

Notes 20: Sampling spanning trees by random walk

1. RANDOM SPANNING TREES

Fix a connected graph $G = (V, E)$ on n vertices. A spanning tree in G is an acyclic subgraph of G containing $n - 1$ edges.

We want to sample a spanning tree of G , (nearly) uniformly at random, as follows.

Random walk on spanning trees

Let T_0 be an arbitrary spanning tree of G

For $t = 0, 1, 2, \dots$

Remove an edge from T_t uniformly at random to obtain F_t

Among all spanning trees containing F_t , uniformly pick one as the new T_{t+1}

This is a random walk/Markov chain on an auxiliary weighted graph \mathcal{T}_G , whose nodes are spanning trees in G , and two nodes in \mathcal{T}_G are adjacent if they share exactly $n - 2$ edges.

For decades, this random walk was conjectured to mix in polynomial time. It was recently proved by Anari, Liu, Oveis Gharan, and Vintzart.

Theorem 1.1. *The above random walk has eigenvalue gap at least $1/(n - 1)$.*

Eigenvalue gap β is the difference $\lambda_1 - \lambda_2$ between the two largest eigenvalues. By results in Notes12, the lazy version of the random walk mixes in polynomial time. Recall that the lazy random walk mixes in time $O((\log |V(\mathcal{T}_G)|)/\beta)$. Since $|V(\mathcal{T}_G)| \leq \binom{n}{n-1} \leq \binom{n^2}{n-1} \leq n^{2(n-1)} = \exp(O(n \log n))$, the lazy random walk mixes in time $O(n^2 \log n)$.

2. BIPARTITE INCLUSION GRAPHS

Since the ambient graph G is fixed, we identify a spanning tree T in G with the set of $n - 1$ edges in T .

Denote by $Y(n - 1)$ the sets of edges of spanning trees in G . More generally, define

$$Y(k) = \left\{ F \in \binom{E}{k} \mid F \text{ is acyclic} \right\} \quad \text{for } 0 \leq k \leq n - 1.$$

Decompose the random walk over spanning trees into two transitions:

- (Down) Go from $T_t \in Y(n - 1)$ to $F_t \in Y(n - 2)$ by removing an edge uniformly at random
- (Up) Then from F_t to $T_{t+1} \in Y(n - 1)$ by choosing T_{t+1} uniformly from among all $T_{t+1} \supset F_t$

Down transition corresponds to matrix D_{n-1} , which is $Y(n - 1)$ -by- $Y(n - 2)$ (recall that a row probability vector multiplies on the left to a transition matrix). Likewise up transition corresponds to matrix U_{n-1} , which is $Y(n - 2)$ -by- $Y(n - 1)$.

$D_{n-1}U_{n-1}$ is the transition matrix for the random walk. We want to bound $\beta(D_{n-1}U_{n-1})$.

We think of up and down transitions as random walk transitions on the auxiliary graph Γ_{n-1} :

Definition 2.1 (Bipartite inclusion graph). For $0 \leq k \leq n - 1$, Γ_k has vertex set $Y(k) \cup Y(k - 1)$. $F \in Y(k)$ is adjacent to $F' \in Y(k - 1)$ if $F \supset F'$.

We will look at Γ_k for $0 \leq k \leq n - 1$ later on to apply induction.

In Γ_k , down transition means moving from $F \in Y(k)$ to a random neighbor $F' \in Y(k - 1)$, uniformly from among all k neighbors of F .

We will define up transition matrix U_k shortly, to move from $F' \in Y(k - 1)$ to a random neighbor $F \in Y(k)$ from some distribution.

Thanks to the following proposition, $D_{n-1}U_{n-1}$ and $U_{n-1}D_{n-1}$ have the same eigenvalue gap, whenever this gap is at most 1.

Proposition 2.2. *Given any matrices A and B , AB and BA have the same non-zero eigenvalues with the same multiplicities.*

This proposition can be proved by showing AB and BA have essentially the same characteristic polynomials. Look up “Characteristic polynomials” on Wikipedia if interested.

Kaufman and Oppenheim came up with a way to relate $\beta(U_{n-1}D_{n-1})$ to $\beta(D_{n-2}U_{n-2})$. Of course, $\beta(D_{n-2}U_{n-2}) = \beta(U_{n-2}D_{n-2})$. One can then inductively bound $\beta(U_k D_k) = \beta(D_k U_k)$.

