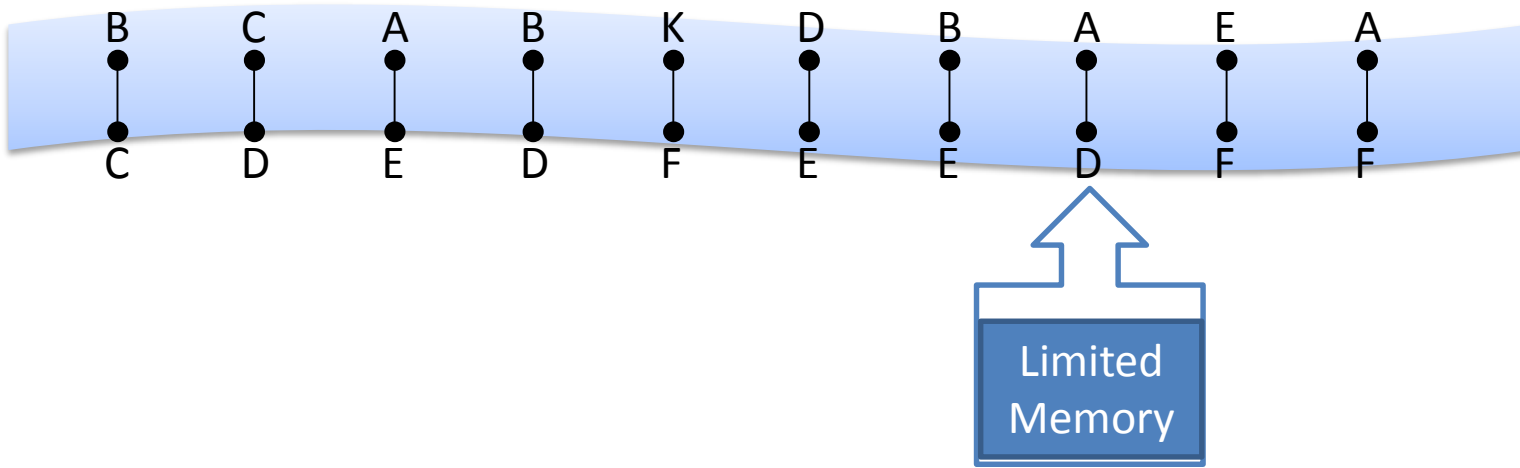


Triangles so far: 2

Graph seen so far:

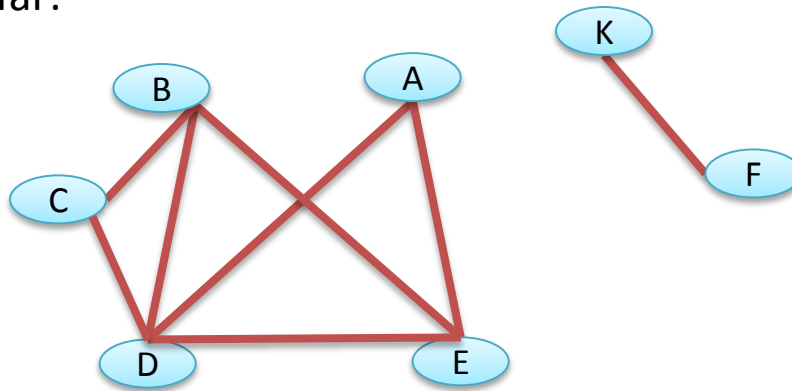


Graph as stream of edges

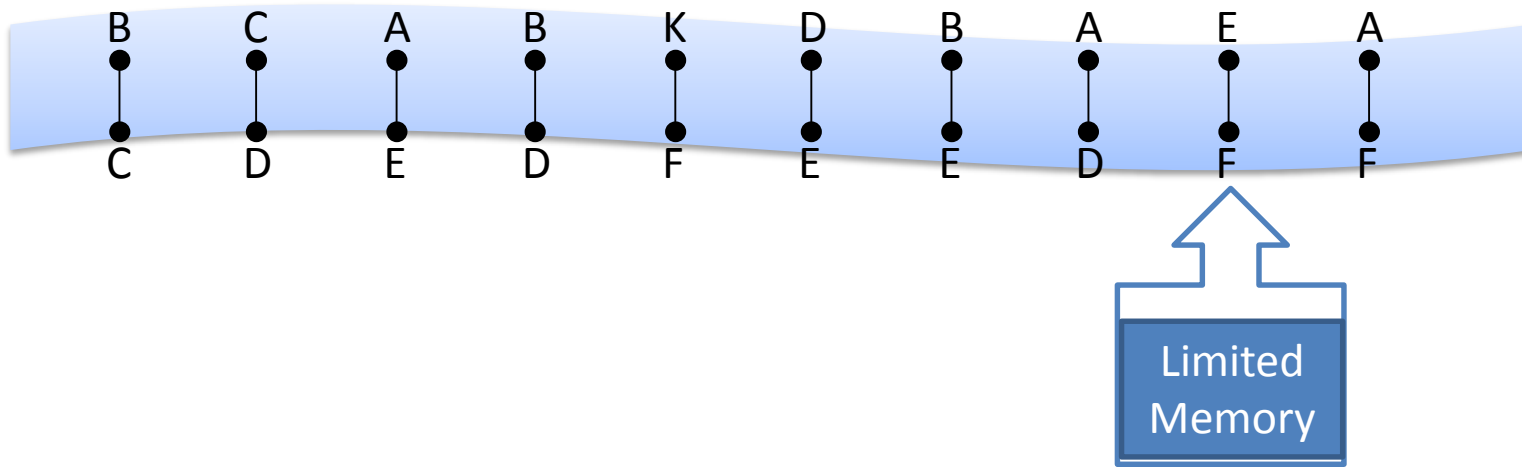


Triangles so far: 3

Graph seen so far:

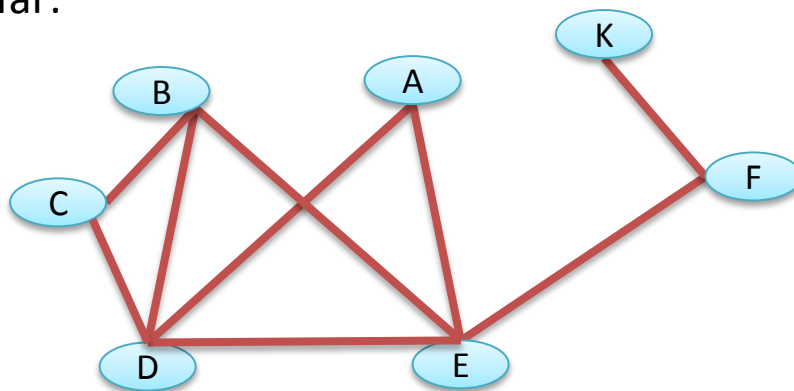


Graph as stream of edges

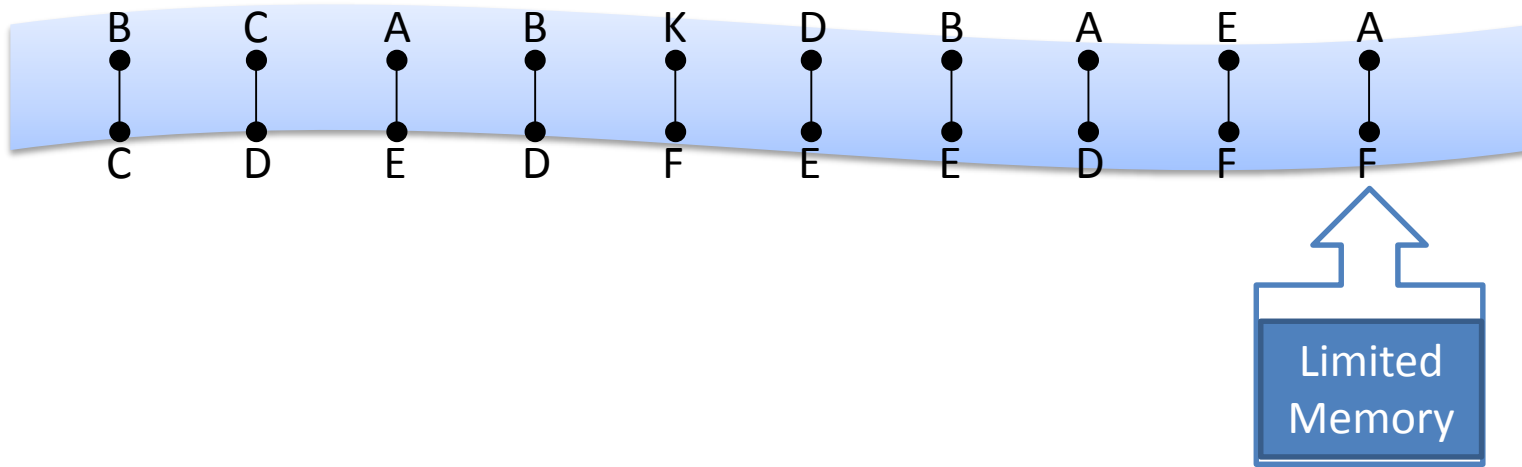


Triangles so far: 3

Graph seen so far:

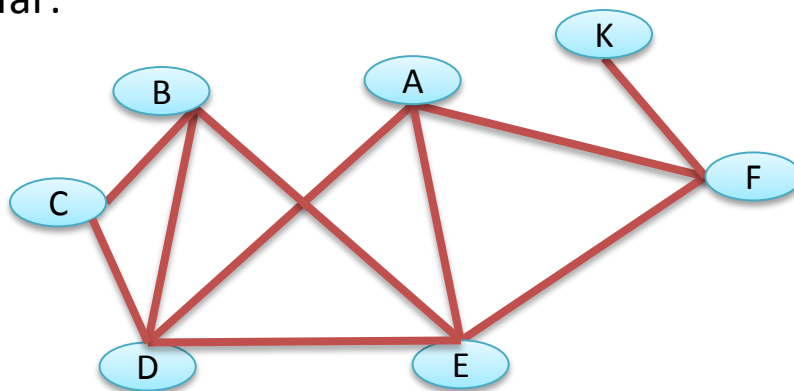


Graph as stream of edges



Triangles so far: 4

Graph seen so far:



Our Contributions : Theoretical

Theorem:

A **single-pass** streaming algorithm (for **arbitrarily ordered edge stream**) which **stores only $O(\sqrt{n})$ edges** (for most real world graphs), requires nearly **constant time update** per edge, and estimates **# triangles** and **transitivity**.

