

CSC2100B Data Structures

Recurrence Relations

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Recurrence Relations

- Recurrence relations are useful in certain counting problems.
- A recurrence relation relates the n -th element of a sequence to its predecessors.
- Recurrence relations arise naturally in the analysis of recursive algorithms.



Sequences and Recurrence Relations

- A (numerical) sequence is an ordered list of number.
 - 2, 4, 6, 8, ... (positive even numbers)
 - 0, 1, 1, 2, 3, 5, 8, ... (the Fibonacci numbers)
 - 0, 1, 3, 6, 10, 15, ... (numbers of key comparisons in selection sort)



Definitions

- A **recurrence relation** for the sequence, a_0, a_1, \dots is an equation that relates a_n to certain of its predecessors a_0, a_1, \dots, a_{n-1} .
- **Initial conditions** for the sequence a_0, a_1, \dots are explicitly given values for a finite number of the terms of the sequence.



Example

- A person invests \$1000 at 12% compounded annually. If A_n represents the amount at the end of n years, find a recurrence relation and initial conditions that define the sequence A_n .
- At the end of $n - 1$ years, the amount is A_{n-1} . After one more year, we will have the amount A_{n-1} plus the interest. Thus $A_n = A_{n-1} + (0.12)A_{n-1} = (1.12)A_{n-1}, n \geq 1$.
- To apply this recurrence relation for $n = 1$, we need to know the value of A_0 which is 1000.



Solving Recurrence Relations

- **Iteration** - we use the recurrence relation to write the n -th term a_n in terms of certain of its predecessors a_{n-1}, \dots, a_0 .
- We then successively use the recurrence relation to replace each of a_{n-1}, \dots by certain of their predecessors.
- We continue until an explicit formula is obtained.



Some Definitions of Linear Second-order recurrences with constant coefficients

- **kth-order**
 - Elements $x(n)$ and $x(n-k)$ are k positions apart in the unknown sequence.
- **Linear**
 - It is a linear combination of the unknown terms of the sequence.
- **Constant coefficients**
 - The assumption that a , b , and c are some fixed numbers.
- **Homogeneous**
 - If $f(x) = 0$ for every n .



Solving Recurrence Relations

- **Linear homogeneous** recurrence relations with constant coefficients - a linear homogeneous recurrence relation of order k with constant coefficients is a recurrence relation of the form

$$a_0 = c_0, a_1 = c_1, \dots, a_{k-1} = c_{k-1}$$

- Notice that a linear homogeneous recurrence relation of order K with constant coefficients, together with the k initial conditions

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}, c_k \geq 0$$

uniquely defines a sequence a_0, a_1, \dots



Example

- Nonlinear

$$a_n = 3a_{n-1}a_{n-2}.$$

- Inhomogeneous

$$a_n - a_{n-1} = 2n.$$

- Homogeneous recurrence relation with nonconstant coefficients

$$a_n = 3n \cdot a_{n-1}.$$



Iteration Example

- We can solve the recurrence relation $a_n = a_{n-1} + 3$ subject to the initial condition $a_1 = 2$, by iteration.
- $a_{n-1} = a_{n-2} + 3.$
- $a_n = a_{n-1} + 3 = a_{n-2} + 3 + 3 = a_{n-2} + 2 \times 3.$
- $a_{n-2} = a_{n-3} + 3.$
- $a_n = a_{n-2} + 2 \times 3 = a_{n-3} + 3 + 2 \times 3 = a_{n-3} + 3 \times 3.$
- $a_n = a_{n-k} + k \times 3 = 2 + 3(n - 1).$



Iteration Example

- In general, to solve $a_n = a_{n-1} + k, a_1 = c$, one obtains $a_n = c + k(n - 1)$.
- We can solve the recurrence relation
 - $a_n = ka_{n-1}, a_0 = c$.
 - $a_n = ka_{n-1} = k(ka_{n-2}) = \dots = k^n a_0 = ck^n$.



Linear Homogeneous Recurrence Example

$$a_n = 5a_{n-1} - 6a_{n-2}, a_0 = 7, a_1 = 16$$

- Since the solution was of the form $a_n = t^n$, thus for our first attempt at finding a solution of the second-order recurrence relation, we will search for a solution of the form $a_n = t^n$.

$$- t^n = 5t^{n-1} - 6t^{n-2}$$

$$- t^2 - 5t + 6 = 0.$$



Example

- Solving the above we obtain, $t = 2, t = 3$.
- At this point, we have two solutions S and T given by
 - $S_n = 2^n, T_n = 3^n$.
- We can verify that S and T are solutions of the above, then $bS + dT$, where b and d are any numbers whatever, is also a solution of the above.



Example

- In our case, if we define the sequence U by the equation

- $U_n = bS_n + dT_n$

- $= b2^n + d3^n$

- To satisfy the initial conditions, we must have

- $7 = U_0 = b2^0 + d3^0 = b + d.$

- $16 = U_1 = b2^1 + d3^1 = 2b + 3d.$



Example

- Solving these equations for b and d , we obtain

- $b = 5, d = 2$

- Therefore, the sequence U defined by

- $U_n = 5 \times 2^n + 2 \times 3^n$

satisfies the recurrence relation and the initial conditions.



Fibonacci Sequence

- **The Fibonacci sequence is defined by the recurrence relation**

- $f_n = f_{n-1} + f_{n-2}, n \geq 3$ and initial conditions

- $f_1 = 1, f_2 = 2.$

- **We begin by using the quadratic formula to solve**

- $t^2 - t - 1 = 0$

- **The solutions are**

- $t = \frac{1 \pm \sqrt{5}}{2}$



Example

- Thus the solution is of the form

$$f_n = b\left(\frac{1+\sqrt{5}}{2}\right)^n + d\left(\frac{1-\sqrt{5}}{2}\right)^n.$$

- To satisfy the initial conditions, we must have

$$b\left(\frac{1+\sqrt{5}}{2}\right) + d\left(\frac{1-\sqrt{5}}{2}\right) = 1,$$

$$b\left(\frac{1+\sqrt{5}}{2}\right)^2 + d\left(\frac{1-\sqrt{5}}{2}\right)^2 = 2.$$



Example

- Solving these equations for b and d , we obtain

$$b = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right), d = -\frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right).$$

- Therefore, an explicit formula for the Fibonacci sequence is

$$f_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^{n+1} - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)^{n+1}.$$



Tower of Hanoi

- Find an explicit formula for a_n , the minimum number of moves in which the n-disk Tower of Hanoi puzzle can be solved.

- $a_n = 2a_{n-1} + 1, a_1 = 1.$

- Applying the iterative method, we obtain

$$a_n = 2a_{n-1} + 1$$

$$= 2(2a_{n-2} + 1) + 1$$

$$= 2^2 a_{n-2} + 2 + 1$$

$$= 2^2 (2a_{n-3} + 1) + 2 + 1$$

$$= 2^3 a_{n-3} + 2^2 + 2 + 1$$

...

$$= 2^{n-1} a_1 + 2^{n-2} + 2^{n-3} + \dots + 2 + 1$$

$$= 2^{n-1} + 2^{n-2} + 2^{n-3} + \dots + 2 + 1$$

$$= 2^n - 1$$



Common Recurrence Types

- **Decrease-by-one**

- $T(n) = T(n - 1) + f(n)$

- **Decrease-by-a-constant-factor**

- $T(n) = T(n/b) + f(n)$

- **Divide-and-conquer**

- $T(n) = aT(n/b) + f(n)$