Proposition 2.3 (Kaufman–Oppenheim). $\beta(D_k U_k) \geq 1/k$ for $1 \leq k \leq n - 1$.

This main proposition implies the main theorem.

3. RANDOM WALK ON BIPARTITE GRAPH

Up and down transitions U_k and D_k in Γ_k are special cases of random walk on bipartite graph.

3.1. General bipartite graph. Consider bipartite graph Γ on vertex set $T \cup B$ (“top” and “bottom”) with weight μ over its edges.

Further assume μ is a distribution: nonnegative weight on edges summing to 1.

Choosing an edge $(t, b) \in T \times B$ from μ induces marginal distribution π_T on T , and marginal π_B on B .

The marginal probability $\pi_T(t)$ coincides with the degree of $t \in T$ (sum of edge weights incident to t). Similarly $\pi_B(b)$ is the degree of $b \in B$.

The simple random walk on Γ with edge weights μ has transition matrix $\begin{pmatrix} O & D \\ U & O \end{pmatrix}$.

D denotes (“down”) transition from T to B ; U denotes (“up”) transition from B to T .

Distribution $\pi = (\frac{1}{2}\pi_T, \frac{1}{2}\pi_B)$ over $T \cup B$ is stationary for the random walk, because its probability masses are proportional to vertex degrees.

$$\left(\frac{1}{2}\pi_T^\top \quad \frac{1}{2}\pi_B^\top\right) = \left(\frac{1}{2}\pi_T^\top \quad \frac{1}{2}\pi_B^\top\right) \begin{pmatrix} O & D \\ U & O \end{pmatrix} = \left(\frac{1}{2}\pi_B^\top U \quad \frac{1}{2}\pi_T^\top D\right), \text{ so}$$

$$\pi_T^\top D = \pi_B^\top \quad \text{and} \quad \pi_B^\top U = \pi_T^\top.$$

3.2. Bipartite inclusion graph Γ_k . Then $T = Y(k)$ and $B = Y(k - 1)$.

Edge weight $\mu = \mu_k$ will be chosen later so that $D = D_k$. That is, random walk from T to B according to μ_k corresponds to uniformly dropping an element from $t \in T$.

We want transition probabilities $\mu_k(t, b)/\pi_T(t)$ from t to any of its neighbor b to be $1/k$, so

$$\mu_k(t, b) = \frac{\pi_T(t)}{k} \quad \text{for every edge } (t, b) \text{ in } \Gamma_k.$$

U_k will be up transition matrix U in Γ_k , i.e. random walk from B to T according to μ_k .

U_k coincides with this transition: Given $b \in B$, choose neighbor t of b with probability proportional to $\pi_T(t)$. This is because

$$U_k(b, t) = \frac{\mu_k(t, b)}{\pi_B(b)} = \frac{1}{\pi_B(b)} \frac{\pi_T(t)}{k} \propto \pi_T(t).$$

4. n -BIPARTITE INCLUSION GRAPH

We can visualize using the n -bipartite inclusion graph $\Gamma_{n-1} \cup \dots \cup \Gamma_1$.

We now define edge weights μ_k 's for every layer so that

- (A) Every μ_k is a distribution over edges in Γ_k
- (B) Down transition in Γ_k according to μ_k is D_k
- (C) The top marginal of μ_{n-1} is the uniform over $Y(n - 1)$
- (D) Top marginal distribution of μ_{k-1} in Γ_{k-1} equals bottom marginal distribution of μ_k in Γ_k (so as to relate $\beta(U_k D_k)$ in Γ_k to $\beta(D_{k-1} U_{k-1})$ in Γ_{k-1})

This is done by the following random process of choosing a path (F_{n-1}, \dots, F_0) on $\Gamma_{n-1} \cup \dots \cup \Gamma_1$:

- (1) Choose $F_{n-1} \in Y(n - 1)$ uniformly at random
- (2) For k from $n - 1$ to 1, drop a uniformly random edge in F_k to get $F_{k-1} \in Y(k - 1)$

F_k in this random path has marginal distribution π_k over $Y(k)$. π_k is both the top marginal of μ_k in Γ_k and bottom marginal of μ_{k+1} in Γ_{k+1} .

The joint marginal of $(F_k, F_{k-1}) \in Y(k) \times Y(k-1)$ induces edge distribution μ_k for Γ_k .

$\pi_k^\top D_k = \pi_{k-1}^\top$ and $\pi_{k-1}^\top U_k = \pi_k^\top$, so both $D_k U_k$ and $U_{k+1} D_{k+1}$ have stationary distribution π_k .