Analysis based on the classic **Birthday Paradox**.

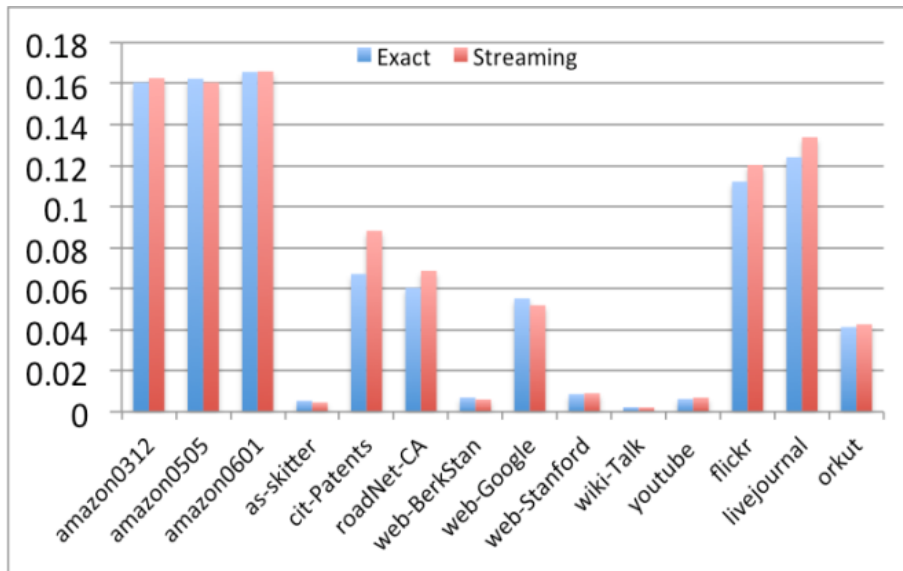
Our Contributions : Practical

- Accurate triangles estimates in low space

Example: On Orkut graph (200 M edges and 0.627 B triangles), our algorithm stores only 40 K edges (2% of graph) and reports 0.658 B triangles (less than 5% relative error).

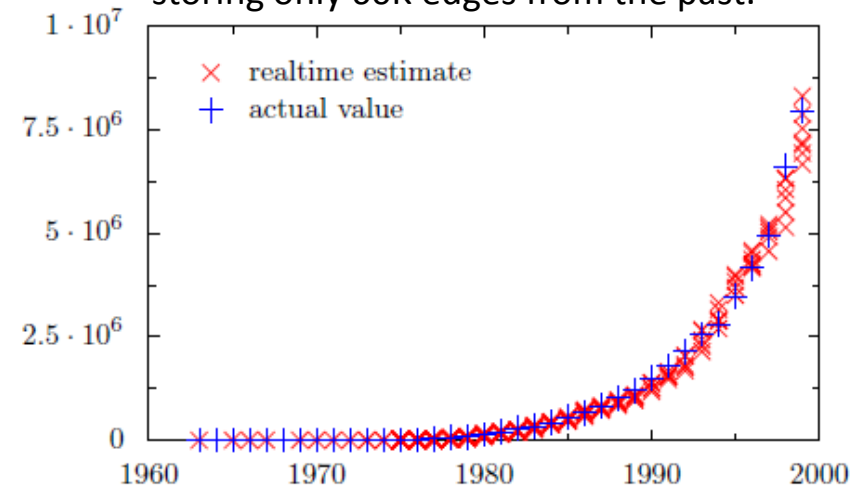
- Accurate transitivity estimates

- Realtime tracking



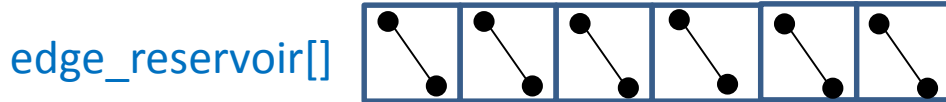
Estimating transitivity on a variety of dataset.
(Our algorithm stores only 40 K edges in all these runs.)

Realtime tracking of # triangles on cit-Patents graph (16M edges), storing only 60K edges from the past.



Data Structures of the Algorithm

Input Parameters: s_e and s_w .



