# Solutions for Assignment 1

#### CSCI2100B

#### October 16, 2013

# 1 Written Assignment

#### Exercise 1.1

(6)  $\sum_{i=1}^{n} ia^{i}$  **Solution:** According to the summation formula of geometric progression,  $T = \sum_{i=1}^{n} a^{i} = \frac{a^{n+1}-a}{a-1}$ . Let  $S = \sum_{i=1}^{n} ia^{i} aS = \sum_{i=1}^{n} ia^{(i+1)} = \sum_{i=1}^{n} (i-1)a^{i} + na^{n+1}$ .  $(a-1)S = na^{n+1} - T \Rightarrow S = \frac{1}{a-1}(na^{n+1} - \frac{a^{n+1}-a}{a-1})$ . (8)  $\sum_{i=0}^{n} i^{2}$ **Solution:** 

$$i^{3} - (i - 1)^{3} = 3i^{2} - 3i + 1$$

$$\Rightarrow \quad n^{3} = \sum_{i=1}^{n} (i^{3} - (i - 1)^{3}) = \sum_{i=1}^{n} (3i^{2} - 3i + 1)$$

$$= \quad 3\sum_{i=1}^{n} i^{2} - 3\frac{n(n+1)}{2} + n$$

$$\Rightarrow \quad \sum_{i=0}^{n} i^{2} = \sum_{i=1}^{n} i^{2} = \frac{n(n+1)(2n+1)}{6}$$

(12) Is  $2^{n+1} = O(2^n)$ ? Solution: Yes.  $2^{n+1} = 2 \cdot 2^n = O(2^n)$ (13) Is  $2^{2n} = O(2^n)$ ? Solution: No.  $2^{2n} = O(4^n)$ 

#### Exercise 1.2

(3) If GCD(a, b) = p and GCD(c, d) = q, is GCD(ac, bd) = pq true for all the a, b, c, d? Either prove it or give a counterexample. Solution: Counterexample: a = 5, b = 4, c = 4, d = 5.

#### Exercise 1.3

(1) T(n) = aT(n-1) + bn, T(1) = 1

**Solution:** First consider the problem of solving T(n) = aT(n-1) + b. An intuitive approach is to let S(n) = T(n) + k and then replace T(n) with S(n) in the original formula to make it S(n) = aS(n-1).

Here the problem changes to T(n) = aT(n-1) + bn. We now let S(n) = T(n) + pn + q. After replacing T(n) with S(n), we obtain S(n) - pn - q = a(S(n-1) - p(n-1) - q) + bn, which leads to S(n) = aS(n-1) + (p - ap + b)n + (q + ap - aq). Thus we let p - ap + b = q - ap - aq = 0, we can get S(n) = aS(n-1). After calculation,  $p = \frac{b}{a-1}$  and  $q = \frac{ab}{(a-1)^2}$ . Since we know S(1) according to T(1) and we have S(n) = aS(n-1), we can solve S(n) easily. Finally, we use

S(n) = T(n) + pn + q to get T(n). One thing you should take care is that when a = 1, T(n) = T(n-1) + bn. Then T(n) = $T(1) + 2b + 3b + \dots + nb = 1 + \frac{(n+2)(n-1)}{2}b.$ (5) Solve  $x_n = x_{n-1} - \frac{1}{4}x_{n-2}$ , with  $x_0 = 1, x_1 = 1/2$ . **Solution:** First solve the quadratic formula  $t^2 - t + \frac{1}{4} = 0$ . The solutions are  $t_1 = t_2 = \frac{1}{2}$ . Thus the solution is of the form  $x_n = a(\frac{1}{2})^n$ . To satisfy the initial conditions, we can obtain a = 1. Thus,  $x_n = (\frac{1}{2})^n$ .

#### Exercise 1.4

(5) Prove 2lq(n!) > nlqn by using Induction, where n is a positive integer greater than 2. **Solution:** Let P(n) be lg(n!) > nlgn, where n is a positive integer. For n = 1, L.H.S = 2lg(1!) > lg(1) = R.H.S. P(1) is true. Assume P(k) is true, i.e. 2lg(k!) > klgk, where k is a positive integer For n = k + 1, L.H.S = 2lg((k+1)!)= 2(lq(k!) + lq(k+1))> klgk + 2lg(k+1)(by assumption)  $(k+1) > e > (1+\frac{1}{k})^k \Rightarrow k^k > (k+1)^{k-1} \text{ for } k \ge 2$ > (k-1)lg(k+1) + 2lg(k+1)= (k+1)lg(k+1)= R.H.S

P(k+1) is also true.