π_{n-1} is uniform over $Y(n-1)$. U_{n-1} defined here agrees with §2 since $U_{n-1}(b, t) \propto \pi_{n-1}(t) \propto 1$.

For some applications, one may want to sample from a nonuniform target distribution π over spanning trees in $Y(n-1)$. For example, G may have edge weights w , and a natural distribution π on spanning trees T would be proportional to the product of edge weights in T , so that $\pi(T) \propto \prod_{e \in T} w(e)$.

We need to change the edge weights μ_k , so that they now satisfy

(C') The top marginal of μ_{n-1} is π

This can be done by changing the first step of the random process of the path (F_{n-1}, \dots, F_0) :

(1') Choose $F_{n-1} \in Y(n-1)$ from π

Step (2) remains the same. Changing initial distribution π in step (1) will affect edge distributions μ_k in Γ_k and marginals π_k on $Y(k)$.

Everything in our analysis still holds given any target distribution π over $Y(n-1)$.

5. UP-WALK AND DOWN-WALK

In Γ_k , layer k of the n -partite inclusion graph, there are two natural two-step walks:

- Up-walk $P_{k-1}^\wedge = U_k D_k$ on a weighted graph on vertex set $Y(k-1)$
- Down-walk $P_k^\vee = D_k U_k$ on a weighted graph on vertex set $Y(k)$

Focusing on Γ_k , abbreviate

$$T = Y(k) \quad B = Y(k-1) \quad \mu = \mu_k$$

5.1. Down-walk. $P_k^\vee = D_k U_k$ induces a random walk on a weighted graph on T .

A step in P_k^\vee also corresponds to a length-2 walk in Γ_k , from T to B to T .

$$P_k^\vee(t, t') = \begin{cases} \sum_{b \in B, b \subset t} \frac{\mu(t, b)}{\pi_k(t)} \frac{\mu(t', b)}{\pi_{k-1}(b)} & t = t' \\ \frac{\mu(t, b)}{\pi_k(t)} \frac{\mu(t', b)}{\pi_{k-1}(b)} & t \cap t' = b \in B \\ 0 & \text{otherwise} \end{cases} .$$

5.2. Up-walk. $P_{k-1}^\wedge = U_{k-1} D_{k-1}$ corresponds to the random walk on a weighted graph on B .

A step in P_{k-1}^\wedge also corresponds to a length-2 walk in Γ_k , from B to T to B .

$$P_{k-1}^\wedge(b, b') = \begin{cases} 1/k & b = b' \\ \frac{\mu(t, b)}{\pi_{k-1}(b)} \frac{\mu(t, b')}{\pi_k(t)} & b \cup b' = t \in T \\ 0 & \text{otherwise} \end{cases} .$$

5.3. Non-lazy up-walk. \tilde{P}_{k-1}^\wedge is the non-lazy version of P_{k-1}^\wedge .

A step in \tilde{P}_{k-1}^\wedge also corresponds to a length-2 path in Γ_k , from B to T to a different vertex in B .

$$P_{k-1}^\wedge = \frac{1}{k} I + \frac{k-1}{k} \tilde{P}_{k-1}^\wedge .$$

π_{k-1} is stationary for both P_{k-1}^\wedge and \tilde{P}_{k-1}^\wedge , as a common left-eigenvector with eigenvalue 1.

6. SPECTRAL COMPARISON

We now upperbound the spectrum of \tilde{P}_k^\wedge by the spectrum of P_k^\vee . It is more convenient to first transform each of them into a symmetric matrix with the same spectrum.

\tilde{P}_k^\wedge and P_k^\vee share a common stationary distribution π_k . These transitions are both of the form $P(b, b') = A_P(b, b')/\pi_k(b)$ for symmetric A_P , so $P = \Pi^{-1}A_P$, where $\Pi = \text{Diag}(\pi_k)$. They both represent a random walk with adjacency matrix $A_P = \Pi P$ and common degree matrix Π .

Definition 6.1. Given matrices P, Q, Π , if $\Pi P, \Pi Q$ are symmetric, we write

$$P \preceq_{\Pi} Q \iff \Pi P \preceq \Pi Q.$$

$P \preceq_{\Pi} Q$ is equivalent to $\Pi^{1/2}P\Pi^{-1/2} \preceq \Pi^{1/2}Q\Pi^{-1/2}$, if Π is symmetric and $\Pi \succ 0$.