An array to store **edges** of size s_e

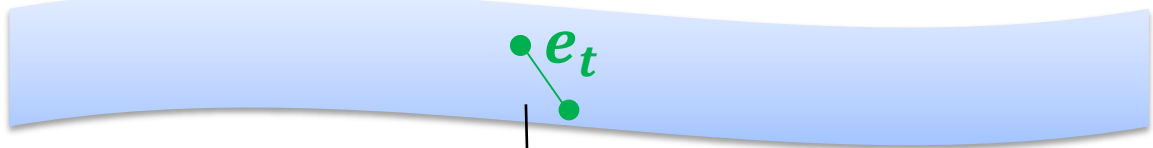


An array to store **wedges** of size s_w

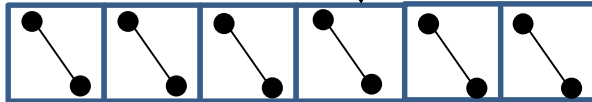


A Boolean array of size s_w

The Algorithm

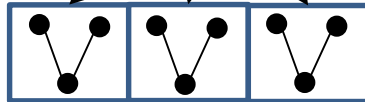
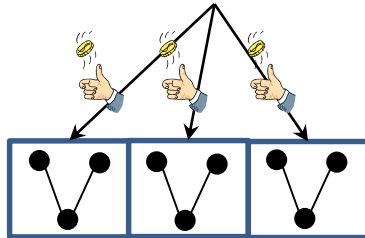


edge_reservoir[]



Update edge_reservoir

wedge_reservoir[]



Update wedge_reservoir

isClosed[]

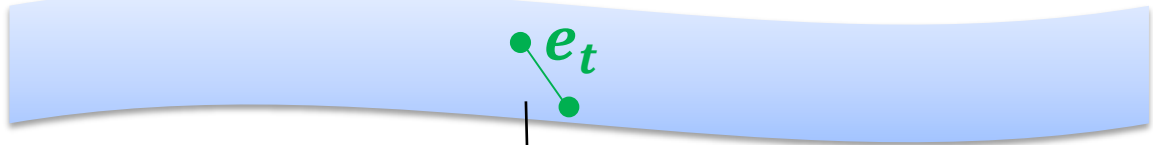


Update isClosed

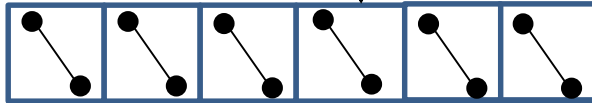
Let p be fraction of 1's in isClosed[]. Output

1. Transitivity, $est-\tau_t = 3p$
2. Triangles, $est-T_t = est-\tau_t \times \text{normalizing-factor}$

The Algorithm

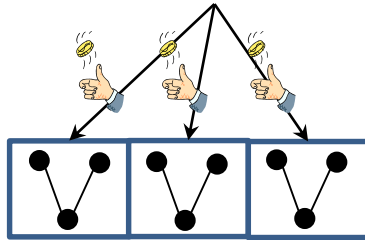


edge_reservoir[]



Update edge_reservoir

wedge_reservoir[]



Update wedge_reservoir

isClosed[]

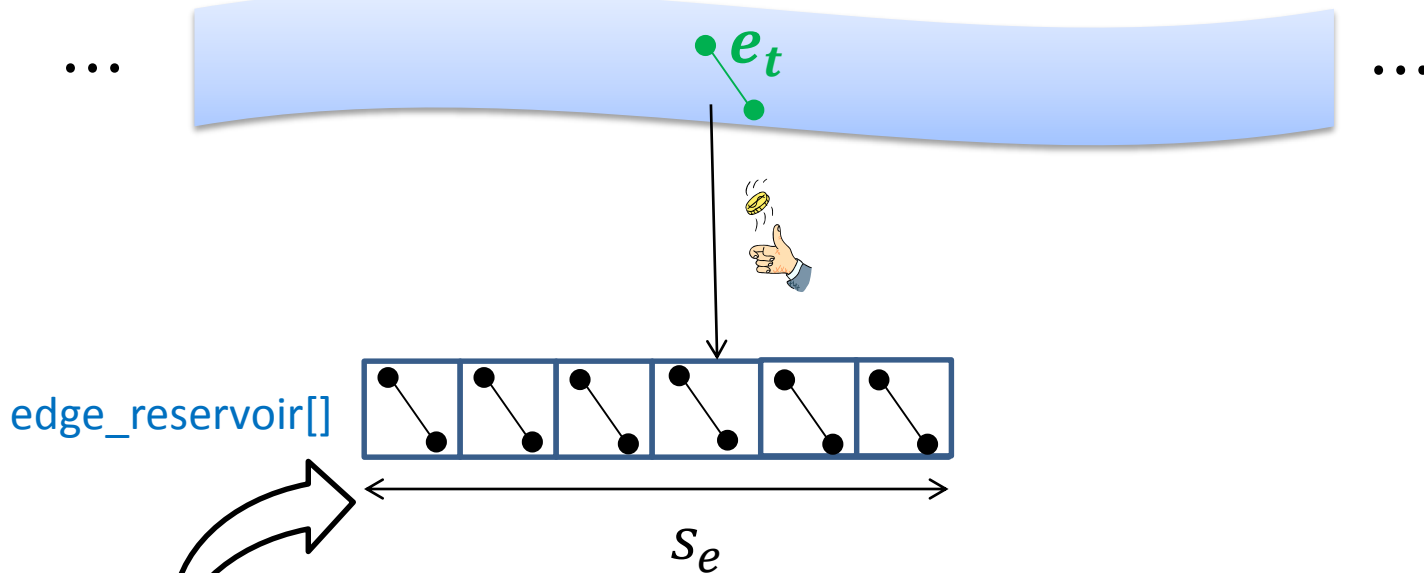


Update isClosed

Updates to edge_reservoir very rare!

