

Quick Sort

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Today, we will discuss another sorting algorithm named **quick sort**. It is a randomized algorithm that runs in $O(n^2)$ time in the **worst** case but $O(n \log n)$ time **in expectation**.

Recall:

The Sorting Problem

Problem Input:

A set S of n integers is given in an array A of length n .

Goal:

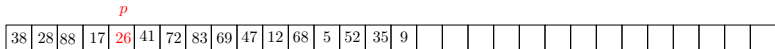
Produce an array that stores the elements of S in ascending order.

Quick Sort

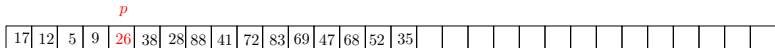
- 1 Pick an integer p in A **uniformly at random**, which is called the **pivot**.
- 2 Re-arrange the integers in an array A' such that
 - all the integers **smaller** than p are **before** p in A' ;
 - all the integers **larger** than p are **after** p in A' .
- 3 Sort the part of A' before p recursively (a subproblem).
- 4 Sort the part of A' after p recursively (a subproblem).

Example

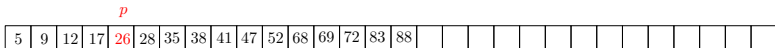
After Step 1 (suppose that 26 was randomly picked as the pivot):



After Step 2:



After Steps 3 and 4:



Analysis of Quick Sort

Quick sort is not attractive in the worst case: its worst case time is $O(n^2)$ (why?). However, quick sort is fast **in expectation**: we will prove that its expected time is $O(n \log n)$. Remember: this holds on **every** input array A .

The rest of the slides will not be tested for CSCI2100.

Analysis of Quick Sort

First, convince yourself that it suffices to analyze the number X of comparisons. The running time is bounded by $O(n + X)$.

Next, we will prove that $\mathbf{E}[X] = O(n \log n)$.

Analysis of Quick Sort

Denote by e_i the i -th smallest integer in S . Consider e_i, e_j for any i, j such that $i \neq j$.

What is the probability that quick sort compares e_i and e_j ?

This question, which seems to be difficult at first glance, has a surprisingly simple answer. Let us observe:

- Every element will be selected as a pivot **exactly** once.
- e_i and e_j are **not** compared, if any element **between** them gets selected as a pivot **before** e_i and e_j .

For example, suppose that $i = 7$ and $j = 12$. If e_9 is the pivot, then e_i and e_j will be separated by e_9 (**think:** why?) and will not be compared in the rest of the algorithm.

Analysis of Quick Sort

Therefore, e_i and e_j are compared if and only if either one is the first among e_i, e_{i+1}, \dots, e_j picked as a pivot.

The probability is $2/(j - i + 1)$.

Analysis of Quick Sort

Define random variable X_{ij} to be 1, if e_i and e_j are compared. Otherwise, $X_{ij} = 0$. We thus have $\Pr[X_{ij} = 1] = 2/(j - i + 1)$. That is, $E[X_{ij}] = 2/(j - i + 1)$.

Clearly, $X = \sum_{i,j} X_{ij}$. Hence:

$$\begin{aligned} E[X] &= \sum_{i,j:i < j} E[X_{ij}] = \sum_{i,j:i < j} \frac{2}{j - i + 1} \\ &= 2 \sum_{i=1}^{n-1} \sum_{j=i+1}^n \frac{1}{j - i + 1} \\ &= 2 \sum_{i=1}^{n-1} O(\log(n - i + 1)) \\ &= 2 \sum_{i=1}^{n-1} O(\log n) = O(n \log n). \end{aligned}$$

Analysis of Quick Sort

The above analysis used the following fact:

$$1 + 1/2 + 1/3 + 1/4 + \dots + 1/n = O(\log n).$$

The left-hand side is called the **harmonic series**.