Since P is similar to the symmetric matrix $\mathcal{A}_P = \Pi^{1/2}P\Pi^{-1/2}$ (the normalized adjacency matrix), they have the same spectra.

In our application, P and Q often represent transitions with a common stationary distribution (the main diagonal of Π). \mathcal{A}_P and \mathcal{A}_Q will have the same top eigenspace (spanned by $\mathbb{1}$). Applying Courant–Fischer to the orthogonal subspace,

$$P \preceq_{\Pi} Q \iff \mathcal{A}_P \preceq \mathcal{A}_Q \implies \lambda_2(\mathcal{A}_P) = \sup_{x \perp \mathbb{1}, \|x\|=1} x^\top \mathcal{A}_P x \leq \sup_{x \perp \mathbb{1}, \|x\|=1} x^\top \mathcal{A}_Q x = \lambda_2(\mathcal{A}_Q),$$

and hence $\lambda_2(P) = \lambda_2(\mathcal{A}_P) \leq \lambda_2(\mathcal{A}_Q) = \lambda_2(Q)$.

Proposition 6.2. For $1 \leq k \leq n - 2$, let $\Pi = \text{Diag}(\pi_k)$. Then $\tilde{P}_k^\wedge \preceq_{\Pi} P_k^\vee$.

This proposition is proved in the next section. We now show how it implies [Proposition 2.3](#).

Proof of [Proposition 2.3](#). $\lambda_1(P_k^\vee) = 1$. We prove by induction that $\lambda_2(P_k^\vee) \leq 1 - \frac{1}{k} = \frac{k-1}{k}$.

When $k = 1$, P_k^\vee has rank 1, so $\lambda_2(P_k^\vee) = 0 \leq \frac{k-1}{k}$.

For $k > 1$, $\tilde{P}_{k-1}^\wedge \preceq_{\Pi} P_{k-1}^\vee$ by [Proposition 6.2](#), so $\lambda_2(\tilde{P}_{k-1}^\wedge) \leq \lambda_2(P_{k-1}^\vee) \leq \frac{k-2}{k-1}$.

Also $\lambda_2(P_{k-1}^\wedge) = \frac{1}{d} + \frac{k-1}{k} \lambda_2(\tilde{P}_{k-1}^\wedge) \leq \frac{k-1}{k}$.

$P_{k-1}^\wedge = U_k D_k$ and $P_k^\vee = D_k U_k$ share the same non-zero eigenvalues, thus $\lambda_2(P_k^\vee) \leq \frac{k-1}{k}$. \square

7. GARLAND'S METHOD

We now discuss [Proposition 6.2](#) that bounds \tilde{P}_k^\wedge by P_k^\vee .

Focusing on adjacent layers $\Gamma_{k+1} \cup \Gamma_k$, abbreviate

$$T = Y(k+1) \quad M = Y(k) \quad B = Y(k-1) \quad (\text{“top” “middle” “bottom”})$$

Garland method decomposes transitions of P_k^\vee and \tilde{P}_k^\wedge into unions of subgraphs.

7.1. Down-walk. A step in $P_k^\vee = D_k U_k$ represents a length-2 walk (m, b, m') from M to B to M . Decompose $P_k^\vee = \sum_{b \in B} P_b^\vee$ based on the bottom vertex b of (m, b, m') .

P_b^\vee also corresponds to transitions on a weighted clique (with self-loops) over $S_b \subseteq M$, where

$$S_b = \{m \in M \mid m \supset b\}.$$

7.2. Non-lazy up-walk. A step in \tilde{P}_k^\wedge represents a length-2 path (m, t, m') from M to T to a different vertex in M . Decompose $\tilde{P}_k^\wedge = \sum_{b \in B} \tilde{P}_b^\wedge$ based on the common intersection $b = m \cap m'$ of this path.

\tilde{P}_b^\wedge also corresponds to transitions on the weighted graph $H_b = (S_b, E_b)$ over S_b with edge set

$$E_b = \{(m, m') \in S_b \times S_b \mid m \cup m' \in T\}.$$

7.3. Spectra. The down-walk and non-lazy up-walk are now decomposed into subgraphs on S_b , over various $b \in B$.

For P_b^\vee , the subgraph (with adjacency matrix ΠP_b^\vee) is a clique with self-loops.

For \tilde{P}_b^\wedge , the subgraph H_b (with adjacency matrix $\Pi \tilde{P}_b^\wedge$) has its edges determined by $Y(k+1)$.