$$\sum_{t \leq m} 1 - (1 - 1/t)^{s_e} \approx \sum_{t \leq m} s_e/t \approx s_e \ln m$$

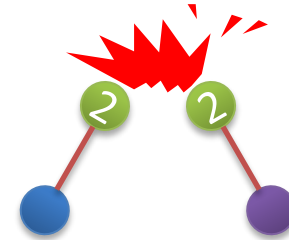
The Algorithm



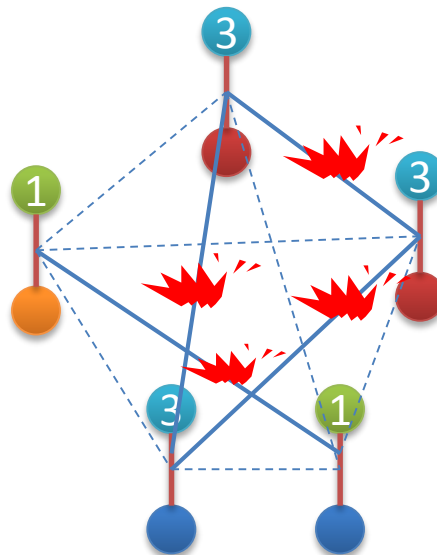
How many wedges are there in a random pool of s_e edges?

The Birthday Paradox to Rescue

Idea: Fundamentally, a wedge is a collision of two edges!



Birthday Paradox $\Rightarrow s_e$ edges give rise to $s_e^2 \cdot \Pr[\text{A single collision}]$

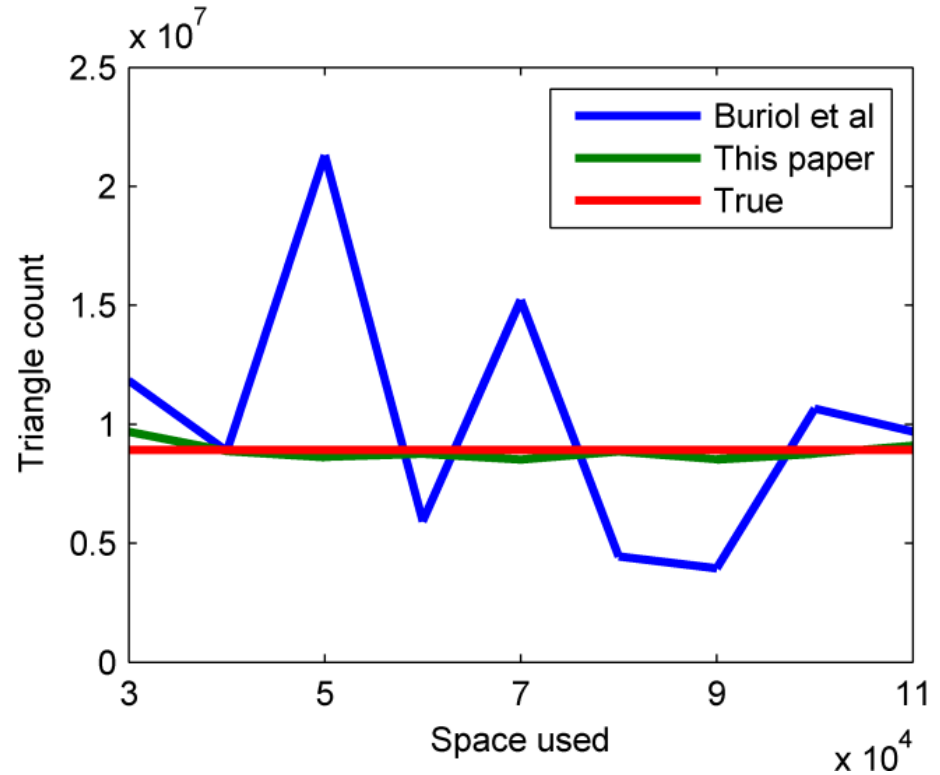
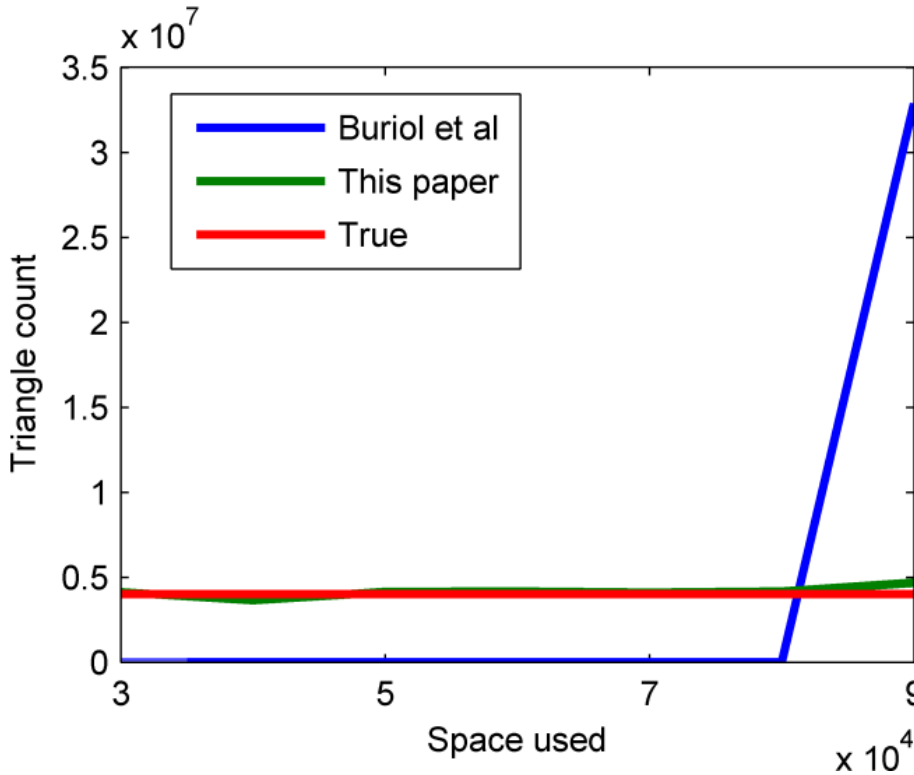


