

# Basic Concepts of Graphs

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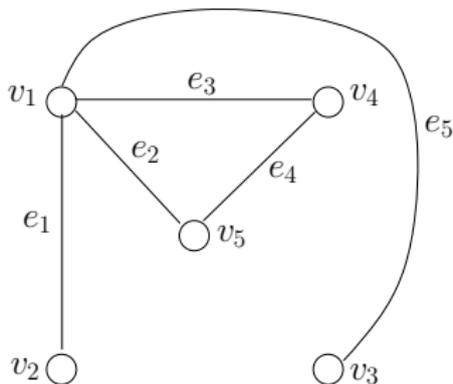
## Undirected Graphs

An **undirected graph** is a pair of  $(V, E)$  where:

- $V$  is a set of elements, each of which called a **node**.
- $E$  is a set of **unordered pairs**  $\{u, v\}$  where  $u$  and  $v$  are nodes.

A node is also called a **vertex**. Each element  $\{u, v\} \in E$  is also called an **edge**. Node  $u$  is a **neighbor** of  $v$ ; the two vertices are **adjacent** to each other.

## Example



This is an undirected graph where there are 5 vertices  $v_1, v_2, \dots, v_5$ , and 5 edges  $e_1, e_2, \dots, e_5$ .

## Directed Graphs

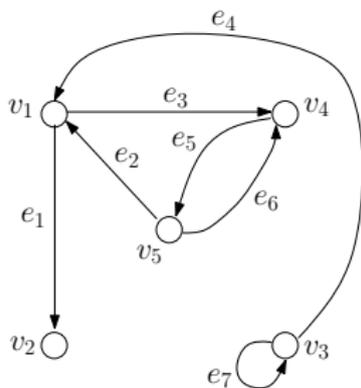
An **directed graph** is a pair of  $(V, E)$  where:

- $V$  is a set of elements, each of which called a **node**.
- $E$  is a set of pairs  $(u, v)$  where  $u$  and  $v$  are nodes in  $V$ .

A node is also called a **vertex**. Each element  $\{u, v\} \in E$  is also called an **edge**.

Each element  $(u, v) \in E$  is a **directed edge**. More specifically, it is an **outgoing** edge of  $u$  and an **incoming** edge of  $v$ . Accordingly,  $v$  is an **out-neighbor** of  $u$  and  $u$  is an **in-neighbor** of  $v$ .

## Example

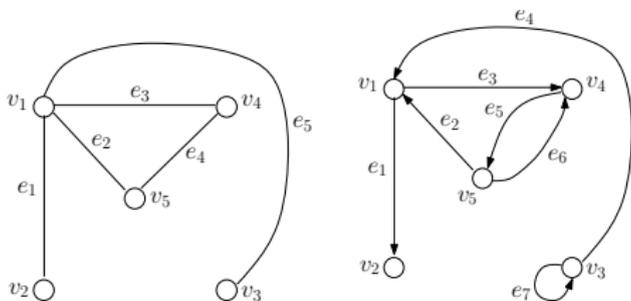


This is a directed graph  $(V, E)$  where there are 5 vertices  $v_1, v_2, \dots, v_5$ , and 7 edges  $e_1, e_2, \dots, e_7$ . Edge  $e_6$ , for instance, is an outgoing edge of  $v_5$  and an incoming edge of  $v_4$ .

## Degrees

- In an undirected graph, the **degree** of a vertex  $u$  is the number of edges of  $u$ .
- In a directed graph, the **out-degree** of a vertex  $u$  is the number of outgoing edges of  $u$ , and its **in-degree** is the number of its incoming edges.

