

SQL 3: Nesting in Where and Having

Yufei Tao

Department of Computer Science and Engineering
Chinese University of Hong Kong

SQL offers several constructs beyond relational algebra to allow users to write more powerful queries. In the previous lecture, we have studied a collection of constructs for statistical analysis. This lecture will introduce another collection for enhancing the functionality of WHERE and HAVING.

We will first discuss WHERE and then HAVING.

Recall that the where clause contains conditions of the form $A \text{ op } v$ where A is an attribute/value, op is an arithmetic operator (e.g., $<$), and v is another attribute/value.

Next, we will learn new conditions where

- v is an SQL statement
- op is an operator that compares a value to the result of v .

In

(A_1, \dots, A_n) in ([an SQL statement])

where each A_i is an attribute/value. (A_1, \dots, A_n) must obey the schema of the table T returned by the SQL statement.

The expression returns:

- true, if tuple (A_1, \dots, A_n) appears in T .
- false, otherwise.

The bracket of “ (A_1, \dots, A_n) ” can be omitted if $n = 1$.

PROF

pid	name	dept	rank	sal
<i>p1</i>	Adam	CS	asst	6000
<i>p2</i>	Bob	EE	asso	8000
<i>p3</i>	Calvin	CS	full	10000
<i>p4</i>	Dorothy	EE	asst	5000
<i>p5</i>	Emily	EE	asso	8500
<i>p6</i>	Frank	CS	full	9000

select pid from PROF
where dept in (select dept from PROF where sal >= 10000)

Result:

pid
<hr/>
<i>p1</i>
<hr/>
<i>p3</i>
<hr/>
<i>p6</i>

PROF

pid	name	dept	rank	sal
<i>p1</i>	Adam	CS	asst	6000
<i>p2</i>	Bob	EE	asso	8000
<i>p3</i>	Calvin	CS	full	10000
<i>p4</i>	Dorothy	EE	asst	5000
<i>p5</i>	Emily	EE	asso	8500
<i>p6</i>	Frank	CS	full	9000

select pid from PROF
 where (dept, rank) in (select dept, rank from PROF where sal >= 10000)

Result:

pid
<i>p3</i>
<i>p6</i>

PROF

pid	name	dept	rank	sal
p1	Adam	CS	asst	6000
p2	Bob	EE	asso	8000
p3	Calvin	CS	full	10000
p4	Dorothy	EE	asst	5000
p5	Emily	EE	asso	8500
p6	Frank	CS	full	9000

select pid from PROF
 where (dept, sal) in (select dept, rank from PROF where sal >= 10000)

Error! (dept, sal) does not match the schema of the table returned by the nested SQL statement.

Not in

(A_1, \dots, A_n) not in ([an SQL statement])

where each A_i is an attribute/value. (A_1, \dots, A_n) must obey the schema of the table T returned by the SQL statement.

The expression returns:

- true, if tuple (A_1, \dots, A_n) does not appear in T .
- false, otherwise.

The bracket of " (A_1, \dots, A_n) " can be omitted if $n = 1$.

PROF

pid	name	dept	rank	sal
<i>p1</i>	Adam	CS	asst	6000
<i>p2</i>	Bob	EE	asso	8000
<i>p3</i>	Calvin	CS	full	10000
<i>p4</i>	Dorothy	EE	asst	5000
<i>p5</i>	Emily	EE	asso	8500
<i>p6</i>	Frank	CS	full	9000

select pid from PROF
 where (dept, rank) not in
 (select dept, rank from PROF where sal >= 10000)

Result:

pid
<i>p1</i>
<i>p2</i>
<i>p4</i>
<i>p5</i>

Set Comparison 1

Some

$A > \text{some} ([\text{an SQL statement}])$

where A is an attribute/value, and must obey the schema of the table T returned by the SQL statement.

The expression returns:

- true, if A is greater than **at least one** tuple in T .
- false, otherwise.