Experimental Results

Our Algorithm vs Buriol et al

Dataset: web-NotreDame

Dataset: amazon0505



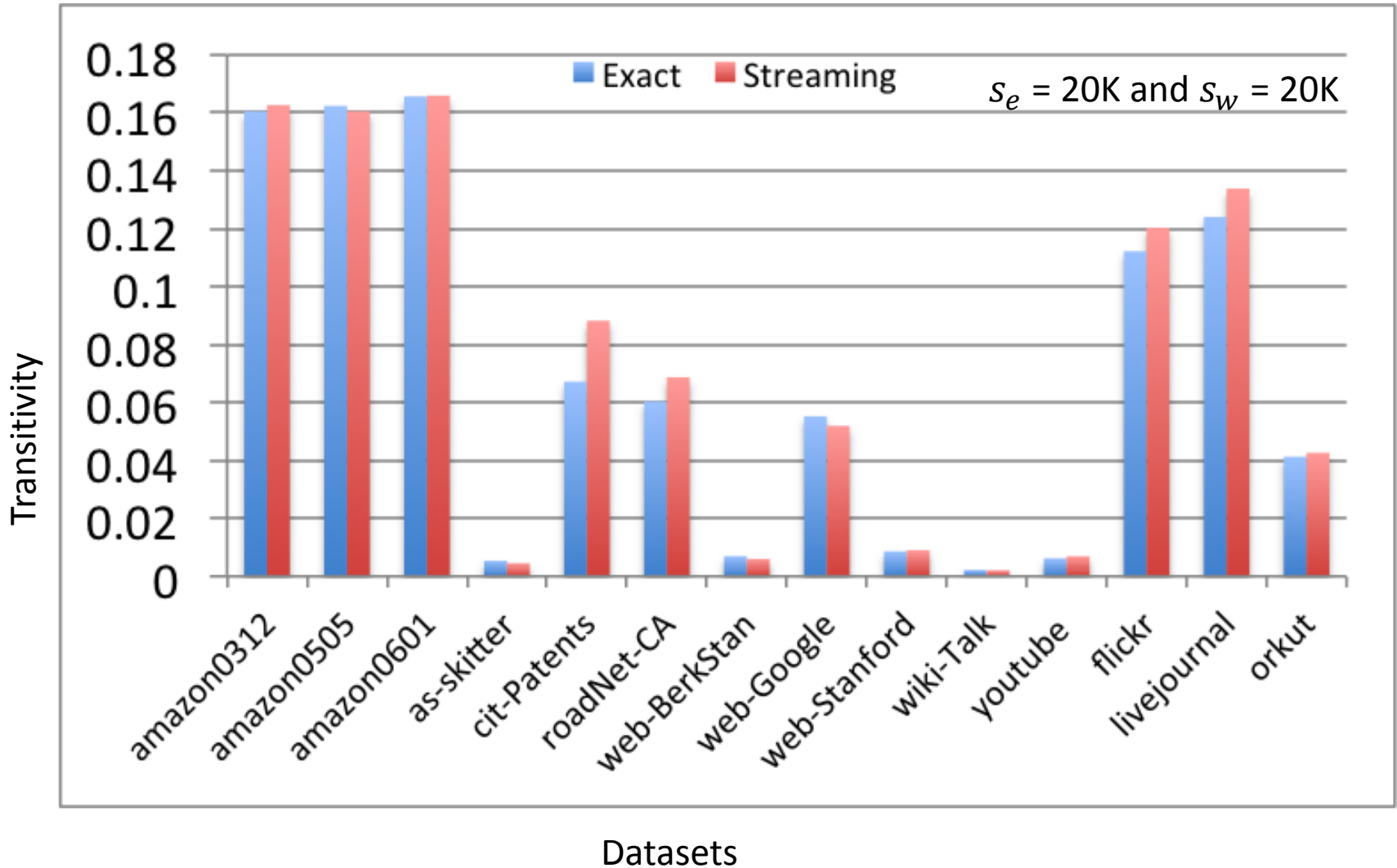
We fix $s_e = 20K$ and vary s_w

Space used in our algorithm: $s_e + s_w$

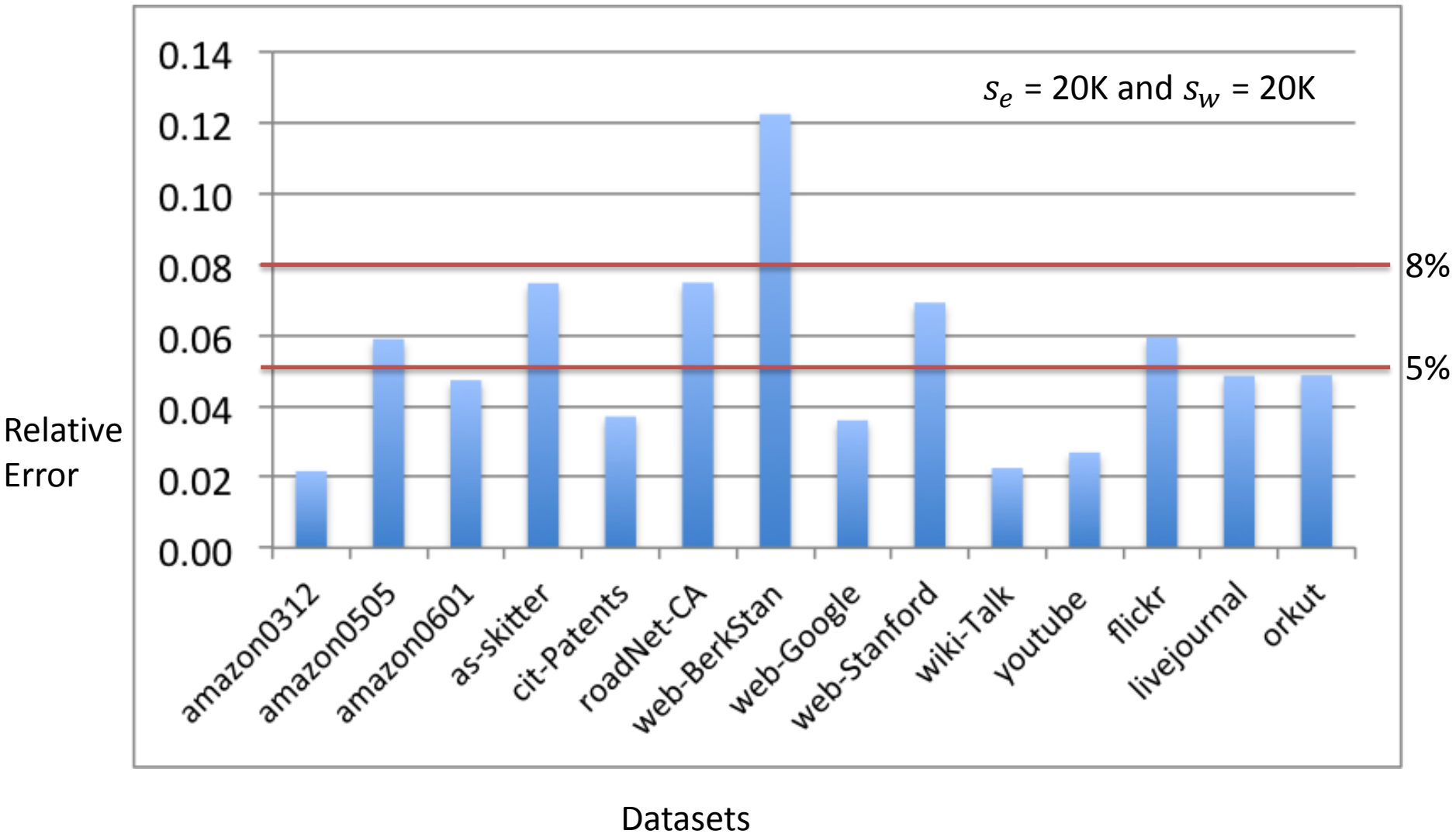
Space used in Buriol et al: number of edges sampled

Note: The results for Buriol et al is consistent with the analysis and experiments of their paper.

