

Notes 6: Weighted Majority

1. ONLINE REGRET BOUND MODEL

e.g. stock market prediction: guessing whether it will go up or down for each day

A sequence of rounds/trials, each being:

- (1) A new unlabeled example x arrives
- (2) n experts reveal their opinions about the label for x (label is either 0 or 1)
- (3) Algorithm predicts 0 or 1 according to experts' opinions
- (4) Algorithm is told correct label for x

Goal: minimize number of mistakes, compared with the best expert

If every "expert" makes many mistakes, algorithm may, too

2. WEIGHTED MAJORITY

Weighted Majority

Fix parameter $0 \leq \beta < 1$

Initialize: $w_1 = \dots = w_n = 1$

On input x , poll opinions from experts

 Compute total weight q_0 of experts predicting 0 and total weight q_1 predicting 1

 Predict according to weighted majority (predict 0 if $q_0 > q_1$; predict 1 otherwise)

On revealing correct label, penalize incorrect experts

 Multiply every incorrect expert i 's weight w_i by β

If $\beta = 0$, Weighted Majority algorithm becomes Halving algorithm

expert	\longleftrightarrow	concept
expert i 's opinion in j th trial	\longleftrightarrow	concept c 's classification for j th sample

No longer assume any expert/concept correctly classifies all samples

Robust to classification noise

Theorem 1. For any trial sequence, if the best expert (out of n experts) makes m mistakes, then number of mistakes of Weighted Majority is at most

$$\frac{\log n + m \log(1/\beta)}{\log(\frac{2}{1+\beta})}$$

e.g. $\beta = 1/2$: $2.41(m + \log n)$

e.g. $\beta = 3/4$: $2.2m + 5.2 \log n$

e.g. $\beta = 1 - \varepsilon$: $\approx (2 + \frac{3}{2}\varepsilon)m + \frac{2}{\varepsilon} \log n$

Proof. let $W = q_0 + q_1 =$ total weight of all experts (initially n)

After each mistake, at least half of W shrinks by factor β

Total weight reduces to $\leq \frac{W}{2} + \beta \frac{W}{2} = \frac{1+\beta}{2} W$

when Weighted Majority makes M mistakes: $W \leq (\frac{1+\beta}{2})^M n$

when best expert makes m mistakes: $w_i = \beta^m$

$w_i \leq W \implies \beta^m \leq (\frac{1+\beta}{2})^M n \iff m \log \beta \leq M \log(\frac{1+\beta}{2}) + \log n$

$\iff M \log(\frac{2}{1+\beta}) \leq \log n + m \log(1/\beta)$

□

Note: The bound can be interpreted as

$$\frac{\log(W_{\text{init}}/W_{\text{final}})}{\log(1/u)} \quad \text{where } u = \frac{1+\beta}{2} = \text{shrink in } W \text{ per mistake}$$

3. RANDOMIZED WEIGHTED MAJORITY

Randomized Weighted Majority

Fix parameter $0 \leq \beta < 1$

Initialize: $w_1 = \dots = w_n = 1$

On input x , poll opinions from experts

Predict according to a random expert i chosen with probability proportional to w_i
i.e. probability w_i/W , where $W = \text{total weight} = \sum_{1 \leq i \leq n} w_i$

On revealing correct label, penalize incorrect experts

Multiply every incorrect expert i 's weight w_i by β

Denote $\varepsilon = 1 - \beta$

Theorem 2. *Given any trial sequence with fixed correct labels, if the best expert (out of n experts) makes m mistakes, then*

$$\mathbb{E}[\#\text{mistakes of RWM}] \leq \frac{\ln n - m \ln(1 - \varepsilon)}{\varepsilon}$$

e.g. $\beta = 1/2$: $1.39m + 2 \ln n$

e.g. $\beta = 3/4$: $1.16m + 4 \ln n$

e.g. $\beta = 1 - \varepsilon$: $\approx (1 + \frac{\varepsilon}{2})m + \frac{1}{\varepsilon} \ln n$

Key benefit: $\approx m$ mistakes (ignoring additive $\log n$), down from $\approx 2m$

Proof. Fix any sequence of T trials together with their correct labels

Let $F_t =$ fraction of total weight on wrong prediction at trial t

Want to bound $\mathbb{E}[\#\text{mistakes of RWM}] = \sum_{1 \leq t \leq T} F_t$

At trial t , probability of mistake is F_t , and εF_t fraction of weight is removed

$$W_{\text{final}} = W_{\text{init}}(1 - \varepsilon F_1) \dots (1 - \varepsilon F_T) \quad (W_{\text{init}} = n)$$

$$\ln W_{\text{final}} = \ln n + \ln(1 - \varepsilon F_1) + \dots + \ln(1 - \varepsilon F_T)$$

Best expert makes m mistakes: $w_i = \beta^m = (1 - \varepsilon)^m$

$$W_{\text{final}} \geq w_i \iff \ln W_{\text{final}} \geq \ln w_i \iff \ln n + \sum_{1 \leq t \leq T} \ln(1 - \varepsilon F_t) \geq m \ln(1 - \varepsilon)$$

Claim: $\ln(1 - x) \leq -x$ for all $x < 1$

Take $x = \varepsilon F_t$ in Claim, we get $\ln(1 - \varepsilon F_t) \leq -\varepsilon F_t$, and

$$\varepsilon \sum_{1 \leq t \leq T} F_t \leq \sum_{1 \leq t \leq T} -\ln(1 - \varepsilon F_t) \leq \ln n - m \ln(1 - \varepsilon) \quad \square$$

Above Claim is true because for all real x $1 - x \leq e^{-x}$