PROF

pid	name	dept	rank	sal
<i>p1</i>	Adam	CS	asst	6000
<i>p2</i>	Bob	EE	asso	8000
<i>p3</i>	Calvin	CS	full	10000
<i>p4</i>	Dorothy	EE	asst	5000
<i>p5</i>	Emily	EE	asso	8500
<i>p6</i>	Frank	CS	full	9000

select pid from PROF
where sal > some (select sal from PROF where dept = 'CS')

Result:

pid
<i>p2</i>
<i>p3</i>
<i>p5</i>
<i>p6</i>

All

$A > \text{all} ([\text{an SQL statement}])$

where A is an attribute/value, and must obey the schema of the table T returned by the SQL statement.

The expression returns:

- true, if A is greater than **all** tuples in T .
- false, otherwise.

PROF

pid	name	dept	rank	sal
<i>p1</i>	Adam	CS	asst	6000
<i>p2</i>	Bob	EE	asso	8000
<i>p3</i>	Calvin	CS	full	10000
<i>p4</i>	Dorothy	EE	asst	5000
<i>p5</i>	Emily	EE	asso	8500
<i>p6</i>	Frank	CS	full	9000

select pid from PROF
where sal > all (select sal from PROF where dept = 'EE')

Result:

pid
<i>p3</i>
<i>p6</i>

The operator $>$ in “ $>$ some (all)” can be replaced with $<$, $<=$, $=$, $<>$ and $>=$.

The semantics in each case agrees with the literal meaning in English. For example, “ $<$ some” means “smaller than some element (in a table)”.

Next, we will learn two more useful conditions of the form:

op ([an SQL statement])

Exists

exists ([an SQL statement])

The expression returns:

- true, if the SQL statement returns a table with **at least one** tuple.
- false, otherwise.

not exists ([an SQL statement])

The expression returns:

- true, if the SQL statement returns an **empty** table.
- false, otherwise.

PROF

pid	name	dept	rank	sal
p1	Adam	CS	asst	6000
p2	Bob	EE	asso	8000
p3	Calvin	CS	full	10000
p4	Dorothy	EE	asst	5000
p5	Emily	EE	asso	8500
p6	Frank	CS	full	9000

select pid from PROF as P
 where not exists
 (select sal from PROF
 where sal > P.sal)

Note

Observe how P is used in the nested query.

Result:

pid
p3

T_1	
pid	cid
p1	c1
p1	c2
p1	c3
p2	c2
p2	c3
p3	c1
p4	c1
p4	c2
p4	c3

T_2
cid
c1
c2
c3

select distinct pid from T1 as P
 where not exists (
 (select cid from T2)
 minus
 (select cid from T1 where pid = P.pid))

Result:

pid
p1
p4

Note that this is another way to do division in SQL.

Next we extend the above discussion to HAVING.

```
select  $A_1, \dots, A_t, agg_1(B_1), \dots, agg_m(B_m)$   
from  $T_1, \dots, T_n$   
where  $P$   
group by  $C_1, \dots, C_g$   
having  $H$ 
```

where

- C_1, \dots, C_g are called **group-by attributes**.
- H is a **group predicate**.

The group predicate can contain conditions of the form:

$$\text{agg}(A) \text{ op (SQL statement)}$$

where

- *agg* is an aggregate function
- *op* can be
 - in, not in
 - < some/all
 - <= some/all
 - = some/all
 - <> some/all
 - > some/all
 - >= some/all
- The nested SQL statement must return a table of a single numeric column.

PROF

pid	name	dept	rank	sal
<i>p1</i>	Adam	CS	asst	6000
<i>p2</i>	Bob	EE	asso	8000
<i>p3</i>	Calvin	CS	full	10000
<i>p4</i>	Dorothy	EE	asst	5000
<i>p5</i>	Emily	EE	asso	8500
<i>p6</i>	Frank	CS	full	9000

select dept from PROF
 group by dept
 having avg(sal) >= some (select avg(sal) from PROF)

Result:

dept
cs

PROF

pid	name	dept	rank	sal
p1	Adam	CS	asst	6000
p2	Bob	EE	asso	8000
p3	Calvin	CS	full	10000
p4	Dorothy	EE	asst	5000
p5	Emily	EE	asso	8500
p6	Frank	CS	full	9000

select dept from PROF
 group by dept
 having avg(sal) >= some (select sal from PROF)

Result:

dept
cs
ee