Accuracy of Transitivity Estimate



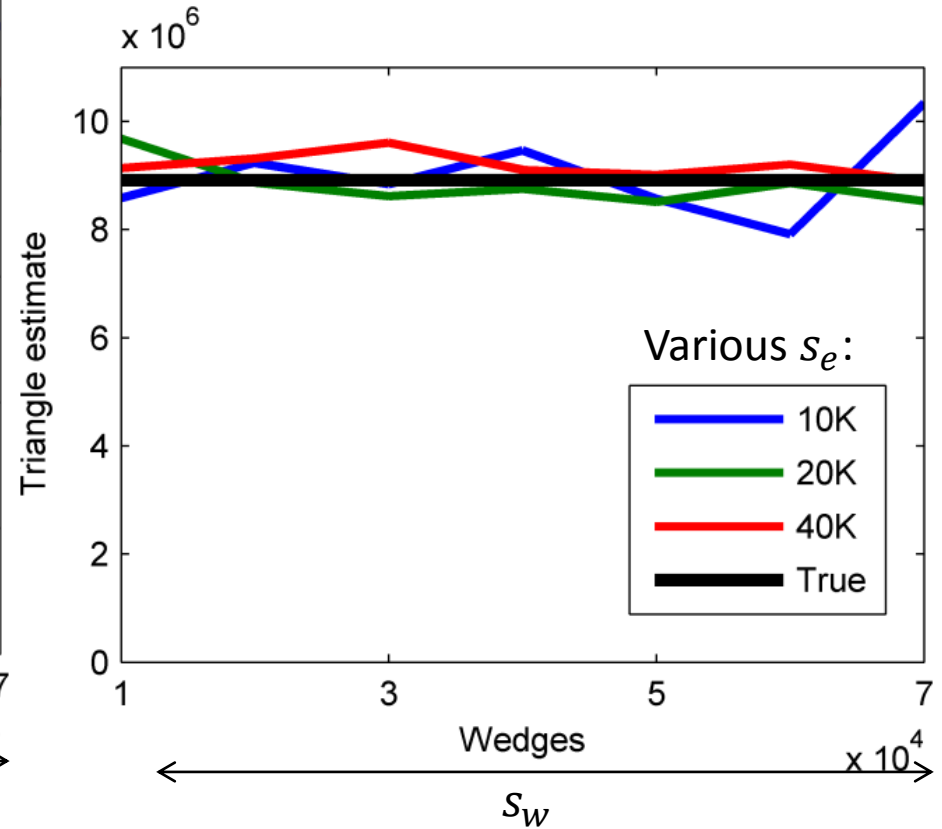
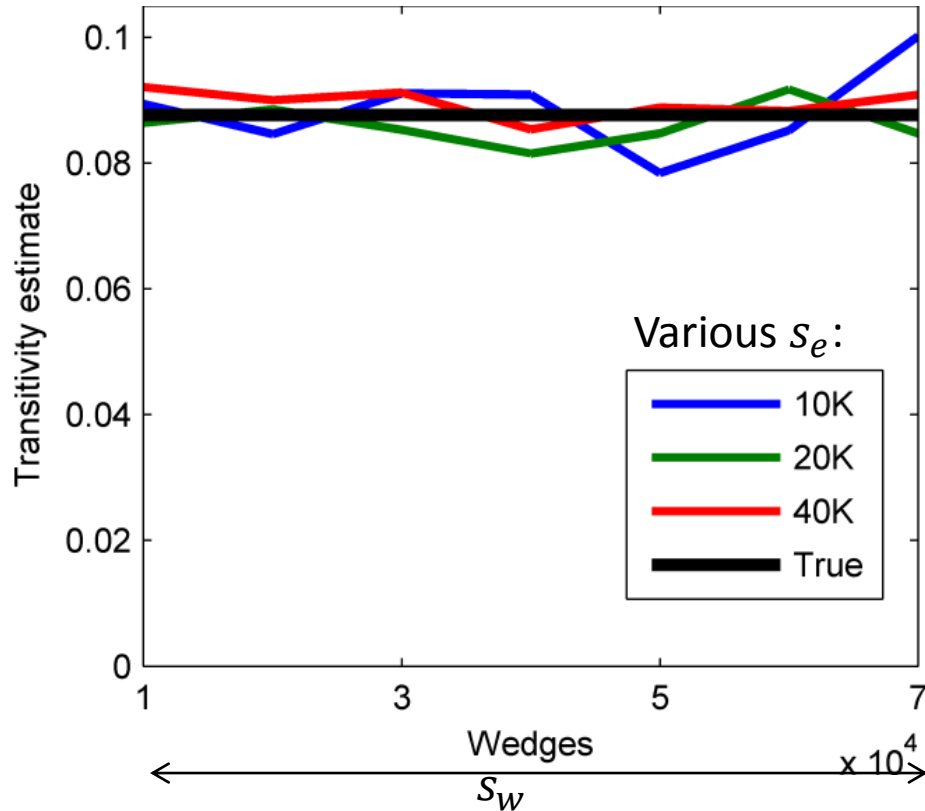
Accuracy of Triangles Estimate



Note: web-BerkStan has very low transitivity 0.007. Therefore, relative error is high.

Convergence of Estimates

Dataset: amazon0505



Future Work

- Can we go below \sqrt{n} space bound?
- Can we prove a lower bound on the space required by a 1-pass streaming algorithm to estimate triangle counts?
- Can we extend this approach to handle edge deletions ?
- Can we compute (and track) degree-wise clustering coefficient?