Therefore, by M.I., P(n) is true for all positive integer n. (6) The number generated by the formula  $n^2 + n + 17$  is prime for  $n \ge 0$ , where n is an integer. Either prove it or disprove it by counterexample. **Solution:** No. Let n = 17. Then  $n^2 + n + 17 = 17 \times (17 + 1 + 1) = 17 \times 19$ 

### Exercise 1.6

(1) for i = 1 to n; for j = 1 to n;x := x + 1;

**Solution:**  $f(n) = n^2$ ,  $g(n) = n^2$ .

(3) for i = 1 to n; for j = i to n; for k = 1 to j; x := x + 1;

**Solution:**  $f(n) = \sum_{i=1}^{n} \sum_{j=i}^{n} \sum_{k=1}^{j} 1 = \sum_{i=1}^{n} \sum_{j=i}^{n} j = \frac{1}{2} \sum_{i=1}^{n} (n+i)(n-i+1) = \frac{1}{2} \sum_{i=1}^{n} (n(n+i))(n-i+1) = \frac{1}{2} \sum_{i=1}^{n} (n$ 

Solution:  $f(n) = \sum_{i=1}^{n} \sum_{j=i}^{n} \sum_{k=1}^{1000} 1 = 1000 \sum_{i=1}^{n} \sum_{j=i}^{n} 1 = 1000 \sum_{i=1}^{n} (n-i+1) = 1000 \sum_{i=1}^{n} i = 1000 \sum_{i=1}^{n} 1 = 1000 \sum_{i=1}^$  $500n(n+1), g(n) = n^2.$ 

## Exercise 1.8

(3) What is the running time of this algorithm with the following assumptions? What is its big-O notation?

Statement	Time Unit
assignment	1
+	1.25
*	1.75
for-next loop set-up	2.3
each loop	1.5

**Solution:** First assignment consumes 1 unit. Setting up for-next loop consumes 2.3 units. For each loop, there are one \*, one +, and one assignment. Thus, (1.75 + 1.25 + 1 + 1.5) \* (n + 1) = 5.5n + 5.5 units will be consumed. In total, the running time is 5.5n + 8.8 units, which is O(n) in big-O notation.

# Exercise 1.9

(2) Calculate the time and space complexity for n = 10, 20, 30, 50, 70, and 100 for each algorithm.

Solution: Just pay attention to different ranges.

(4) Come up with a strategy that you would use to minimize the time and space complexity individually?

**Solution:**  $t(n) = min\{t_A(n), t_B(n)\}, s(n) = min\{s_A(n), s_B(n)\}$ . Thus,

 $t(n) = \begin{cases} n & \text{if } 1 \le n < 50\\ n^2 & \text{if } 50 \le n < 70\\ n^3 & \text{if } 70 \le n \le 100\\ s(n) = \begin{cases} n & \text{if } 1 \le n < 20\\ 1.5n & \text{if } 20 \le n < 50\\ 0.5n & \text{if } 50 \le n \le 100 \end{cases}$ 

# 2 Programming Assignment

#### Exercise 1.17

Analysis: The function isPrime() is responsible for checking the input n is a prime. In the main function, the program uses array P to store all 50 smallest primes. Pay attention to the first for-loop in the main function. The sample code is shown below.

# Exercise 1.20

**Analysis:** All you need to care is how to read input. This problem tests your ability of dealing with strings. The sample code is shown below.

```
#include <stdio.h>
#include <stdlib.h>
int main()
{
        char str[30];
        int T, i, j;
        scanf("%d\n", &T);
        for (i = 0; i < T; i++) {
                scanf("%s", str);
                for (j = 0; j < strlen(str); j++) {</pre>
                         if (str[j] < '0' || str[j] > '9') putchar(str[j]);
                }
                putchar('\n');
        }
        return 0;
}
```