Nondeterministic Finite Automata

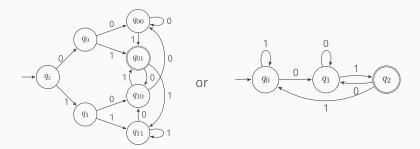
CSCI 3130 Formal Languages and Automata Theory

Siu On CHAN Fall 2019

Chinese University of Hong Kong

Example from last lecture with a simpler solution

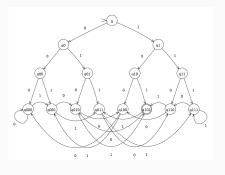
Construct a DFA over {0,1} that accepts all strings ending in 01



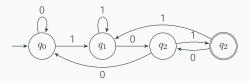
Three weeks later: DFA minimization

Another example from last lecture

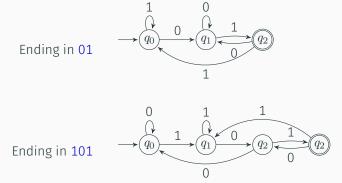
Construct a DFA over $\{0,1\}$ that accepts all strings ending in 101



or



String matching DFAs

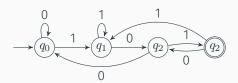


Fast string matching algorithms to turn a pattern into a string matching DFA and execute the DFA:

Boyer-Moore (BM) and Knuth-Morris-Pratt (KMP) (won't cover in class)

Nondeterminism

In a few lectures

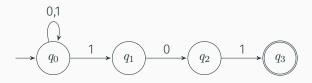


What problems can finite state machines solve?

We'll answer this question in the next few lectures Useful to consider hypothetical machines that are nondeterministic

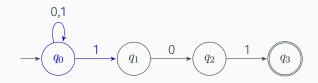
Nondeterministic finite automata

A machine that is nondeterministic (and effectively making guesses)



Each state can have zero, one, or more outgoing transitions labeled by the same symbol

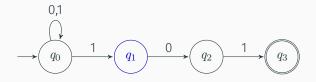
Choosing where to go



State q_0 has two transitions labeled 1

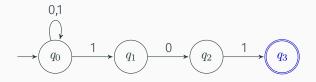
Upon reading 1, we have the choice of staying at \emph{q}_0 or moving to \emph{q}_1

Ability to choose



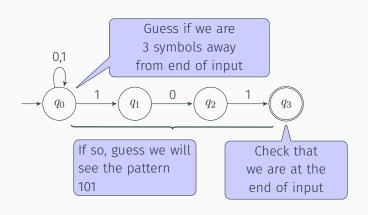
State q_1 has no transition labeled 1 Upon reading 1 at q_1 , die; upon reading 0, continue to q_2

Ability to choose

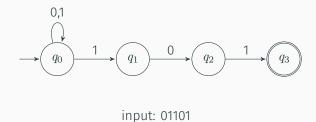


State q_1 has no transition going out Upon reading 0 or 1 at q_3 , die

Meaning of NFA



How to run an NFA



The NFA can have several active states at the same time NFA accepts if at the end, one of its active states is accepting

Example

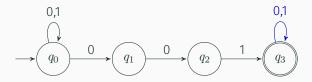
Construct an NFA over alphabet {0,1} that accepts all strings containing the pattern 001 somewhere

11001010, 001001, 111001 ε , 000, 010101

should be accepted should not

Example

Construct an NFA over alphabet {0,1} that accepts all strings containing the pattern 001 somewhere



Definition

A nondeterministic finite automaton (NFA) is a 5-tuple $(Q, \Sigma, \delta, q_0, F)$ where

- $\cdot Q$ is a finite set of states
- \cdot Σ is an alphabet
- $\delta: Q \times (\Sigma \cup \{\varepsilon\}) \rightarrow \text{subsets of } Q \text{ is a transition function}$
- $q_0 \in Q$ is the initial state
- $F \subseteq Q$ is a set of accepting states

Differences from DFA:

- transition function δ can go into several states
- allows ε -transitions

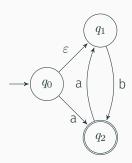
Language of an NFA

The NFA accepts string x if there is some path that, starting from q_0 , ends at an accepting state as x is read from left to right

The language of an NFA is the set of all strings accepted by the NFA

ε -transitions

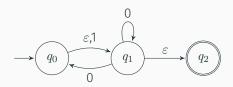
ε -transitions can be taken for free:



accepts a, b, aab, bab, aabab, ...

rejects ε , aa, ba, bb, ...

Example

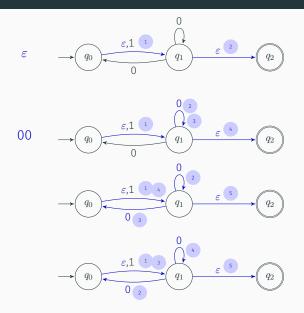


alphabet $\Sigma=\{0,1\}$ states $Q=\{q_0,q_1,q_2\}$ initial state q_0 accepting states $F=\{q_2\}$

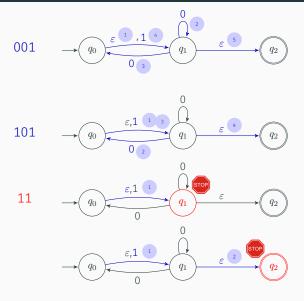
table of transition function δ

		inputs		
		0	1	ε
states	q_0	Ø	$\{q_1\}$	$\{q_1\}$
	q_1	$\{q_0,q_1\}$	Ø	$\{q_2\}$
	q_2	Ø	Ø	Ø

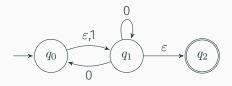
Some computational paths of the NFA



Some computational paths of the NFA



Language of this NFA



What is the language of this NFA?

Example of ε -transitions

Construct an NFA that accepts all strings with an even number of 0s or an odd number of 1s

Example of ε -transitions

Construct an NFA that accepts all strings with an even number of 0s or an odd number of 1s