## Example



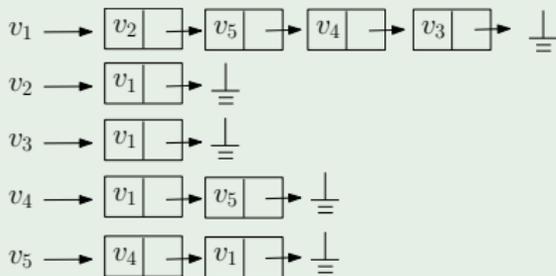
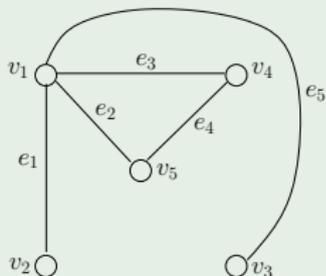
In the left graph, the degree of  $v_5$  is 2. In the right graph, the out-degree of  $v_3$  is 2 and its in-degree is 1.

Next, we discuss two common ways to store a graph: **adjacency list** and **adjacency matrix**. In both cases, we represent each vertex in  $V$  using a unique id in  $1, 2, \dots, |V|$ .

## Adjacency List – Undirected Graphs

Each vertex  $u \in V$  is associated with a linked list that enumerates all the vertices adjacent to  $u$ .

### Example

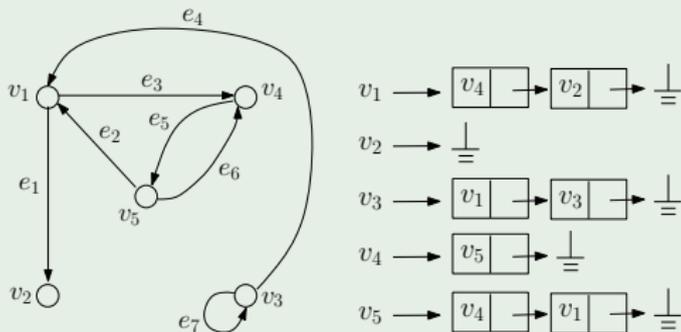


Space =  $O(|V| + |E|)$ .

## Adjacency List – Directed Graphs

Each vertex  $u \in V$  is associated with a linked list that enumerates all the out-neighbors of  $u$ .

### Example

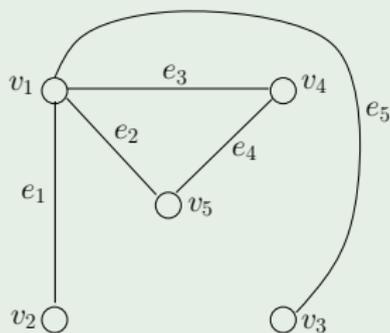


Space =  $O(|V| + |E|)$ .

## Adjacency Matrix – Undirected Graphs

A  $|V| \times |V|$  matrix  $A$  where  $A[u, v] = 1$  if  $(u, v) \in E$ , or 0 otherwise.

### Example



	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$
$v_1$	0	1	1	1	1
$v_2$	1	0	0	0	0
$v_3$	1	0	0	0	0
$v_4$	1	0	0	0	1
$v_5$	1	0	0	1	0

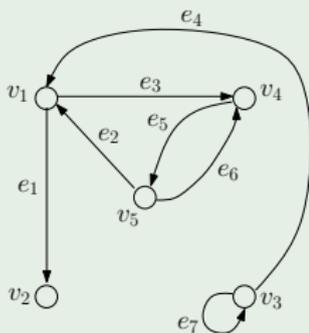
- $A$  must be symmetric.
- Space =  $O(|V|^2)$ .

**Think:** How to store  $A$  so that, for any vertices  $u, v \in V$ , we can find out if they have an edge in constant time?

## Adjacency Matrix – Directed Graphs

Defined in the same way as in the undirected case.

### Example



	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$
$v_1$	0	1	0	1	0
$v_2$	0	0	0	0	0
$v_3$	1	0	1	0	0
$v_4$	0	0	0	0	1
$v_5$	1	0	0	1	0

- $A$  may **not** be symmetric.
- Space =  $O(|V|^2)$ .