Kaufman and Oppenheim proved that **Proposition 6.2** still holds “when restricted to these two subgraphs on S_b ”.

Lemma 7.1. *For $1 \leq k \leq n-2$ and $b \in Y(k-1)$, let $\Pi = \text{Diag}(\pi_k)$. Then $\tilde{P}_b^\wedge \preceq_\Pi P_b^\vee$.*

This lemma is proved in the next lecture. Summing over $b \in Y(k-1)$ yields **Proposition 6.2**.

For P_b^\vee : It has rank 1 since $\Pi P_b^\vee = \mu_b \mu_b^\top / \pi_{k-1}(b)$ (μ_b is the vector of edge weights incident to b). Therefore $\lambda_2(P_b^\vee) = \dots = \lambda_n(P_b^\vee) = 0$, and ΠP_b^\vee is a clique that mixes perfectly in one step.

For \tilde{P}_b^\wedge : Turns out \tilde{P}_b^\wedge and P_b^\vee share the same top eigenvector with the same eigenvalue.

Therefore $\tilde{P}_b^\wedge \preceq_\Pi P_b^\vee$ is equivalent to $\lambda_2(\tilde{P}_b^\wedge) \leq 0$.

In particular \tilde{P}_b^\wedge must be a weighted expander.

8. VARIATIONS

Suppose you want apply the same random walk algorithm to a different setting, such as

- Uniformly sample a path of length d in a graph G ; or
- Uniformly sample a clique of size d in a graph G .

Does the same random walk mix quickly in these settings?

All the constructions still make sense in those settings ($(k+1)$ -partite inclusion graph, up- and down-walks, Garland’s decomposition). But **Lemma 7.1** may or may not hold.

Turns out P_b^\vee from the previous section will still has rank 1 and represents a perfectly mixing weighted clique with $\lambda_2(P_b^\vee) = \dots = \lambda_n(P_b^\vee) = 0$.

But now \tilde{P}_b^\wedge may not represent weighted expanders and may violate $\lambda_2(\tilde{P}_b^\wedge) \leq 0$, because the subgraphs H_b depend crucially on $Y(k+1)$ (and also indirectly on distribution π at the top layer).

They might only satisfy the weaker bound $\lambda_2(\tilde{P}_b^\wedge) \leq \alpha$ for some $0 \leq \alpha \leq 1$, depending on the situation. So \tilde{P}_b^\wedge are only α -expanders. This α upperbounds the error when you approximate the clique of P_b^\vee by the α -expander of \tilde{P}_b^\wedge . We will discuss this general situation in the next lecture.

If you can only show $\lambda_2(\tilde{P}_b^\wedge) \leq \alpha$ for all $b \in Y(k)$ and $1 \leq k \leq d-1$, then the conclusion of **Proposition 6.2** is weakened to be $\tilde{P}_k^\wedge \preceq_\Pi P_k^\vee + \alpha I$. And in the proof of **Proposition 2.3**, you get the weaker bound $\lambda_2(P_k^\vee) \leq 1 - \frac{1}{k} + \alpha(k-1)$, since the error accumulates in the induction. This might still give you something useful if α is tiny (at most roughly $1/k^2$).

With a more careful analysis, Alev et al can get a nontrivial upperbound on $\lambda_2(P_k^\vee)$ even when $\alpha = \Theta(1/k)$.

9. MATROID

One general situation in which **Lemma 7.1** holds (and hence fast mixing) is when the set system $Y(0) \cup \dots \cup Y(d)$ is a matroid.

A matroid is a family \mathcal{I} of subsets over a ground set U that is:

- (1) Nonempty: $\mathcal{I} \neq \emptyset$
- (2) Downward closed: If $A \in \mathcal{I}$ and $B \subseteq A$, then $B \in \mathcal{I}$
- (3) Exchangable: If $A, B \in \mathcal{I}$ and $|A| > |B|$, then there is $e \in A \setminus B$ such that $B \cup \{e\} \in \mathcal{I}$

Turns out all maximal $A \in \mathcal{I}$ have the same size d (called the rank of the matroid).

The family of acyclic edges in a graph G is an example of a matroid. The ground set U is the set E of edges in G . A subset $F \subseteq U$ belongs to \mathcal{I} if F is acyclic. Maximal $F \in \mathcal{I}$ are spanning trees in G . This matroid has rank $n-1$, where n is the number of vertices in G . Fast mixing over spanning trees is thus a special case.